Jiří Rosický Concerning binding categories

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# CONCERNING BINDING CATEGORIES

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A category A is binding if any algebraic category can be fully embedded into it (see [1]). By an algebraic category we mean in this paper any equationally definable category of algebras with finitary operations. J. SICHLER has found in [6] a finite category C such that a category A satisfying the following conditions (0)-(6) is binding if and only if C can be fully embedded into A:

- (0) there exists a faithful functor  $U: A \to Ens$  (it means that (A, U) is a concrete category),
- (1) there are a class E of epis and a class M of monics of A such that A is a bicategory in the sense of Isbell with respect to these two classes,
- (2) U(m) is one-to-one mapping for every  $m \in M$ ,
- (3) for every object a of A and for every bijection b : Ua → x there is an isomorphism i of A such that U(i) = b,
- (4) A has and U preserves equalizers,
- (5) A is cocomplete,
- (6) if  $D: S \to A$  is a diagram and  $a \in A$  is its colimit with the colimiting cone  $\tau: D \to a$ , then

$$Ua = \bigcup_{s \in S} U(\tau_s) (UDs).$$

He has proved it in the following way. Let  $\overline{G}$  be the category of all undirected graphs and their compatible mappings. Let  $| : \overline{G} \to Ens$  be the usual forgetful functor. Denote successively by  $\underline{1}, \underline{2}, \underline{3}$  and  $\underline{4}$  the full graph without diagonal having one, two, three and four vertices. An undirected graph is 3-colourable if it has a compatible mapping into  $\underline{3}$ . Let G be the full subcategory of  $\overline{G}$  consisting of all 3-colourable graphs. The category G is binding. Let C be the full subcategory of  $\overline{G}$ determined by graphs  $\underline{1}, \underline{2}$  and  $\underline{4}$ . C is dense (left adequate) in  $G \cup {\underline{4}}$  and cogenerates  $G \cup {\underline{4}}$  because  $\underline{4}$  cogenerates G itself. J. Sichler has shown that if C can be fully embedded into a category A satisfying (0)-(6), then there exists a full embedding  $T: C \to A$  such that a left Kan extension  $L_0$  of T is a full embedding of the binding category of all connected graphs from  $G \cup \{4\}$  into A.

On the other hand, let B be a category,  $B_0$  a full subcategory of B which is small and cogenerates B, and A a cocomplete and co-well-powered category. If  $T: B_0 \to A$ is a full embedding, then beginning with a left Kan extension  $L_0$  of T we can transfinitely construct a functor  $L_*: B \to A$  extending T such that whenever  $B_0$  is dense in B and A has enough isomorphic copies of each of its objects, then  $L_*$  is a full embedding if and only if a full embedding extending T exists (see [5]).

If we take  $L_*$  instead of  $L_0$  in the previous situation, we can enlarge the class of categories tested for bindability by a small category. First, we can show that any co-well-powered category A satisfying (0), (3), (5) and (6) is tested for bindability by C, again. For instance, such categories A cover all comonadic categories. Restrictive is the condition (6). However, if we replace C by a certain small category  $C_0$ , (6) can be weakened to a condition satisfied by any algebraic category. In this way we shall solve the problem set in [6] whether there is a small category testing the bindability of any algebraic category.

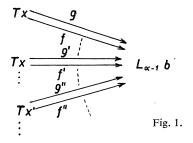
All necessary concepts of the category can be found in [2].

### 1. THE CONSTRUCTION

We shall describe the construction of  $L_*$  on objects. Let  $B_0$  be a small full subcategory of B which cogenerates B, A a cocomplete co-well-powered category and  $T: B_0 \to A$  a full embedding. Let  $b \in B$  and denote by  $P: (B_0 \downarrow b) \to B_0$  the projection of the comma category  $(B_0 \downarrow b)$  into  $B_0$ . Then  $L_0 b$  is a colimit of the functor

$$(B_0 \downarrow b) \xrightarrow{P} B_0 \xrightarrow{T} A$$
.

Suppose that we have a functor  $L_{\alpha-1}$  for an ordinal  $\alpha$ . Then  $L_{\alpha}b$  is a colimit of the following diagram.



