Jaroslav Guričan; Pavol Zlatoš Archimedean and geodetical biequivalences

Commentationes Mathematicae Universitatis Carolinae, Vol. 26 (1985), No. 4, 675--698

Persistent URL: http://dml.cz/dmlcz/106407

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COMMENTATIONES MATHEMATICAE UNIVERSITATIS CAROLINAE

26,4 (1985)

ARCHIMEDEAN AND GEODETICAL BIEQUIVALENCES Jaroslav GURIČAN and Pavol ZLATOŠ

Abstract: This paper contributes to the topological problematics in the AST. The central role in it is due to the concept of a biequivalence introduced in [G-Z 1]. A metrization theorem for biequivalences is established. Two properties of biequivalences bearing upon the connectedness of galaxies named in the title are formulated and characterized. The notions of a path and a motion of point appear as powerful tools in formulations and proofs of the results.

Key words: Biequivalence, path, motion, compact, connected, metric, galaxy, Archimedean, direct, geodetical.

Classification: Primary 54J05 Secondary 54D05, 54E35

This paper is a direct continuation of [G-Z 1] contributing to the topological problematics in the AST. The central role in it is played by the concept of a biequivalence introduced in [G-Z 1]. The article joins results of two areas of "biequivalence problematics" originally occurring rather independent.

The first part is devoted to the characterization of biequivalences $\langle \pm, \leftrightarrow \rangle$ such that for each mean bound R and each pair $x \leftrightarrow y$ there is a finite R-path from x to y (Archimedean biequivalences). The formulation and the proof of the result itself are preceeded by a section dealing with paths and metions.

The second part of our work was iniciated by the unexpectedly easy (using the results of [M 2]) proof of the metrization theorem for arbitrary biequivalences. From this theorem some results, analo--675 - gous to those from the classical topology, on embeddings of a biequivalence with a set domain u into the linear space RN^U endowed with the componentwise biequivalence easily follow. If one would like to generalize these results to arbitrary biequivalences, he will find unavoidable to extend suitably the field of rational numbers. That's why we sketch the construction of hyperreal numbers in the AST.

Keeping the fact that every biequivalence is induced by some metric, in mind, there arise several questions under what conditions such a metric abundant in some further useful properties can be found. One particular problem of this type is solved in the paper. Namely, biequivalences which can be induced by a metric H such that each pair of accessible points can be joined by a direct motion with respect to H are characterized (geodetical biequivalences).

Both the notions of Archimedean and geodetical biequivalences illustrate the "restriction principle to galaxies" mentioned in [G-Z 1]. Via the concept of a motion of point they bear upon some questions concerning the connectedness of galaxies of a biequivalence, as well.

The reader is assumed to be acquainted with [V] and [G-Z 1]. Most of the notions and results from these two sources will be used even without any explicit referring to them.

1. Paths, motions and connectedness

Z denotes the set-theoretically definable class of all integers and FZ stands for finite integers. Variables $\boldsymbol{\omega}, \boldsymbol{\beta}, \boldsymbol{\gamma}, \boldsymbol{\lambda}, \boldsymbol{\mu}, \boldsymbol{\nu}, \ldots$ (k,m,n) are used sometimes for arbitrary (finite) integers, not just natural numbers. The interval of integers between $\boldsymbol{\mu}, \boldsymbol{\nu}$ is denoted by $[\boldsymbol{\mu}, \boldsymbol{\nu}] = \{\lambda \boldsymbol{\varepsilon} Z; \boldsymbol{\mu} \boldsymbol{\epsilon} \lambda \boldsymbol{\epsilon} \boldsymbol{\nu}\}$. In particular $[\boldsymbol{\mu}, \boldsymbol{\nu}] = \emptyset$ - 676 - if $\mu > \nu$, $[\mu, \mu] = \{\mu\}$ and $\nu = [0, \nu - 1]$ for each $\nu \in \mathbb{N}$.

In the whole paper $\langle \doteq, \leftrightarrow \rangle$ denotes the usual biequivalence on the class RN of all rational numbers (see [G-Z 1] Example 3). For any set $u \langle \doteq^{u}, \leftrightarrow^{u} \rangle$ denotes the biequivalence on $RN^{u} = \{f; dom(f) = u \ddagger rng(f) \le RN\}$ arising from $\langle \doteq, \leftrightarrow \rangle$ componentwise ([G-Z 1] Example 5). For rational numbers a,b we put $a \le b \equiv a \le b \lor a \doteq b$ and $a \lt b \equiv a \le b \And a \neq b$. The formulation of the basic properties of the relations \le and $\lt \cdot$ is left to the reader.

Let us record a result for the future.

<u>Lemma</u>. Let R be a π -relation and $u \leq \text{dom}(R)$ be a set. Then there is a set function $f \subseteq R$ such that dom(f) = u.

<u>Proof.</u> If R is set-theoretically definable, the statement can be easily proved by induction. Let $\{R_n; n \in FN\}$ be a decreasing sequence of set-theoretically definable relations whose intersection is R. For each n there is a function $f_n \subseteq R_n$ with domain u. Then the result follows by the axiom of prolongation.

Let R be an arbitrary relation. A (set) function p such that dom(p) = $[\eta, \vartheta]$ is a nonvoid interval of integers is called an R-path provided for each $\alpha \in [\eta, \vartheta - 1]$ holds $\langle p(\alpha), p(\alpha+1) \rangle \in \mathbb{R}$. Then the set rng(p) is called the trace of p. If $x = p(\eta)$ and $p(\vartheta) = y$ then p is called an R-path from x to y. In most cases the domains of the paths considered will be of form $[0, \vartheta] = \vartheta + 1$ where $\vartheta \in \mathbb{N}$; in such a case p will be called an R-path in the time ϑ . Thus $\langle x, y \rangle \in \mathbb{R}^{\vartheta}$ iff there is an R-path from x to y in the time ϑ .

If \pm is a π -equivalence then any (\pm) -path is called a motion of point in \pm . If $\{B_n; n \in \mathbb{N}\}$ is a generating sequen-- 677 - ce of \pm then obviously p is a motion of point in \pm iff p is an R_n-path for each n. We will frequently say "a motion" instead of "a motion of point in \pm ", mainly in the case when the \approx -equivalence \pm will be clear from the context.

