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## THE DISTANCE BETWEEN A GRAPH AND ITS COMPLEMENT

BOHDAN ZELINKA, Liberec

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In [3] the distance between isomorphism classes of graphs was introduced. Here we shall investigate this distance between a graph and its complement.

An isomorphism class of graphs is the class of all graphs which are isomorphic to a given graph.

Now let *n* be a positive integer and let  $\mathscr{G}_n$  be the set of all isomorphism classes of graphs with *n* vertices. Let  $\mathfrak{G}_1 \in \mathscr{G}_n$ ,  $\mathfrak{G}_2 \in \mathscr{G}_n$ . Let *p* be the maximum number of vertices of a graph which is isomorphic simultaneously to an induced subgraph of a graph  $G_1 \in \mathfrak{G}_1$  and to an induced subgraph of a graph  $G_2 \in \mathfrak{G}_2$ . We put  $\delta(\mathfrak{G}_1, \mathfrak{G}_2) = n - p$  and call this number the distance between the isomorphism classes  $\mathfrak{G}_1, \mathfrak{G}_2$ .

For the sake of brevity we shall (not quite accurately) speak about the distance between graphs instead of the distance between isomorphism classes of graphs. By the distance  $\delta(G_1, G_2)$  of the graphs  $G_1, G_2$  (with the same number of vertices) we mean the distance  $\delta(\mathfrak{G}_1, \mathfrak{G}_2)$  of the isomorphism classes  $\mathfrak{G}_1, \mathfrak{G}_2$  such that  $G_1 \in \mathfrak{G}_1$ ,  $G_2 \in \mathfrak{G}_2$ . By a common induced subgraph of  $G_1$  and  $G_2$  we shall mean a graph which is isomorphic simultaneously to an induced subgraph of  $G_1$  and to a an induced subgraph of  $G_2$ .

In this paper we shall study the distance  $\delta(G, \overline{G})$  between a graph G and its complement  $\overline{G}$ . As the complement  $\overline{G}$  is uniquely determined by the graph G, the distance  $\delta(G, \overline{G})$  is a numerical invariant of G; we denote it by  $\overline{\delta}(G)$ .

We shall consider only finite undirected graphs without loops and multiple edges. Obviously  $\overline{\delta}(G) = 0$  if and only if G is a self-complementary graph, i.e. a graph isomorphic to its own complement. These graphs were studied by G. Ringel [1] and H. Sachs [2]; these authors have (mutually independently) proved that a self-complementary graph with n vertices exists if and only if  $n \equiv 0 \pmod{4}$  or  $n \equiv 1 \pmod{4}$ .

**Theorem 1.** Let n be an integer,  $n \ge 2$ . If  $n \equiv 0 \pmod{4}$  or  $n \equiv 1 \pmod{4}$ , then for any graph G with n vertices

$$0 \leq \delta(G) \leq n-1$$

holds and for any integer d such that  $0 \le d \le n - 1$  there exists a graph G with n vertices such that  $\overline{\delta}(G) = d$ . If  $n \equiv 2 \pmod{4}$  or  $n \equiv 3 \pmod{4}$ , then for any graph

G with n vertices

$$1 \leq \bar{\delta}(G) \leq n-1$$

holds and for any integer d such that  $1 \leq d \leq n - 1$  there exists a graph G with n vertices such that  $\overline{\delta}(G) = d$ .

Proof. As it was mentioned above, for  $n \equiv 0 \pmod{4}$  and for  $n \equiv 1 \pmod{4}$ there exist self-complementary graphs with *n* vertices, i.e. graphs *G* for which  $\overline{\delta}(G) =$ = 0. For  $n \equiv 2 \pmod{4}$  and for  $n \equiv 3 \pmod{4}$  such graphs do not exist, but in [4] it was proved that there exist almost self-complementary graphs with *n* vertices. An almost self-complementary graph is a graph *G* with the property that it can be transformed into a graph isomorphic to  $\overline{G}$  by adding or deleting one edge. Thus consider such an almost self-complementary graph *G* with *n* vertices. Let *e* be the edge by whose adding or deleting from *G* a graph isomorphic to  $\overline{G}$  is obtained, let *u* be one of its end vertices. Then the graph obtained from *G* by deleting *u* is an induced subgraph of a graph isomorphic to  $\overline{G}$  and thus  $\overline{\delta}(G) = \delta(G, \overline{G}) = 1$ . This gives the lower bound. Any non-empty graph contains a subgraph consisting of one isolated vertex, hence  $\overline{\delta}(G) \leq n - 1$ .

Now let an integer d be given,  $0 \le d \le n - 1$ . The case d = 0 was yet considered; thus suppose  $1 \le d \le n - 1$ . If  $n - d \equiv 0 \pmod{4}$  or  $n - d \equiv 1 \pmod{4}$ , we take sets V,  $V_0$  of vertices such that  $V_0 \subset V$ ,  $|V_0| = n - d$ , |V| = n. We construct a self-complementary graph  $G_0$  on  $V_0$ . Now the graph G is the graph obtained from  $G_0$ by adding the vertices of  $V - V_0$  as isolated vertices. The subgraphs of G and  $\overline{G}$ induced by  $V_0$  are both isomorphic to  $G_0$ . Any subgraph of G having more than n - dvertices contains at least one isolated vertex, while such a subgraph of  $\overline{G}$  has not. Therefore  $\delta(G, \overline{G}) = n - (n - d) = d$ . If  $n - d \equiv 2 \pmod{4}$ , then we take the vertex sets  $V_0$ , V such that  $V_0 \subset V$ ,  $|V_0| = n - d + 1$ , |V| = n, construct an almost self-complementary graph  $G_0$  on  $V_0$  and proceed further as in the preceding case. If  $n - d \equiv 3 \pmod{4}$ , then we take again  $V_0$  and V so that  $V_0 \subset V$ ,  $|V_0| = n - d + 1$ , |V| = n, construct a self-complementary graph on  $V_0$  and add an edge to it to obtain  $G_0$ ; then we proceed as in the preceding case.