Arrows of this diagram are all arrows of A with the domain in  $TB_0$  and the codomain  $L_{\alpha-1}b$ . Arrows  $f, g: Tx \to L_{\alpha-1}b$  have the same domain in this diagram if and only

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if  $L_{\alpha-1}(h) \cdot f = L_{\alpha-1}(h) \cdot g$  for every arrow  $h : b \to y$  and every  $y \in B_0$ . Denote by  $\lambda_b^{\alpha-1,\alpha}$  the component of the colimiting cone with the domain  $L_{\alpha-1}b$ . Further, let  $\alpha$  be limit and consider the diagram

$$(**) L_0 b \xrightarrow{\lambda_b^{0,1}} L_1 b \xrightarrow{\lambda_b^{1,2}} \dots L_{\beta} b \xrightarrow{\lambda_b^{\beta,\beta+1}} L_{\beta+1} b \xrightarrow{\lambda_b^{\beta+1,\beta+2}} \dots$$

having objects  $L_{\beta}b$  and arrows  $\lambda_{b}^{\beta,\beta+1}$  for  $\beta < \alpha$ . Then  $L_{\alpha}b$  is defined to be a colimit of this diagram and by  $\lambda_{b}^{\beta,\alpha}: L_{\beta}b \to L_{\alpha}b$  we denote components of the colimiting cone. This process stops at some ordinal  $\gamma$  and we put  $L_{*}b = L_{\gamma}b$ . We get functors  $L_{\beta}$ and  $L_{*}$  extending T. Moreover, our  $\lambda$ 's determine natural transformations  $\lambda^{\beta,\alpha}$ :  $:L_{\beta} \to L_{\alpha}$  and  $\lambda^{\alpha}: L_{\alpha} \to L_{*}$  for any  $\beta < \alpha$ . It holds  $\lambda^{\beta,\alpha} \cdot \lambda^{\delta,\beta} = \lambda^{\delta,\alpha}$  for any  $\delta < \beta < \alpha$ and therefore  $L_{\alpha}b$  is a colimit of a slight modified diagram (\*\*'), which we obtain from (\*\*) taking all  $\lambda_{b}^{\delta,\beta}, \delta < \beta < \alpha$  as arrows.

The detailed description of this construction can be found in [5].

**Lemma 1.** Let  $x \in B_0$ ,  $b \in B$  and  $f, g : L_*x \to L_*b$  in A such that  $L_*(h)f = L_*(h)g$ for any  $h : b \to y$  and any  $y \in B_0$ . Then f = g.

Proof immediately follows from the construction of  $L_*b$ .

## **Lemma 2.** In addition, let $B_0$ generate B. Then $L_*$ is faithful.

Proof. Let  $x \in B_0$ ,  $b \in B$  and  $f \neq g : x \to b$ . Since  $B_0$  cogenerates B, there exist  $y \in B_0$  and  $h : b \to y$  such that  $hf \neq hg$ . Since T is faithful,  $L_*(hf) \neq L_*(hg)$  and therefore  $L_*(f) \neq L_*(g)$ . Now, faithfulness of  $L_*$  follows similarly from the fact that  $B_0$  generates B (see [5] Prop. 1).

Now, we are going to give some sufficient conditions for  $L_*$  to be a full embedding. A subcategory  $B_0$  of a concrete category (B, | |) projectively generates B if for any  $b, c \in B$  and any mapping  $f: |b| \to |c|, f = |f_1|$  for an arrow  $f_1: b \to c$  of B if and only if for every  $x \in B_0$  and every arrow  $h: c \to x$  of B there is an arrow  $h': b \to x$  of B such that |h'| = |h|. f. Choose the following classes of small categories:  $\mathscr{C}_1$  is the class of all well-ordered sets without the greatest element taken as categories,  $\mathscr{C}_2$  consists of all connected small categories S containing an object t such that any non-identity arrow of S has the codomain t and finally  $\mathscr{C}_3(B_0, B)$  consists of all comma categories  $(B_0 \downarrow b)$  for  $b \in B$ . Put  $\mathscr{C}(B_0, B) = \mathscr{C}_1 \cup \mathscr{C}_2 \cup \mathscr{C}_3(B_0, B)$ . Clearly any colimit needed in the construction of  $L_*$  is a colimit of a diagram  $D: S \to A$  with  $S \in \mathscr{C}(B_0, B)$ . Namely, (\*) for  $S \in \mathscr{C}_2$  and (\*\*') for  $S \in \mathscr{C}_1$ .

The crucial part of the following proof, the proof of fulness of  $L_*$  with respect to arrows of A having a domain in  $L_*B_0$ , is a modification of Lemma 4 from [6].