From the results in [V] it follows directly:

<u>Theorem 1</u>. Let [±] be a X-equivalence and u be a set. The following conditions are equivalent:

- (1) for each nonempty proper subset v of u there are two points $x \in v$, $y \in u - v$ such that $x \stackrel{+}{=} y$;
- (2) there is a motion p such that u = rng(p);
- (3) for all $x, y \in u$ there is a motion p from x to y such that $rng(p) \subseteq u$.

According to Theorem 1 we addopt the following definition: A class X is connected in the \mathscr{X} -equivalence $\stackrel{\pm}{=}$ if for all x, y $\in X$ there is a motion p from x to y such that $rng(p) \subseteq X$.

<u>Theorem 2</u>. Let \pm be a x-equivalence and X be a class. If X is connected then Fig(X) is also connected. If X is a x-class and Fig(X) is connected then X is connected, as well.

<u>Proof</u>. Let X be connected and $a^{\pm}x$, $b^{\pm}y$ where $x, y \in X$. If p is a motion from x to y in the time $\sqrt[3]{}$ such that $\operatorname{rng}(p) \subseteq X$, then $q = p \cup \{\langle a, -1 \rangle, \langle b, \sqrt[3]{+1} \rangle\}$ is a motion from a to b and $\operatorname{rng}(q) \subseteq \operatorname{Fig}(X)$. Now, let X be a x-class with connected figure and $x, y \in X$. There is a motion p from x to y in the time $\sqrt[3]{}$ such that $\operatorname{rng}(p) \subseteq \operatorname{Fig}(X)$. Put

 $R = \left\{ \langle \mathbf{x}, \mathbf{0} \rangle, \langle \mathbf{y}, \mathbf{\tilde{v}} \rangle \right\} \cup \left\{ \langle \mathbf{z}, \mathbf{\alpha} \rangle; 0 < \mathbf{\alpha} < \mathbf{\tilde{v}} \} \mathbf{z} \in \mathbb{I} \ \mathbf{z} \neq \mathbf{p}(\mathbf{\alpha}) \right\}.$ Then R is a X-relation and $[0, \mathbf{\tilde{v}}] = \operatorname{dom}(\mathbf{R})$. The choice function $\mathbf{f} \in \mathbf{R}$ with domain $[0, \mathbf{\tilde{v}}]$ existing by the virtue of the Lemma is a motion from x to y and $\operatorname{rng}(\mathbf{f}) \leq \mathbb{X}$. - 678 - Each motion p in the X-equivalence \pm induces two X-equivalences on its domain. The first one seems to be prima facie more natural: $\alpha \stackrel{\pm}{\leftarrow} \beta = p(\alpha) \stackrel{\pm}{=} p(\beta)$.

<u>Theorem 3</u>. Let p be a motion in \ddagger . The \mathscr{F} -equivalence $\ddagger_{(P)}$ is compact iff rng(p) is compact in \ddagger .

<u>Proof</u>. Clearly, $\stackrel{\pm}{(p)}$ is a \Re -equivalence. Let $\stackrel{\pm}{=}$ be compact and $u \subseteq \operatorname{rng}(p)$ be infinite. If for all α , $\beta \in p^{-1} = v$ $p(\alpha) \neq p(\beta)$ implied $p(\alpha) \not\equiv p(\beta)$, v would be an infinite set of pairwise discernible elements contradicting the compactness of $\stackrel{\pm}{=}$. Thus there are α , $\beta \in v$ such that $p(\alpha) \neq p(\beta)$ and $p(\alpha) \stackrel{\pm}{=} p(\beta)$. Now, assume that $\operatorname{rng}(p)$ is compact in $\stackrel{\pm}{=}$ and $v \in \operatorname{dom}(p)$ is infinite. If $p^n v$ is finite then there are at least two elements $\alpha, \beta \in v$ such that even $p(\alpha) = p(\beta)$. If $p^n v$ is infinite then there are two $\alpha, \beta \in v$ such that $p(\alpha) \stackrel{\pm}{=} p(\beta)$.

The second \mathscr{F} -equivalence induced by the motion p in \ddagger on its domain is even of more importance:

$$\stackrel{\pm}{\mathbf{p}} \mathfrak{G} = (\forall \mathbf{J}, \delta \in [\min\{\alpha, \beta\}, \max\{\alpha, \beta\}]) \mathfrak{p}(\mathbf{J}) \stackrel{\pm}{=} \mathfrak{p}(\delta).$$

The motion p is called compact if the \mathcal{F} -equivalence \ddagger is compact. Obviously, \ddagger is finer than (\ddagger) , thus the compactness of \ddagger implies the compactness of (\ddagger) .

Let us recall from [V] that a motion p oscillates between points x and y (sets u and v) if $x \not\equiv y$ (Fig(u) \cap Fig(v) = \emptyset) and there are sequences $\{\alpha_n; n \in \mathbb{FN}\}, \{\beta_n; n \in \mathbb{FN}\}$ of elements of dom(p) such that for each n holds $\alpha_n < \beta_n < \alpha_{n+1}$ and $p(\alpha_n) \stackrel{\pm}{=} x$, $p(\beta_n) \stackrel{\pm}{=} y$ ($p(\alpha_n) \in u$, $p(\beta_n) \in v$).

We omit the proof of the following slight generalization of the result from [V].:

<u>Theorem 4</u>. The following conditions are equivalent for any - 679 - motion p in the π -equivalence \pm :

- (1) p is compact;
- (2) the trace of p is convact in ± and for each point x p⁻¹"Mon(x) is compact in ±;
- (3) p has a compact trace and there are no points x and y
 (sets u and v) such that p oscilates between x and
 y (u and v).

<u>Remark</u>. Given a \mathcal{X} -equivalence \ddagger and any set-theoretically definable function F the relation $a \ddagger_{(F)}^{\pm} b \rightleftharpoons F(a) \ddagger F(b)$ is still a \mathcal{X} -equivalence on its domain. Theorem 3 remains valid without the assumption that F is a motion, as well. Similarly, if F is a set-theoretically definable function and \leq is a set-theoretical lattice ordering of dom(F), the definitions of the \mathcal{X} -equivalence \ddagger_{F} on dom(F) and of the oscillation extend directly. A careful analysis of the proofs in [V] shows that Theorem 4 still holds for such an F.