Now we shall investigate graphs with the property that all of their connected components are cliques. Their complements are the so-called complete multipartite graphs.

**Theorem 2.** Let G be a graph with n vertices having q connected components, all of which are cliques, let r be the maximum number of vertices of a connected component of G. Then  $\overline{S}(G) = r = \min(q, r)$ 

$$\delta(G) = n - \min\{q, r\}.$$

Proof. Denote  $s = \min \{q, r\}$ . First suppose s = q. Then  $s \leq r$  and both G and  $\overline{G}$  contain subgraphs which are complete graphs with s vertices. Now consider a subgraph H of G with more than s vertices. All connected components of H are complete graphs and at least one of them has more than one vertex. If H is a complete graph, then no induced subgraph of  $\overline{G}$  is isomorphic to H, because the largest

clique in  $\overline{G}$  has s vertices. If H contains at least two connected components, then also no induced subgraph of  $\overline{G}$  is isomorphic to it, because each disconnected induced subgraph of  $\overline{G}$  consists of isolated vertices. Hence  $\overline{\delta}(G) = n - s$ . Now let s = r. Then  $s \leq q$  and both G and  $\overline{G}$  contain induced subgraphs consisting of s isolated vertices. Now consider a subgraph H of G with more than s vertices. Then this graph is disconnected. If it contains an edge, it is isomorphic to no induced subgraph of  $\overline{G}$ as it was mentioned above. If H consists of isolated vertices, it is also isomorphic to no induced subgraph of  $\overline{G}$ , because the maximum number of vertices of an independent set in  $\overline{G}$  is s. Again  $\overline{\delta}(G) = n - s$ .

**Theorem 3.** For a graph G with n vertices  $\overline{\delta}(G) = n - 1$  if and only if G is a complete graph or consists of isolated vertices.

Proof. The sufficiency follows from Theorem 2, where q = 1, r = n or q = n, r = 1. The necessity follows from the fact that any graph which neither is complete, nor consists of isolated vertices contains both possible types of two-vertex subgraphs.

**Theorem 4.** For a graph G with n vertices  $\overline{\delta}(G) = n - 2$  if and only if G is a graph of someone of the following types:

- (a) complete bipartite graph;
- (b) graph consisting of two connected components being cliques;
- (c) graph consisting of connected components being cliques at which the maximum number of vertices of a clique is 2;
- (d) the complement of a graph of the type (c).

Proof. The graphs of the types (b) and (c) are graphs described in Theorem 2 for q = 2 or r = 2, the graphs of the types (a) and (d) are their complements. This implies the sufficiency. Now let G be a graph which does not belong to the types (a), (b), (c), (d); then evidently  $\overline{G}$  also does not belong to them. Suppose that all connected components of G are cliques. If each of them consists of one vertex or there exists only one connected component, then Theorem 3 holds for G. Otherwise there are at least three connected components and at least one of them has at least three vertices. Then both G and  $\overline{G}$  contain triangles and  $\delta(G) \leq n - 3$ . If all connected components of  $\overline{G}$  are cliques, the proof is analogous. Finally, if both G and  $\overline{G}$ contain a connected component which is not a complete graph, then they both contain an induced subgraph being a path of the length 2 and again  $\overline{\delta}(G) \leq n - 3$ .

At the end we shall study paths and circuits. By  $P_n$  we denote the path of the length n, i.e. with n edges and n + 1 vertices. By  $C_n$  we denote the circuit of the length n.

**Theorem 5.** For the paths there is

$$\begin{split} \bar{\delta}(P_1) &= 1 ,\\ \bar{\delta}(P_2) &= 1 ,\\ \bar{\delta}(P_3) &= 0 ,\\ \bar{\delta}(P_n) &= n - 4 \quad \text{for} \quad n \geq 4 . \end{split}$$

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Proof. The assertions for  $P_1$  and  $P_2$  are evident. The path  $P_3$  is a self-complementary graph. If  $n \ge 4$ , then  $P_n$  contains an induced subgraph isomorphic to  $P_3$ ; the subgraph induced by the same vertex set in  $\overline{P}_n$  is also isomorphic to  $P_3$ . The graph  $P_3$  has four vertices and thus  $\delta(P_n) \le n - 4$ . On the other hand, each induced subgraph of  $P_n$  with at least five vertices contains an independent set with three vertices; hence the subgraph of  $\overline{P}_n$  induced by the same set contains a triangle, while  $P_n$  contains no triangle. This implies  $\delta(P_n) = n - 4$ .

**Theorem 6.** For the circuits there is

$$\overline{\delta}(C_3) = 1 , \overline{\delta}(C_4) = 2 , \overline{\delta}(C_5) = 0 , \overline{\delta}(C_n) = n - 4 \text{ for } n \ge 6 .$$

Proof. The assertions for  $C_3$  and  $C_4$  follow from Theorem 2. The circuit  $C_5$  is a self-complementary graph. The assertion for  $n \ge 6$  can be proved in the same way as the assertion for  $n \ge 4$  in Theorem 5.

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Author's address: 461 17 Liberec 1, Studentská 1292, Czechoslovakia (katedra tváření a plastů VŠST).