**Theorem 1.** Let (B, | |) be a concrete category and  $B_0$  a small dense full subcategory of B which cogenerates and projectively generates B. Let  $B_0$  contains an object e such that card |e| = 1 and for any  $x \in B$  and any mapping  $u' : |e| \to |x|$  there exists  $u : e \to x$  in  $B_0$  with |u| = u'. For any  $z \in B_0$ ,  $b \in B$  and  $g : z \to b$  in B let there exist  $d \in B_0$  with the following properties:

- a) d is a cogenerator of B,
- b) for any permutation s' of |d| interchanging two elements of |d| there exists  $s: d \rightarrow d$  in B with |s| = s',
- c) there is an  $h_0: b \to d$  in B such that for any  $h: b \to d$  in B, an  $s: d \to d$  in B with  $hg = sh_0g$  can be found,
- d) card  $(|d| |h_0g|(|z|)) > 1$ .

Let (A, U) be a co-well-powered concrete category having colimits of functors  $D: S \to A$  for  $S \in \mathscr{C}(B_0, B)$ , satisfying (6) for these colimits and fulfilling the condition (3). Let  $T: B_0 \to A$  be a full embedding.

Then  $L_*: B \to A$  is a full embedding.

Proof. By the previous lemma and Corollary 2 from [5] it suffices to prove that for any  $y \in B_0$ ,  $b \in B$  and  $f: L_*y \to L_*b$  in A there is an arrow  $f': y \to b$  in B with  $L_*(f') = f$ .

Denote by  $u_{x,i}$  the arrow  $u_{x,i}: e \to x$  of B for which  $|u_{x,i}|(|e|) = \{i\}$ , where  $x \in B$  and  $i \in |x|$ . Let us have a  $d \in B_0$  satisfying a)-d). Let  $p \in UL_*e$  such that  $UL_*(u_{d,i})(p) = UL_*(u_{d,j})(p)$  for some  $i, j \in |d|$ ,  $i \neq j$ . Take  $k \in |d|$ ,  $i \neq k \neq j$  and  $s: d \to d$  such that |s| is the permutation of |d| interchanging i and k. We get  $UL_*(u_{d,k})(p) = UL_*(su_{d,i})(p) = UL_*(su_{d,j})(p) = UL_*(su_{d,j})(p)$ . Thus there exists  $p \in UL_*e$  such that  $UL_*(u_{d,i})(p) = UL_*(su_{d,j})(p) = UL_*(u_{d,j})(p)$ . Thus there exists  $p \in UL_*e$  such that  $UL_*(u_{d,i})(p) = UL_*(u_{d,j})(p)$  for some  $x \in B_0$  and  $i, j \in |x|$ ,  $i \neq j$ . Suppose that  $UL_*(u_{x,i})(p) = UL_*(u_{x,j})(p)$  for some  $x \in B_0$  and  $i, j \in |x|$ ,  $i \neq j$ . Since d is a cogenerator, there exists  $h: x \to d$  in B such that  $u_{d,/h/(i)} = hu_{x,i} \neq hu_{x,j} = u_{d,/h/(i)}$ . Further,  $UL_*(u_{d,/h/(i)})(p) = UL_*(u_{d,/h/(i)})$ , which is a contradiction.

Let  $b \in B$  and  $f: L_*e \to L_*b$ . Since  $L_*b$  is defined by colimits of functors  $D: S \to A$ , where  $S \in \mathscr{C}(B_0, B)$  and any  $\lambda^{\beta,\alpha}$  is a natural transformation, (6) enables us to find  $z \in B_0$ ,  $g: z \to b$  in B and  $q \in UL_*z$  such that  $U(f)(p) = UL_*(g)(q)$ . Take d for this g. We shall denote  $u_{d,i}$  for d just taken briefly by  $u_i$ .