Thus particularly Theorems 3 and 4 apply to arbitrary set functions defined on intervals of integers (V^2 -paths). Given such a function p and a \hat{r} -equivalence \pm we put for α . $\beta \in \text{dom}(p)$

$$\alpha \stackrel{*}{=} \beta = \alpha < \beta \lor \alpha \stackrel{*}{=} \beta \qquad \text{and} \\ \alpha < + \beta = \alpha < \beta & \alpha \stackrel{*}{=} \beta .$$

2. Archimedean biequivalences

For each upper bound R of the T-equivalence ± the least equivalence

 $[R] = \bigcup \{R^n; n \in FN\}$

containing R raises to a biequivalence $\langle \pm, [R] \rangle$. This con-- 680 - struction was already used in the proof of Theorem 10 in [G-Z 1]. Similarly, for each mean bound R of the bisquivalence $\langle \pm, \pm \rangle$ one obtains a bisquivalence $\langle \pm, [R] \rangle$ which is tighter than $\langle \pm, \pm \rangle$. A bisquivalence is called Archimedean if for each its mean bound R holds $[R] = (\pm)$.

<u>Theorem 5</u>. Let $\langle \pm, \leftrightarrow \rangle$ be a biequivalence. The following conditions are equivalent:

- (1) $\langle \pm, \leftrightarrow \rangle$ is Archimedean;
- (2) for each mean bound R of (±, ↔) and all x,y such that x ↔ y there is a finite R-path from x to y;

<u>Proof.</u> (1) = (2) and (1) \Rightarrow (4) \Rightarrow (5) \Rightarrow (1) are trivial. (2) \Rightarrow (3): Let $x \leftrightarrow y$, $y \in N-FN$ and $\{R_n; n \in FZ\}$ be a bix generating sequence of $\langle \pm, \leftrightarrow \rangle$. For each $n \notin 0$ there is a finite R_n -path p_n from x to y. By the prolongation axiom there is a motion p from x to y such that $p \gtrsim y$. (3) \Rightarrow (2): Let R be a mean bound of $\langle \pm, \leftrightarrow \rangle$ and $x \leftrightarrow y$. Then the set-theoretically definable class $\{y \in N; \langle x, y \rangle \in R^*\}$ contains all infinite natural numbers. Hence it contains also a finite n.

None of conditions (4) and (5) ensures that for an Archimedean biequivalence $\langle \pm, \pm \rangle$ there does not exist any biequivalence - 681 -

strictly tighter than $\langle \pm, \leftrightarrow \rangle$. The reader will easily find examples of biequivalences on RN which are strictly tighter than the compatible Archimedean biequivalence $\langle \pm, \leftrightarrow \rangle$.

Condition (3) suggests that the Archimedean property is a kind of compactness concerning connectedness of galaxies of a biequivalence.

<u>Theorem 6</u>. Let $\langle \pm, \leftrightarrow \rangle$ be a bisquivalence such that for any accessible x,y there is a motion with compact trace from x to y. Then $\langle \pm, \leftrightarrow \rangle$ is Archimedean and has connected galaxies.

<u>Proof.</u> We will prove that $\langle \pm, \leftrightarrow \rangle$ satisfies condition (2) of Theorem 5. Let R be a mean bound of $\langle \pm, \leftrightarrow \rangle$, $x \leftrightarrow y$, and p be a motion with compact trace from x to y in the time \checkmark . We put G(0) = x and $G(\alpha+1) = p(\psi)$ where $\psi = \max \{ \forall \in \Im; \langle G(\alpha), p(\psi) \rangle \in R \}$ if $G(\alpha) \neq y$, and $G^{\alpha} \{ \alpha+1 \} = \emptyset$ if $G(\beta) = y$ for some $\beta \leq \alpha$. Then G is a set-theoretically definable function and its domain is a section in the linearly ordered class $\langle N, \leq \rangle$. For any $\alpha, \beta \in \text{dom}(G)$ $\alpha < \beta$ implies $G(\alpha) \not\equiv G(\beta)$ (with perhaps one exception $\beta = \alpha + 1$ and $G(\beta) = y$). Since $\operatorname{rng}(G) \subseteq \operatorname{rng}(p)$ and the latter is compact, the former has to be a finite set. As G is one-one, it is a finite set function and a finite R-path from x to y. Essentially the same argument works to establish that for any motion p with compact trace and for all $\alpha, \beta \in \operatorname{dom}(p)$ holds $p(\alpha) \leftrightarrow p(\beta)$. Thus $\langle \pm, \leftrightarrow \rangle$ has connected galaxies.

<u>Corollary</u>. Every compatible bisquivalence with connected galaxies is Archimedean.

As a byproduct of the proof of Theorem 6 one obtains:

<u>Theorem 7</u>. Let $\langle \pm, \pm \rangle$ be a biequivalence and X be a pseudocompact connected class in \pm . Then for all x, y $\in X$ holds - 682 - х 📥 у.

Notice that both the notions of pseudocompactness and connectedness are defined purely in terms of the \mathscr{E} -equivalence \ddagger . Thus Theorem 7 applies to any \mathscr{E} -equivalence $(\ddagger) \ge (\ddagger)$.

<u>Example 1</u>. Put $A = \{ \langle x, y \rangle \in \mathbb{RN}^2; x^2y^2 = 1 \}$. Then the biequivalence $\langle \doteq^2 \uparrow A, \leftrightarrow^2 \uparrow A \rangle$ is compatible, Archimedean and its domain A is connected. However, Gal($\langle 1, 1 \rangle$) is not connected. Moreover, there is no motion with compact trace from $\langle -1, 1 \rangle$ to $\langle 1, 1 \rangle$.

<u>Example 2</u>. For every set u the biequivalence $\langle \doteq^{u}, \xleftarrow^{u} \rangle$ on RN^u is Archimedean with connected galaxies and for any pair $f \Leftrightarrow^{u} g$ there is even a compact motion from f to g (it can be defined for any $\gamma \in N-FN$ by $p(\alpha)(x) = (\alpha/\gamma)f(x) + (1-\alpha/\gamma)g(x)$ for $\alpha \in \gamma + 1$, $x \in u$). Nevertheless, for infinite $u \langle \doteq^{u}, \leftrightarrow^{u} \rangle$ is not compatible.

<u>Example 3</u>. Let \ll be an infinite natural number. Put $R_n = \{\langle x, y \rangle \in RN^2; |x - y| < \alpha^n \}$ for each $n \in FZ$. Then the biequivalence with the bigenerating sequence $\{R_n; n \in FZ\}$ has connected galaxies, connected domain and is neither compatible nor Archimedean.

