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# DIGRAPHS MAXIMAL WITH RESPECT TO CONNECTIVITY 

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The digraphs considered in this paper are finite without loops and multiple arcs.
The strong (unilateral, weak) connectivity $\chi^{3}=\chi^{3}(G)\left(\varkappa^{2}=\chi^{2}(G), \varkappa^{1}=\chi^{1}(G)\right)$ of a digraph $G$ is the minimum number of points whose removal results in a not strong (unilateral, weak) or trivial digraph.

We shall describe constructively and determine the number of digraphs maximal with respect to the strong or unilateral or weak connectivity, respectively.

Maximal graphs with a given vertex (edge) connectivity have been studied in [1].
Let digraphs $G_{1}$ and $G_{2}$ have disjoint sets $V_{1}$ and $V_{2}$ of points and disjoint arc sets $E_{1}$ and $E_{2}$, respectively.

Their union is the digraph $G=G_{1} \cup G_{2}$, which has the point set $V=V_{1} \cup V_{2}$ and the arc set $E=E_{1} \cup E_{2}$.

Their join $G_{1}+G_{2}$ consists of $G_{1} \cup G_{2}$ and all arcs joining $V_{1}$ with $V_{2}$.
Their directional join $G_{1} \oplus G_{2}$ consists of $G_{1} \cup G_{2}$ and all arcs going from $V_{1}$ to $V_{2}$.

It is clear that the directional join is not a commutative operation.
Definition 1. Let $G$ be a not complete digraph and $n$ a nonnegative integer. Then $G$ is called $\varkappa_{n}^{i}$-maximal if $\varkappa^{i}(G)=n$ and $\varkappa^{i}(G+x)>\varkappa^{i}(G)$ holds for every $\operatorname{arc} x \in E(\bar{G})($ for $i=1,2,3)$.

All notions not defined here will be used in the sense of [2]. The symbol $K_{n}$ denotes here the complete diagraph on $n$ points.

Theorem 1. Let $G$ be a diagraph and $n$ be a natural number. Then $G$ is $x_{n}^{i}$-maximal if and only if $G \simeq K_{n}+D$, where $D$ is a $\varkappa_{0}^{i}$-maximal digraph (for $i=1$, 2, 3).

Proof. Let $G \simeq K_{n}+D$ and be $\chi_{0}^{i}$-maximal. Let us denote $V(G)=A \cup B$, where $A=V\left(k_{n}\right), B=V(D)$.

Since $\varkappa^{i}(G-A)=\varkappa^{i}(D)=0$, we have $\varkappa^{i}(G) \leqslant n$. If $C \subset V(G),|C| \leqslant n, C \neq A$, then by $G-C=K_{m}+D^{\prime}$ it follows that the digraph $G-C$ is strong. Thus $\varkappa^{i}(G)=n$.

Let $x \in E(\bar{G})$. Then $x^{i}((G+x)-A)=x^{i}(D+x)>0$, because the digraph $D$ is $\chi_{0}^{i}$-maximal. Now $\varkappa^{i}((G+x)-C) \geqslant \chi^{i}(G-C)>0$ implies $\varkappa^{i}(G+x)>n$, i.e. the digraph $G$ is $\chi_{n}^{i}$-maximal.

Let $G$ be $\varkappa_{n}^{i}$-maximal. Then there exists a set $A$ of points of $G$ such that $|A|=n$ and $x^{i}(G-A)=0$. We denote by $D$ the digraph $G-A$. From the $x_{n}^{i}$-maximality of $G$ it follows that $G \simeq K_{n}+D$. To finish our proof we must show that the digraph $D$ is $\chi_{0}^{i}$-maximal. We shall prove it indirectly.

Let $x \in E(\bar{G})$ and $\chi^{i}(D+x)=0$. Then $\chi^{i}(G+x)=n$ and this is a conradiction because $G$ is $x_{n}^{i}$-maximal. Q.E.D.

Theorem 2. Let $G$ be a digraph. Then $G$ is
a) $x_{0}^{1}$-maximal if and only if $G \simeq K_{a} \cup K_{b}$,
b) $\chi_{0}^{2}$-maximal if and only if either $G \simeq K_{a} \cup K_{b}$, or $G \simeq K_{c} \oplus\left(K_{a} \cup K_{b}\right)$ or $G \simeq\left(K_{a} \cup K_{b}\right) \oplus K_{c}$ or $G \simeq K_{d} \oplus\left(\left(K_{a} \cup K_{b}\right) \oplus K_{c}\right)$,
c) $x_{0}^{3}$-maximal if and only if $G \simeq K_{a} \oplus K_{b}$.