Consider  $h_0$  from c). We can find  $u_k$  such that  $L_*(h_0) \cdot f = L_*(u_k)$ . Suppose that  $k \notin |h_0g|(|z|)$ . Let  $s: d \to d$  be an arrow in B such that |s| is the permutation of |d| interchanging k with an element  $i \in |h_0g|(|z|)$ . Then there exists  $u_{z,n}: e \to z$  in B with  $u_k = sh_0gu_{z,n}$ . It holds  $UL_*(u_k)(p) = U(L_*(h_0)f)(p) = UL_*(h_0g)(q)$ . By d) and b) there exists an  $s': d \to d$  such that |s'| is the permutation of |d| interchanging i with some  $j \in |d| - (|h_0g|(|z|) \cup \{k\})$ . We have  $UL_*(u_k)(p) = UL_*(s'u_k)(p) = UL_*(s'h_0g)(q)$  and therefore  $UL_*(u_i)(p) = UL_*(su_k)(p) = UL_*(s'h_0g)(q) = UL_*(u_k)(p)$ , which is a contradiction. Thus  $k \in |h_0g|(|z|)$  and there exists a  $u_{z,n}: e \to z$  such that  $u_k = h_0gu_{z,n}$ . Put  $f' = gu_{z,n}$ .

Suppose that there exists  $x \in B_0$  and  $h: b \to x$  in B such that  $L_*(h)f \neq L_*(hf')$ . Let  $u: e \to x$  with  $L_*(h)f = L_*(u)$ . There is  $h': x \to d$  with the property  $h'u \neq h'hf'$  and  $s: d \to d$  such that  $h'hg = sh_0g$  (see c)). Thus  $h'hf = sh_0f'$ . Hence  $UL_*(su_k)(p) = UL_*(sh_0f')(p) = UL_*(h'hf')(p) + UL_*(h'u)(p) = U(L_*(h'h)f)(p) = UL_*(h'hg)(q) = UL_*(sh_0g)(q) = UL_*(su_k)(p), \text{ a contradiction. Therefore } L_*(h)f = L_*(h) L_*(f') \text{ for any } x \in B_0 \text{ and any } h : b \to x \text{ and by Lemma } 1, f = L_*(f').$ 

Now, let  $y \in B_0$ ,  $b \in B$  and let  $f: L_*y \to L_*b$  be an arrow in A. Define  $\overline{f}: |y| \to |b|$ by  $fL_*(u_{y,i}) = L_*(u_{b,\overline{f}(i)})$  for  $i \in |y|$ . Let  $x \in B_0$  and  $h: b \to x$ . Then  $L_*(h) f = L_*(t)$ for some  $t: y \to x$  and we have  $L_*(tu_{y,i}) = L_*(h) fL_*(u_{y,i}) = L_*(hu_{b,\overline{f}(i)})$ . Thus  $|tu_{y,i}| = |h|\overline{f}|u_{y,i}|$  for any  $i \in |y|$ . Hence  $|h|\overline{f} = |t|$  and  $\overline{f} = |f'|$  for an arrow  $f': y \to b$  because  $B_0$  projectively generates B. Moreover,  $L_*(hf') = L_*(t) =$  $= L_*(h) f$  and Lemma 1 yields  $L_*(f') = f$ . The proof is complete.

Note. Using p from the previous proof we may define a natural monotransformation  $\alpha : | | \rightarrow UL_*$  by  $\alpha_b(i) = UL_*(u_{b,i})(p)$ . If we want to avoid the axiom of choice for classes, which is used in the proof of Corollary 2 of [5], we can suppose that  $|t| = id_{|b|}$  implies  $t = id_b$  for any isomorphism  $t : b \rightarrow b$  in B and apply Lemma 1.5 of [4].

## 2. TESTING CATEGORIES

It is easy to see that  $B = G \cup \{\underline{4}\}$  and  $B_0 = C$  fulfil all suppositions of Theorem 1. Indeed,  $e = \underline{1}, \underline{4}$  is the only d and  $\underline{4}$  projectively generates  $G \cup \{\underline{4}\}$  by Lemma 1 of [6].

**Theorem 2.** A co-well-powered category A satisfying (0), (3), (5) and (6) is binding if and only if C can be fully embedded into it.

We shall define a small category  $C_0$  testing bindability of any algebraic category. Let  $\underline{\aleph}_0$  be the full graph with the diagonal having countably many vertices. Let  $C_0$  be the full subcategory of  $\overline{G}$  containing  $\underline{\aleph}_0$ , 4 and all graphs decomposing into a finite number of components of the form 1 or 2.

**Lemma 3.** A comma category  $(C_0 \downarrow x)$  is filtered for any graph  $x \in G$ .

Proof. Let  $x \in G$  and put  $S = (C_0 \downarrow x)$ . We have to prove that any diagram in S of the form

is a base for a cone.