<u>Example 4</u>. Let $\gamma \in N-FN$. Let us endow $RN^{\gamma+1}$ with the structure of a linear space over RN in the obvious way defining the vector addition and the multiplication by scalars componentwise. Put

 $A = \{ f \in RN^{\nu+1}; f(\nu) = 1 \& rng(f) \subseteq \{-1,1\} \},$ $B = \{ tf + (1-t)g; t \in RN \& 0 \leq t \leq 1 \& f, g \in A \\ g \{ \lambda \leq \nu ; f(\lambda) \neq g(\lambda) \} \stackrel{\sim}{\approx} 1 \},$ - 683 = -

$$T = \left\{ \langle g, f \rangle \in \mathbb{RN}^{\nu+1} \times \mathbb{RN}^{\nu+1}; g(\nu) \ge 1 \ \mathfrak{g}(\nu) \ge 1 \ \mathfrak{g}(\nu) = f(\lambda) \right\}.$$

Then the set-theoretically definable class $X = B \cup T^*A$ consists of the edges of the ν -dimensional hypercube with vertices coordinates ± 1 situated in the hyperplane $f(\nu) = 1$ in $RN^{\nu+1}$ and of parts of arcs of the hyperbolas running through the vertices of the cube to the common asymptote $f(0) = \ldots = f(\nu-1) = 0$. Then the biequivalence $\langle = \nu^{\nu+1} | X, \leftrightarrow \nu^{\nu+1} | X \rangle$ is Archimedean with connected galaxies and connected domain. However, there is no motion with compact trace from $\{1\} \times (\nu+1)$ to $(\{-1\} \times \nu) \cup \{\langle 1, \nu \rangle\}$ in X.

Example 5. In this Example $[a,b] = \{x \in \mathbb{R}N; a \leq x \leq b\}$ always denotes the interval of rationals botween $a,b \in \mathbb{R}N$. Put $I_0 = [0,1] \times \{0\}$ and $I_{\alpha} = [0,1] \times \{1/\alpha\}$ for $\alpha \in \mathbb{N} - \{0\}$, $J_{\alpha} = \begin{cases} \{0\} \times [1/(\alpha+1), 1/\alpha] & \text{for even } \alpha \in \mathbb{N} - \{0\} \end{cases}$ $\{1\} \times [1/(\alpha+1), 1/\alpha] & \text{for odd } \alpha \in \mathbb{N} - \{0\} \end{cases}$

Finally $A = I_0 \cup \bigcup \{ I_a \cup J_a ; a \in N - \{0\} \}$ is a set-theoretically definable class. Then $\doteq^2 \uparrow A$ is a compact \mathfrak{A} -equivalence (the biequivalence $\langle \doteq^2 \uparrow A, A^2 \rangle$ is compatible) with connected domain A (its single galaxy). Thus each motion in $\doteq^2 \uparrow A$ has a compact trace and, in particular, $\langle \doteq^2 \uparrow A, A^2 \rangle$ is an Archimedean biequivalence. Nevertheless, there is no compact motion from $\langle 0, 0 \rangle$ to $\langle 0, 1 \rangle$ in $\doteq^2 \uparrow A$ since every such a motion has to oscillate between the points $\langle 0, 0 \rangle$ and $\langle 1, 0 \rangle$.

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3. The metrization theorem and embeddings of biequivalences

For the rest of the article Sd_V^* denotes a fixed revealment of the codable class Sd_V of all set-theoretically definable classes (see [S-V 2]). Everything one needs to know is that Sd_V^* is a fully revealed codable class (i.e. there is a code $\langle K,S \rangle$ of Sd_V^* such that the class $K \times \{0\} \cup S \times \{1\}$ is fully revealed) satisfying the following conditions:

- (1) $\operatorname{Sd}_{v} \leq \operatorname{Sd}_{v}^{\#}$ and each class $X \in \operatorname{Sd}_{v}^{\#}$ is fully revealed;
- (2) for each set u and each $X \in Sd_V^*$ un X is a set;
- (3) for each normal formula $\varphi(x_0, X_0)$ of the language FL_V and each $X \in Sd_V^*$ holds $\{x; \varphi(x, X)\} \in Sd_V^*;$
- (4) if $X \in Sd_V^*$ and $X \cap N \neq \emptyset$ then there is the least element of $X \cap N$ in the natural ordering of N;
- (5) if $X \in Sd_V^{\#}$, $\emptyset \in X$ and $(\forall x, y)(x \in X \Rightarrow x \cup \{y\} \in X)$ then X = V (induction);
- (6) if X ∈ Sd^{*}_V and (∀x)(x ⊆ X ⇒ x ∈ X) then X = V (∈ -induction);
- (7) if $\{X_n; n \in FN\}$ is a sequence of classes from $Sd_V^{\#}$ then there is an $R \in Sd_V^{\#}$ such that $R^{\#}\{n\} = X_n$ for each $n \in FN$ (prolongation).

According to (1) - (7), Sd_V^* should be understood as a "well behaved" system of "well behaved" classes conveniently extending Sd_V admitting "well behaved" prolongations of countable sequences of its members.

On the base of Sd^{*}_y such notions as "X-class, "&-class, "X- and "&-equivalence, "generating sequence of a "X- or of a *&-equivalence, "bigenerating sequence, et cetera, can be de-

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fined in the obvious way. (E.g. X is a **-class if there is a sequence $\{X_n; n \in \mathbb{N}\}$ of classes from Sd_V^* such that $X = \bigcap \{X_n; n \in \mathbb{N}\}$; or a *bigenerating sequence is a sequence $\{R_n; n \in \mathbb{F}Z\}$ of reflexive and symetric relations from Sd_V^* such that for each n holds $R_n \circ R_n \subseteq R_{n+1}$.) The reader should think over that any result concerning the "star-free" notions from [V], [G-Z 1] or from this article remains true under an appropriate "starification". Finally, notice that the restriction of a *biequivalence to a set is always a biequivalence.

For the sake of transparency we shall deal also with "biequivalences with domains different from V. A triple $\langle X, \pm, \leftrightarrow \rangle$ where $\langle \pm, \leftrightarrow \rangle$ is a "biequivalence and $\emptyset \neq X \in \operatorname{Sd}_V^*$ is its domain will be called a "biequivalence space.