Proof. One can easily verify that the sufficient condition in all three assertions holds.

Let $G$ be $x_{0}^{i}$-maximal. If $S$ is a strong component of the diagraph $G$, then the $\chi_{0}^{i}$-maximality of $G$ implies $S \simeq K_{a}$.

Let $S_{1}, S_{2}, \ldots, S_{n}$ be the strong components of $G$.
a) If $i=1$, then the digraph $G$ is disconnected and from the $x_{1}^{1}$-maximality of $G$ it follows that $G$ consists of exactly two strong components. Thus $G \simeq K_{a} \cup K_{b}$.
b) If $i=2$, then the digraph $G$ is not unilateral, hence there exist two points $u \in S_{r}, v \in S_{t}$ such that $u$ cannot be reached from $v$ and $v$ cannot be reached from $u$.

Let us put

$$
\begin{gathered}
\mathcal{M}=\left\{S_{i}, j \neq r, j \neq t, S_{i} \text { can be reached from } S_{t} \text { or from } S_{r}\right\} \\
\mathcal{N}=\left\{S_{i}, j \neq r, j \neq t, S_{i} \notin \mathscr{M}\right\} .
\end{gathered}
$$

The $\chi_{0}^{2}$-maximality of $G$ implies that $\mathcal{M}$ (and analogously $\mathcal{N}$ ) is empty or it contains exactly one strong component.

We have to consider only four cases.

1. Let $\mathcal{M}=\emptyset, \mathcal{N}=\emptyset$. Then $G \simeq K_{a} \cup K_{b}$.
2. Let $\mathcal{M}=\emptyset, \mathcal{N} \neq \emptyset$. Then $G \approx K_{c} \oplus\left(K_{a} \cup K_{b}\right)$.
3. Let $\mathcal{M} \neq \emptyset, \mathcal{N}=\emptyset$. Then $G \simeq\left(K_{a} \cup K_{b}\right) \oplus K_{c}$.
4. Let $\mathcal{M} \neq \emptyset, \mathcal{N} \neq \emptyset$. Then $G \simeq K_{d} \oplus\left(\left(K_{a} \cup K_{b}\right) \oplus K_{c}\right)$.

We can prove $1-4$ by the $\chi_{0}^{2}$-maximality of $G$.
c) If $i=3$, then the diagraph $G$ is unilateral. The strong components of a unilateral diagraph $G$ can be denoted in such a way that $S_{c}$ can be reached from $S_{d}$ if and only if $c \geqslant d$ (see [2], p. 200). Let $n \geqq 3$, then the digraph $G+x$ (where
$\left.x \in E(\bar{G}), x=u v, v \notin S_{1}\right)$ is not strong. This is a contradiction, because $G$ is $\chi_{0}^{3}$ - maximal. Hence $n=2$ and from the $\chi_{0}^{3}$-maximality of $G$ it follows that $G \simeq K_{a} \oplus K_{b}$. Q.E.D.

By using Theorem 1 and 2 we prove the following inequalities.
Theorem 3. Let $G$ be a digraph with $p$ points and $q$ arcs. Let $\chi^{i}(G)=n$. Then we have:
a) $q(G) \leqslant(p-1)(p-2)+2 n \quad($ for $i=1,2)$,
b) $q(G) \leqslant(p-1)^{2}+n \quad($ for $i=3)$.

Proof. Let $\chi^{i}(G)=n$. Then there exists an $\varkappa_{n}^{i}$-maximal diagraph $H$ such that $G$ is a factor of $H$.

By Theorem 1 we have $H \simeq K_{n}+D$, where $D$ is $\chi_{0}^{i}$-maximal. Then

$$
\begin{equation*}
q(G) \leqslant q(H)=n(n-1)+2 n(p-n)+q(D) . \tag{1}
\end{equation*}
$$

a) Let $D$ be $\chi_{0}^{1}$-maximal. According to Theorem $2 D \simeq K_{a} \cup K_{b}$. Then $q(D)$ $=a(a-1)+b(b-1)$, where $a+b=p-n$.
The maximum of the function $q(D)$ (for $1 \leqslant a \leqslant p-n-1$ ) is reached in the digraph $K_{1} \cup K_{p-n-1}$. We have

$$
\begin{equation*}
q(D) \leqslant(p-n-1)(p-n-2) . \tag{2}
\end{equation*}
$$

If $\chi^{1}(G)=n$, then from (1) and (2) it follows that

$$
q(G) \leqslant(p-1)(p-2)+2 n
$$

b) Let $D$ be $\chi_{0}^{2}$-maximal. According to Theorem 2 either $D \simeq K_{a} \cup K_{b}$ or $D \simeq\left(K_{a} \cup K_{b}\right) \oplus K_{c}$ or $D \simeq K_{a} \oplus\left(K_{b} \cup K_{c}\right)$ or $D \simeq K_{a} \oplus\left(\left(K_{b} \cup K_{c}\right) \oplus K_{d}\right)$.