Let  $f, f' \in S$ . It means that  $f: z \to x$  and  $f': z' \to x$  are compatible mappings and  $z, z' \in C_0$ . Since x is 3-colourable,  $4 \neq z, z' \neq \underline{\aleph}_0$ . Let z" be the coproduct of z and z' in G with injections  $i: z \to z'', i': z' \to z''$  and  $f'': z'' \to x$  the unique arrow of G such that f''i = f and f''i' = f'. Since  $|z''| = |z| \cup |z'|$  and any edge of z" is an edge of z or z', we have  $z'' \in C_0$ . Therefore  $f'' \in S$  and  $i : f \to f''$ ,  $i' : f' \to f''$  are arrows of S.

Let  $f, g \in S, g : z \to x$  and let  $i, j : f \to g$  be a parallel pair of arrows of S. Define a compatible mapping  $k : z \to z$  such that two components of z have the same image in k if and only if they have the same image in g. There is an  $h : z \to x$  such that g = hk, which means that  $k : g \to h$  is an arrow in S and moreover, ki = kj holds.

**Theorem 3.** Let A be a co-well-powered category satisfying (0), (3) which has colimits of functors  $D: S \to A$  for S filtered or  $S \in \mathcal{C}_2$  and fulfils (6) for these colimits. Then A is binding if and only if  $C_0$  can be fully embedded into it.

Proof. Put  $B = G \cup \{4, \underline{\aleph}_0\}$  and  $B_0 = C_0$ . Then all suppositions of Theorem 1 are satisfied. Namely, e = 1 and  $\aleph_0$  is the only *d*. Since any category belonging to  $\mathscr{C}_1$  is filtered and by Lemma 3 the same holds for  $\mathscr{C}_3(B_0, B)$ , Theorem 3 follows from Theorem 1.

**Lemma 4.** Let (A, U) be a concrete category having kernel pairs and let Uf be epi for any coequalizer f in A. Then A satisfies (6) for any  $S \in \mathscr{C}_2$ .

Proof. Let  $S \in \mathscr{C}_2$  and let  $a \in A$  be a colimit of a functor  $D : S \to A$  with the colimiting cone  $\tau$ . Denote by f the component of  $\tau$  with the domain t. Let  $f_1, f_2 : a_0 \to t$  be a kernel pair of f. We shall prove that f is a coequalizer of  $f_1$  and  $f_2$ . Suppose that  $gf_1 = gf_2$  for an arrow g in A. Let  $h_1, h_2 : s \to t$  be a parallel pair of arrows of S. Since  $f \cdot D(h_1) = f \cdot D(h_2)$ , there is a unique arrow  $k : D(s) \to a_0$  in A such that  $f_i k = D(h_i)$  for i = 1, 2. Therefore  $g \cdot D(h_1) = g \cdot D(h_2)$  and g determines a cone from the base D. Thus, there is a unique arrow k' in A with k'f = g. Hence f is a coequalizer and Uf is epi. But it means that (6) is satisfied.

**Corollary 1.** An algebraic category A is binding if and only if  $C_c$  can be fully embedded into it.

Proof. The usual forgetful functor  $U: A \rightarrow Ens$  preserves filtered colimits (see [2, p. 209]). Further, A is cocomplete, complete and Uf is epi for any coequalizer f in A. The last fact can be found in [3], Lemma 6. The result follows from Lemma 4 and Theorem 3.

The infiniteness of  $C_0$  plays no role because we can take instead of  $C_0$  the full subcategory C' of  $\overline{G}$  having objects  $\underline{1}, \underline{4}, \underline{\aleph}_0$  and the graph having countably many copies of  $\underline{1}$  and countably many copies of  $\underline{2}$  as components. In this case we have to use only the graphs from G having at least one edge instead of the whole G. After computation that C' is dense in  $G \cup {\underline{4}, \underline{\aleph}_0}$  and that  $(C' \downarrow x)$  is filtered for any  $x \in G$  having at least one edge, we obtain that this four-object category C' tests bindability of any algebraic category. It seems to be an interesting problem to find such a testing category with the smallest possible number of objects. The method presented does not work for categories of algebras of arbitrary arities because categories from  $\mathscr{C}_1$  needed for the construction of  $L_*$  are not m-filtered for a cardinal  $\mathfrak{m} > \aleph_0$ .

Added in proof. This difficulty is avoided and further results are given in the following author's papers: On extensions of full embeddings and binding categories (Comm. Math. Univ. Carol. 15 (1974), 631-653) and Codensity and binding categories (to appear in Comm. Math. Univ. Carol.).

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