Let $\{R_n; n \in FZ\}$ be a "bigenerating sequence (of some "biequivalence $\langle \pm, \stackrel{+}{\longleftrightarrow} \rangle$). The sequence $\{R_y; y \in [e, \tau]\}$ where -e, $\tau \in N$ -FN is called a prolongation of the "bigenerating sequence $\{R_n; n \in FZ\}$ in Sd_V^* if the class $S = \bigcup \{R_y \times \{y\}; y \in [\sigma, \tau]\}$ belongs to Sd_V^* , for each $y \in [e, \tau]$ $R_y \in Sd_V^*$ is a reflexive and symetric relation and $R_y \circ R_y \subseteq R_{y+1}$, $S^*\{n\} = R_n$ holds for each $n \in FZ$, and $R_F = Id$, $R_\tau = V^2$.

The following theorem is a direct consequence of the prolongation condition (7):

<u>Theorem 8</u>. Every * bigenerating sequence has a prolongation in Sd_v^* .

Let X be a nonempty class, H be a function with domain X^2 and $rng(H) \subseteq HN$. Then H is called a metric on X if for all x,y,z $\in X$ holds $H(x,y) \ge 0$, H(x,y) = H(y,x), $H(x,y) + H(y,z) \ge H(x,z)$ and $H(x,y) = 0 \equiv x = y$. The weakening - 686 - of the last condition to mere H(x,x) = 0 leads to the notion of pseudometric. When also metrics taking values in other ordered fields than RN will be considered, we will refer to the notion just defined as to a rational metric.

If H is a metric on X then the pair $\langle X,H \rangle$ is called a (rational) metric space. Given a metric space $\langle X,H \rangle$ we put for $x,y \in X$

 $x \doteq_{H} y \equiv H(x,y) \doteq 0$ and $x \leftrightarrow_{H} y \equiv H(x,y) \leftrightarrow 0.$

The next theorem shows that the AST succeeded in a natural way completely to exclude the pathologies of nonmetrizable spaces from our study and to recure the balance between the topology and "measuring of distances" both on the discernibility and accessibility horizons. From this point of view the indiscernibility and accessibility equivalences occur as mere certain invariants of metric spaces.

<u>Theorem 9.</u> If $\langle X,H \rangle \in Sd_V^*$ is metric space (i.e. $H \in Sd_V^*$ is a metric on X) then $\langle \stackrel{\cdot}{=}_H, \stackrel{\leftarrow}{\leftrightarrow}_H \rangle$ is a "biequivalence with domain X. Conversely, for every "biequivalence space $\langle X, \stackrel{\pm}{=}, \stackrel{\leftarrow}{\leftrightarrow} \rangle$ there is a metric $H \in Sd_V^*$ on X such that $\langle \stackrel{\pm}{=}, \stackrel{\leftarrow}{\to} \rangle = \langle \stackrel{\cdot}{=}_H, \stackrel{\leftarrow}{\leftrightarrow}_H \rangle$.

<u>Proof</u>. The first assertion is trivial. The converse follows directly from the existence of a "bigenerating sequence for $\langle \pm, \leftrightarrow \rangle$ (see [G-Z 1]) and from the \mathcal{T} - and/or \mathcal{C} -valuation lemma ([M 2]). All one has to do is to take a suitable prolongation of an appropriate "bigenerating sequence of $\langle \pm, \leftrightarrow \rangle$ in Sd^{*}_V. We omit the precious proof which is, in fact, implicitly contained in [M 2]. The reader will be made amends in the next section where for a more specific class of (")biequivalences a metric subject to - 687 - some additional properties will be constructed using ideas similar to the dropped ones.

The "biequivalence $\langle \vdots_{H}, \longleftrightarrow_{H} \rangle$ will be called the "biequivalence induced by the metric $H \in Sd_{V}^{*}$ (of course, $\langle \vdots_{H}, \hookleftarrow_{H} \rangle$ can be induced by many different metrics). Obviously, every metric $H \in Sd_{V}$ induces a biequivalence on its domain, though the converse is not true: there are biequivalences which cannot be induced by any set-theoretically definable metric.

The reader can easily verify that for any set u the biequivalence $\langle \doteq^{u}, \leftrightarrow^{u} \rangle$ is induced by the set-theoretically definable metric

 $D(f,g) = \max \{ |f(t) - g(t)|; t \in u \}.$

A function E: $X \rightarrow X'$ is called an isometry of the metric space $\langle X,H \rangle$ into the metric space $\langle X,H' \rangle$ if for all $x,y \in X$ holds H(x,y) = H'(E(x),E(y)).

The following result is fairly expected in the light of the classical topology:

<u>Theorem 10</u>. Let $\langle u,h \rangle$ be a metric space (u and h are sets). Then the function e: $u \rightarrow RN^{u}$ given by e(x)(t) = h(x,t)for $x,t \in u$ is an isometry of $\langle u,h \rangle$ into $\langle RN^{u},D \rangle$, and for each $x \in u$ the function $e(x) \in RN^{u}$ satisfies

 $t \stackrel{*}{=}_{h} z \rightarrow e(x)(t) \stackrel{*}{=} e(x)(z),$

 $t \leftrightarrow_h z \rightarrow e(x)(t) \leftrightarrow e(x)(z)$

for all t,z € u.

<u>Proof</u>. The fact that e is an isometry follows from the computation

$$D(e(x), e(y)) = \max \{ |h(x, t) - h(y, t)|; t \in u \}$$

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$$4 h(x,y) = |e(x)(y) - e(y)(y)|$$

 $4 D(e(x),e(y))$

The rest of the Theorem follows from the inequality $|e(x)(t) - e(x)(z)| \leq e(t,z).$

Thus in particular according to the results from [G-Z 1] each function e(x) is uniformly continuous from \doteq_h to \doteq . Let $\langle X, \pm, \leftrightarrow \rangle$, $\langle X', \pm, \leftrightarrow \rangle$ be two "biequivalence spaces. A one-one map E: $X \rightarrow X'$ is called an embedding of $\langle X, \pm, \leftrightarrow \rangle$ into $\langle X', \pm, \leftrightarrow \rangle$ iff for all $x, y \in X$ holds $x \pm y = E(x) \pm E(y)$ and $x \leftrightarrow y = E(x) \leftrightarrow E(y)$.

An embedding of a "*n*-equivalence space $\langle x, \pm \rangle$ into another

* π -equivalence space $\langle X, \dot{=} \rangle$ can be treated as an embedding of the "biequivalence space $\langle X, \dot{=}, X^2 \rangle$ into the "biequivalence space ce $\langle X', \dot{=}, X'^2 \rangle$. Obviously, every isometry of the metric space $\langle X, H \rangle \in Sd_V^*$ into the metric space $\langle X', H' \rangle \in Sd_V^*$ is an embedding of the "biequivalence space $\langle X, \dot{=}_H, \dot{=}_H \rangle$ into the "biequivalence space $\langle X', \dot{=}_H', \dot{=}_H' \rangle$.