Without loss of generality we can suppose that $b \leqslant c$. One can easily verify that

$$
\begin{gather*}
q\left(K_{a} \oplus\left(\left(K_{b} \cup K_{c}\right) \oplus K_{d}\right)\right) \leqslant q\left(\left(K_{b} \cup K_{c}\right) \oplus K_{a+d}\right)=  \tag{3}\\
\quad=q\left(K_{a+d} \oplus\left(K_{b} \cup K_{c}\right)\right) \leqslant q\left(K_{b} \cup K_{a+c+d}\right) .
\end{gather*}
$$

If $\varkappa^{2}(G)=n$, from (1), (2), (3) it follows that

$$
q(G) \leqslant(p-1)(p-2)+2 n .
$$

c) Let $D$ be $x_{0}^{3}$-maximal. According to Theorem $2 D \simeq K_{a} \oplus K_{b}$. Then $q(D)$ $=a(a-1)+b(b-1)+a b$, where $a+b=p-n$. The maximum of the function $q(D)$ (for $1 \leqslant a \leqslant p-n-1$ ) is reached in the digraph $K_{1} \oplus K_{p-n-1}$. We have

$$
\begin{equation*}
q(D) \leqslant(p-n-1)^{2} \tag{4}
\end{equation*}
$$

If $\chi^{3}(G)=n$, then (1) and (4) imply

$$
q(G) \leqslant(p-1)^{2}+n
$$

By using Theorem 1 and 2 we determine the number of maximal digraphs with respect to the connectivity.

Theorem 4. The number of nonisomorphic $\chi_{n}^{i}$-maximal digraphs with p points is
a) $\left[\frac{m}{2}\right]$
if $i=1, m \geqslant 2$,
b) $\frac{1}{24}\left(2 m^{3}+3 m^{2}-5 m+6\left[\frac{m}{2}\right]\right)$,
if $i=2, m \geqslant 2$,
c) $m-1$
if $i=3, m \geqslant 2$,
where $m=p-n$.
Proof. One can easily verify that parts a) and c) hold. Let $|V(G)|=p$. There are
$\left[\frac{m}{2}\right]$ nonisomorphic ${\chi_{n}^{2}}^{2}$-maximal digraphs of the form $K_{a} \cup K_{b}$,
$\sum_{i=2}^{m-1}\left[\frac{i}{2}\right]=\frac{1}{4}\left(m^{2}-m-2\left[\frac{m}{2}\right]\right)$ nonisomorphic $\mathcal{K}_{n}^{2}$-maximal digraphs of the form $K_{a} \oplus\left(K_{b} \cup K_{c}\right)$ and of the form $\left(K_{a} \cup K_{b}\right) \oplus K_{c}$, too, and $\sum_{i=3}^{m-1} \frac{1}{4}\left(j^{2}-j-2\left[\frac{j}{2}\right]\right)$ $=\frac{1}{24}\left(2 m^{3}-9 m^{2}+7 m+6\left[\frac{m}{2}\right]\right)$ nonisomorphic $x_{n}^{2}$-maximal digraphs of the form $K_{a} \oplus\left(\left(K_{b} \cup K_{c}\right) \oplus K_{d}\right)$, where $m=p-n$. Thus there are $\frac{1}{24}\left(2 m^{3}+3 m^{2}\right.$ $-5 m+6\left[\frac{m}{2}\right]$ ) nonisomorphic $\chi_{n}^{2}$-maximal digraphs with $p=m+n$ points. Q.E.D.

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## ОРГРАФЫ МАКСИМАЛЬНЫЕ ОТНОСИТЕЛЬНО СВЯЗНОСТИ

Петер Горак
Резюме
Сильной (односторонней, слабой) связностью $\chi^{3}=\chi^{3}(K) \chi^{2}=\chi^{2}(K), \chi^{1}=\chi^{1}(K)$ орграфа $K$ называется наименьшее число вершин, удаление которых приводит к не сильному (одностороннему, слабому), или же тривиальному орграфу.

Конструктивно описано и определено число орграфов максималшных относительно сильнои или односторонней или слабой связности.