Then Theorems 9 and 10 have the following consequence:

<u>Theorem 11</u>. For every biequivalence space $\langle u, \pm, \leftrightarrow \rangle$ there is an embedding e of $\langle u, \pm, \leftrightarrow \rangle$ into $\langle RN^{u}, \pm^{u}, \leftrightarrow^{u} \rangle$ such that all the functions $e(x) \in RN^{u}$ $(x \in u)$ are uniformly continuous from \pm to \pm .

Let $\stackrel{x}{=}$ be the **x**-equivalence on RN×u given by $\langle a,x \rangle \stackrel{x}{=} \langle b,y \rangle \cong a \stackrel{z}{=} b & x \stackrel{z}{=} y$. We also put for $f \in RN^{U}$

 $\|\mathbf{f}\| = \sum_{\mathbf{t} \in \mathbf{u}} |\mathbf{f}(\mathbf{t})| / Card(\mathbf{u}).$

The function e from Theorem 10 can be also regarded as an - 689 - embedding of the \mathbf{x} -equivalence space $\langle u, \pm \rangle$ into various \mathbf{x} -equivalence spaces with domain RN^U using the results from [G-Z 1].

<u>Theorem 12</u>. Let $\langle u, \pm \rangle$ be a \mathcal{F} -equivalence space. Then there is a one-one set-function e: $u \rightarrow RN^{u}$ such that for all $x, y \in u$ the following conditions are equivalent:

(1) $x \stackrel{\pm}{=} y$; (2) $e(x) \stackrel{\pm}{=} u e(y)$; (3) $\operatorname{Fig}^{x}(e(x)) = \operatorname{Fig}^{x}(e(y))$. If $u \in \mathbb{N}$ and $\alpha \stackrel{\pm}{=} \beta \equiv \alpha/u \stackrel{\pm}{=} \beta/u$ holds for all $\alpha, \beta \in u$ then the conditions (1) - (3) are equivalent to

(4) $\|e(x) - e(y)\| = 0.$

If one would like to generalize the above embedding results to arbitrary "biequivalences (however, to deal with biequivalences with domain V is quite sufficient), he will find unavoidable to extend the ordered field of all rational numbers in such a way that every nonempty subclass of RN belonging to Sd_V^* which has an upper bound in RN had the supremum in the extension.

A nonempty proper subclass C of RN is called a cut of $\langle RN, \leq \rangle$ if it is a section of $\langle RN, \leq \rangle$ without the greatest element.

Then the fully revealed codable class HR of all cuts in RN belonging to Sd_V^* can be given the structure of an ordered field in the obvious way. It will be called the field of all hyperreal numbers. Using an appropriate coding of HR, of the equality relation on HR and of the operations and order relation on HR, one can work with it as if it were a class from the extended universe. RN can be naturally embedded as an ordered subfield into HR. Likewise ER can be endowed with a pair of relations $\langle \pm, \leftrightarrow \rangle$ behaving as a biequivalence with domain HR prolonging the biequivalence

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from RN denoted by the same symbol in such a way that HR = Fig(RN). A hyperreal number will be called set-theoretically definable if it is determined by a set-theoretically definable cut. The set-theoretically definable hyperreals contain all rationals and form an ordered subfield of the hyperreals. Each nonempty subclass of RN with an upper bound belonging to Sd_V has the supremum in that field. The reason why the field of all set-theoretically definable hyperreal numbers is an unsatisfactory extension of RN is that one cannot apply the prolongation technics in it. HR can be also obtained as a revealment of the field of all set-theoretically definable cuts in RN.

Each nonempty class $X \in Sd_V^*$, $X \subseteq RN$, with an upper (lower) bound in RN has the supremum sup X (infimum inf X) in HR. The ordered field HR is determined by its properties with respect to Sd_V^* uniquely up to an isomorphism. Also the hyperreal numbers constructed on the base of another revealment of Sd_V , say Sd_V^* , are isomorphic to "our" HR via the automorphism of the universe mapping Sd_V^* onto Sd_V^* (see [S-V 2]).

Let us denote just for a moment

 $RN^{X} = \{F \in Sd_{V}^{\#}; dom(F) = X \ \# rng(F) \subseteq RN \ \# (\exists a \in RN) (\forall x \in X) \ |F(x)| < a \}$

the codable class of all bounded rational functions with domain X belonging to Sd_V^* (clearly $\operatorname{RN}^X \neq \emptyset$ iff $X \in \operatorname{Sd}_V^*$). Then for each $X \in \operatorname{Sd}_V^*$, RN^X can be converted into a metric space endowed with a hyperreal metric

 $D(F,G) = \sup \{ | F(x) - G(x) | ; x \in X \}$ Notice that for X being a set RN^X and D coincide with the original ones.

The generalization of Theorems 10 - 12 to arbitrary "biequiva-- 691 - lences using the hyperreal numbers is quite straightforward, now. It is left to the reader.

From the matter just indicated it should follow that the hyperreal numbers will play rather an auxiliary role of a technically convenient extension of the rationals in our study. From this point of view the irrational numbers in HR, and the more, the not set--theoretically definable ones, seem much more curious and odder than the infinitesimally small and infinitely large rationals.

4. Geodetical biequivalences

Let $\langle \pm, \pm \rangle$ be a biequivalence (with domain V). We already know that there is a metric $H \in Sd_v^{\#}$ on V inducing $\langle \pm, + \rangle$. Using the metric H a ternary relation "t lies between x and y" can be defined by the equality H(x,t) + H(t,y) = H(x,y). Similarly. one can define the ternary relation "t lies nearly between x and y" by H(x,t) + H(t,y) = H(x,y). According to some results in [G] concerning classical metric spaces one can show that the biequivalence $\langle \pm, \leftrightarrow \rangle$ can be induced by a metric $H \in Sd_u^*$ such that for all x,y,t holds t lies between x and y iff t = xor t = y, and t lies nearly between x and y iff $t \doteq x$ or t ± y. The reader will probably agree that such a metric is rather a "bad" one. According to a "good" metric H at least for any accessible pair x,y there should be a connected set u containing both x and y such that each t \in u lies between (or at least nearly between) x and y. This section is devoted to the precisation of the notion of a "good" metric and to the characterisation of biequivalences which can be induced by such metrics.

Let $H \in Sd_V^*$ be a rational metric on V and p be a path [1.e. • V²-path) with domain $[4, \sqrt[3]]$. The rational number - 692 - $L_{H}(p) = \sum_{\alpha=1}^{\beta-1} H(p(\alpha), p(\alpha+1))$

is called the length of the path p with respect to the metric H. Then p is called a direct (nearly direct) path with respect to H if $L_{H}(p) = H(p(q), p(q))$ ($L_{H}(p) \doteq H(p(q), p(q))$). When the metric H is clear from the context the attribute "with respect to H" can be omitted from the notions just introduced. The length of the path p will be denoted L(p) in such a case.

In the following three theorems $H \in \operatorname{Sd}_V^*$ denotes a fixed metric on V and $\langle \pm, \pm \rangle$ is the *biequivalence induced by it.

<u>Theorem 13</u>. Let p be a path in the time $rac{1}{2}$.

- (1) p is a direct path iff for all $\alpha \leq \beta \leq \delta'$ holds H(p(α),p(β)) = L(p^{(α}, β]);
- (2) p is a nearly direct path iff for all $\alpha \leq \beta \leq \vartheta$ holds H(p(α),p(β)) \doteq L(p [α , β]).

<u>Proof</u>. We will prove only the second claim which is a bit less trivial. Let p be nearly direct. Then $H(p(\alpha),p(\beta)) \leq$ $L(p \upharpoonright [\alpha,\beta])$ for all $\alpha \leq \beta \leq \sqrt{2}$. Assume that $H(p(\alpha),p(\beta)) < \cdot$ $L(p \upharpoonright [\alpha,\beta])$ for some α,β . Then

$$\begin{split} L(p) &= L(p \upharpoonright [0, \alpha]) + L(p \upharpoonright [\alpha, \beta]) + L(p \upharpoonright [\beta, \beta]) \\ & \rightarrow H(p(0), p(\alpha)) + H(p(\alpha), p(\beta)) + H(p(\beta), p(\beta)) \\ & \geqslant H(p(0), p(\beta)). \end{split}$$

Thus p were not nearly direct. The remaining implication is trivial.

<u>Corollary</u>. If p is a (nearly) direct path from x to y then each $t \in rng(p)$ lies finearly) between x and y.

Theorem 14. Let p be a path in the time
$$\sqrt[7]{}$$
 and $\ll, \beta \in \sqrt[7]{+1}$.
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- (1) If p is a direct path then $\alpha \leq \beta \equiv H(p(0),p(\alpha)) \leq H(p(0),p(\beta))$ $\alpha = \beta \equiv H(p(0),p(\alpha)) = H(p(0),p(\beta)).$
- (2) If p is a nearly direct path then $\alpha \stackrel{2}{\underset{p}{\leftarrow}} \beta \equiv H(p(0), p(\alpha)) \stackrel{2}{\leftarrow} H(p(0), p(\beta))$ $\alpha \stackrel{+}{\underset{p}{\leftarrow}} \beta \equiv H(p(0), p(\alpha)) \stackrel{2}{=} H(p(0), p(\beta)).$

<u>Proof</u>. We will prove only (2) again. Let p be nearly direct and $d \stackrel{>}{\geq} \beta$. Then either $d \stackrel{<}{\leq} \beta$ and by the preceeding Theorem

$$H(p(0),p(\alpha)) \leq H(p(0),p(\alpha)) + H(p(\alpha),p(\beta))$$

$$\stackrel{*}{=} L(p \upharpoonright [0, \alpha]) + L(p \upharpoonright [\alpha, \beta])$$

$$= L(p \upharpoonright [0, \beta]) \stackrel{*}{=} H(p(0),p(\beta)),$$

or $\alpha > \beta$, $p(\alpha) \stackrel{\pm}{=} p(\beta)$ and

$$H(p(0),p(\alpha)) \doteq L(p \upharpoonright [0, \alpha]) = L(p \upharpoonright [0, \beta]) + L(p \upharpoonright [\beta, \alpha])$$
$$\doteq H(p(0),p(\beta)) + H(p(\beta),p(\alpha))$$
$$\doteq H(p(0),p(\beta)).$$

Now assume that $H(p(0),p(\alpha)) \leq H(p(0),p(\beta))$. If $\alpha \leq \beta$, there is nothing to be proved. So let $\alpha > \beta$. Then as already proved also $H(p(0),p(\beta)) \leq H(p(0),p(\alpha))$. If there were a $\gamma \in [\beta, \alpha]$ such that $p(\beta) \neq p(\gamma)$, the following computation would yield a contradiction:

$$H(p(0),p(\beta)) \stackrel{*}{=} H(p(0),p(\alpha)) \stackrel{*}{=} L(p \upharpoonright [0, \alpha])$$

= $L(p \upharpoonright [0, \beta]) + L(p \upharpoonright [\beta, \gamma]) + L(p \upharpoonright [\gamma, \alpha])$
> $H(p(0),p(\beta)) + H(p(\beta),p(\gamma)) + H(p(\gamma),p(\alpha))$
 $\rightarrow H(p(0),p(\beta)).$

The second equivalence in (2) is a direct consequence of the first one.

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Corollary. Let p be a nearly direct path. Then

- (1) the equivalences $\stackrel{\pm}{=}$ and $\stackrel{\mp}{=}$ on dom(p) coincide;
- (2) p is a compact path (i.e. \ddagger is compact) iff p has a compact trace iff $L(p) \leftrightarrow 0$.

<u>Theorem 15</u>. Let p be a nearly direct path. Then rng(p) is a connected set iff p is a motion.

<u>Proof.</u> Obviously, the trace of a motion is a connected set. Assume that **p** is a nearly direct path in the time \mathcal{F} which is not a motion. Then there is an $\alpha < \mathcal{F}$ such that $p(\alpha) \not\equiv p(\alpha+1)$. Then for all $\beta, \gamma \leq \mathcal{F}$ $\beta \leq \alpha$ and $\alpha < \mathcal{F}$ imply $\beta \not\equiv \gamma$. By the previous results $p(\beta) \not\equiv p(\gamma)$. Thus rng(p) is not connected.

Theorems 13 - 15 and their Corollaries justify the following definition:

A metric $H \in Sd_V^*$ is called (nearly) geodetical if for all x,y such that $x \leftrightarrow_H y$ there is a (nearly) direct motion (with respect to H) from x to y. A biequivalence $\langle \pm, \leftrightarrow \rangle$ is called (nearly) geodetical if it can be induced by a (nearly) geodetical metric.

An immediate consequence of this definition and of the preceeding results is the following:

<u>Theorem 16</u>. Let $\langle \pm, \pm \rangle$ be a nearly geodetical biequivalence. Then for every pair $x \leftrightarrow y$ there is a compact motion from x to y. In particular $\langle \pm, \pm \rangle$ is Archimedean and has connected galaxies.

<u>Theorem 17</u>. Let $\langle \pm, \pm \rangle$ be a biequivalence. The following conditions are equivalent:

- (1) $\langle \pm, \leftrightarrow \rangle$ is geodetical;
- (2) $\langle \pm, \pm \rangle$ is nearly geodetical;

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(3) there is a *bigenerating sequence $\{R_n; n \in FZ\}$ of $\langle \pm, \leftrightarrow \rangle$ such that for each n hold $R_{n+1} \subseteq R_n \circ R_n \circ (\pm)$ and $(\pm) \circ R_n = R_n \circ (\pm);$

(4) there is a "bigenerating sequence
$$\{S_n; n \in FZ\}$$

of $\langle \pm, \leftrightarrow \rangle$ such that for each n holds $S_n \circ S_n = S_{n+1}$;

(5) for some $d \in FN$, $d \ge 2$, there is a *bigenerating sequence ce $\{S_n; n \in FZ\}$ such that for each n holds $S_n^d = S_{n+1}$.

<u>Proof.</u> (1) \Rightarrow (2) and (4) \Rightarrow (5) are trivial.

(2) \Rightarrow (3): Let $H \in Sd_V^*$ be a nearly geodetical metric inducing $\langle \pm, \pm \rangle$. We put $R_n = \{\langle x, y \rangle ; H(x, y) \leq 2^n \}$. Obviously, $\{R_n; n \in FZ\}$ is a "bigenerating sequence of $\langle \pm, \pm \rangle$. Let $\langle x, y \rangle \in R_{n+1}$ and p be a nearly direct motion from x to y in the time \mathfrak{F} . Let $d \leq \mathfrak{F}$ be the greatest natural number such that $\langle x, p(\alpha) \rangle \in R_n$, and $\beta \in [\alpha, \mathfrak{F}]$ be the greatest natural number such that $\langle p(\alpha), p(\beta) \rangle \in R_n$. It is routine to check that $p(\beta) \pm y$ since p is a nearly direct motion. Thus $R_{n+1} \subseteq R_n \circ R_n \circ (\pm)$. Now, assume that $\langle x, y \rangle \in (\pm) \circ R_n$. Let p be a nearly direct motion from x to y in the time \mathfrak{F} and $d \leq \mathfrak{F}$ be the greatest natural $\langle x, y \rangle \in (\pm) \circ R_n$. Let p be a nearly direct motion from x to y in the time \mathfrak{F} and $d \leq \mathfrak{F}$ be the greatest natural number such that $\langle x, p(\alpha) \rangle \in R_n$. Let p is a nearly direct motion from x to y in the time \mathfrak{F} and $d \leq \mathfrak{F}$ be the greatest natural number such that $\langle x, p(\alpha) \rangle \in R_n$. The reader can easily verify that $p(\alpha) \pm y$. Thus $(\pm) \circ R_n \in R_n \circ (\pm)$. The remaining inclusion follows by a symetric argument. (3) \Rightarrow (4): One can easily verify that

$$(\forall_n \in \mathbb{PZ})(\forall_k \in \mathbb{FN}) \quad \mathbb{R}_n^{2^k} \subseteq \mathbb{R}_{n+k} \subseteq \mathbb{R}_n^{2^k} \circ (\stackrel{+}{=}).$$

Therefore

$$(\forall m, n \in \mathbb{F}\mathbb{Z})(m \leq n \Rightarrow \mathbb{R}_m^{2^{m-m}} \subseteq \mathbb{R}_n \subseteq \mathbb{R}_m^{2^{m-m}+1}).$$

Hence there is a prolongation $\{R_{\gamma}; \gamma \in [\sigma-1,\tau+1]\}$ of the *bigenerating sequence $\{R_n; n \in FZ\}$ in Sd_{ψ}^{ψ} such that for each $\gamma \in [\sigma, t]$ holds - 696 - $R_{\varphi}^{2^{y-\varphi}} \subseteq R_{y} \subseteq R_{\varphi}^{2^{y-\varphi}+1} .$ Finally, we put $S_{n} = R_{\varphi}^{2^{n-\varphi}}$ for each $n \in FZ$. Then $\{S_{n}; n \in FZ\}$ is a *bigenerating sequence of $\langle \pm, \pm \rangle$ and for each n holds $S_{n} \circ S_{n} = S_{n+1} .$ $(5) \Rightarrow (1)$: Let $\{S_{n}; n \in FZ\}$ be a *bigenerating sequence of $\langle \pm, \pm \rangle$ such that $(\forall n) S_{n}^{d} = S_{n+1}$ where $d \in FN - \{0, 1\}$, and $\{s_{y}; \forall \in [e-1, t+1]\}$ be a prolongation of this sequence in Sd_{φ}^{*} such that for each $\forall \in [e, t-1]$ holds $S_{\varphi}^{d} = S_{y+1}$. Then for each $\forall \in [e, t]$ holds $S_{\varphi}^{d^{y-\varphi}} = S_{y}$. Then the function $H(x,y) = d^{\varphi} \min(\{\mu; \langle x, y \rangle \in S_{\varphi}^{*}\} \cup \{d^{Y-\varphi+1}\})$

obviously belongs to Sd_V^* and is a metric on V. Let us show that the "biequivalence induced by H is indeed $\langle \pm, \leftrightarrow \rangle$. For all x,y the following conditions are equivalent: $x \pm y;$ ($\forall n$) $\langle x, y \rangle \in S_n = S_e^{d^{n-e^*}};$ ($\forall n$) $H(x,y) \leq d^n;$ $H(x,y) \doteq 0.$ Similarly, changing " $\forall n$ " to " $\exists n$ ", one obtains ($\forall x, y$)($x \leftrightarrow y = H(x, y) \leftrightarrow 0$).

It remains to prove that H is geodetical. But from the construction of H it follows even more. Namely, for every pair $\langle x,y \rangle \in S_{+}$ there is a direct S_{-} -path from x to y.

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(Oblatum 12.2. 1985)