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A REMARK ON SYSTEMS OF MAXIMAL CLIQUES OF A GRAPH

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In [1] a concept of a  $\tau$ -covering of a set was defined and applied to the study of tolerances, i.e. reflexive and symmetric binary relations. The aim of this paper is to translate the results from [1] into the language of the graph theory.

Let  $S$  be a non-empty set. A covering of  $M$  by subsets is a system  $\mathfrak{M}$  of subsets of  $S$  such that  $S \subseteq \bigcup_{M \in \mathfrak{M}} M$ . A covering  $\mathfrak{M}$  of the set  $S$  by subsets is called a  $\tau$ -covering of  $S$ , if and only if

(i) for each  $M_0 \in \mathfrak{M}$  and each  $\mathfrak{N} \subseteq \mathfrak{M}$  the following implication holds:

$$M_0 \subseteq \bigcup_{M \in \mathfrak{N}} M \Rightarrow \bigcap_{M \in \mathfrak{N}} M \subseteq M_0 ;$$

(ii) if  $R \subseteq S$  and  $R$  is contained in no set from the covering, then  $R$  contains a two-element subset of the same property.

**Theorem 1.** *Let  $G$  be an undirected graph with the vertex set  $V$ , let  $\mathfrak{M}$  be the system of vertex sets of all maximal cliques of  $G$ . Then  $\mathfrak{M}$  is a  $\tau$ -covering of  $V$ . Conversely, if a set  $V$  and its  $\tau$ -covering  $\mathfrak{M}$  are given, then there exists a graph  $G$  such that  $V$  is its vertex set and  $\mathfrak{M}$  is the system of vertex sets of all maximal cliques of  $G$ .*

*Proof.* Let  $\mathfrak{M}$  be the system of vertex sets of all maximal cliques of  $G$ . Let  $M_0 \in \mathfrak{M}$ ,  $\mathfrak{N} \subseteq \mathfrak{M}$  and let  $M_0 \subseteq \bigcup_{M \in \mathfrak{N}} M$ . The set  $M_0$  is the vertex set of a maximal clique  $C_0$  of  $G$ . Suppose that  $\bigcap_{M \in \mathfrak{N}} M \not\subseteq M_0$ . Denote  $N_0 = \bigcap_{M \in \mathfrak{N}} M$ . Then  $N_0 - M_0 \neq \emptyset$ . Let  $u \in N_0 - M_0$ . If  $v \in M_0$ , then there exists a set  $N(v) \in \mathfrak{N}$  such that  $v \in N(v)$ . As  $N(v) \in \mathfrak{N}$ ,  $u \in N_0$ , we have  $u \in N(v)$ . The set  $N(v)$  is the vertex set of a maximal clique of  $G$ , therefore  $u$  and  $v$  are joined by an edge. As  $v$  was chosen arbitrarily, the vertex  $u$  is joined with all vertices of  $M_0$ . Thus  $M_0 \cup \{u\}$  is the vertex set of a clique  $C_1$  containing  $C_0$  as a proper subgraph, which is a contradiction with the assumption that  $C_0$  is a maximal clique. Therefore (i) is fulfilled. Now let  $R \subseteq V$  and let  $R$  be contained

in no set from  $\mathfrak{M}$ . This means that  $R$  is contained in the vertex set of no maximal clique. The subgraph  $G(R)$  of  $G$  induced by  $R$  is not a clique, because every clique is contained in a maximal clique. Thus  $R$  contains two vertices which are not adjacent in  $G$ ; this means that the set consisting of these two vertices is a subset of the vertex set of no clique and (ii) is fulfilled. The covering  $\mathfrak{M}$  is a  $\tau$ -covering of  $V$ .

Now let a set  $V$  and its  $\tau$ -covering  $\mathfrak{M}$  be given. We join by an edge any two distinct elements of  $V$  to which a set from  $\mathfrak{M}$  exists containing both of them; thus we obtain a graph  $G$ . Each set from  $\mathfrak{M}$  is the vertex set of a clique of  $G$ ; it remains to prove that this clique is maximal. Suppose that there exists  $M_1 \in \mathfrak{M}$  such that the clique whose vertex set is  $M_1$  is not maximal. Then there exists a vertex  $x \notin M_1$  which is joined by edges with all vertices of  $M_1$ . Let  $y \in M_1$ ; as  $y$  is joined with  $x$ , there exists a set  $M(y) \in \mathfrak{M}$  such that  $x \in M(y)$ ,  $y \in M(y)$ . Thus we can assign  $M(y)$  to each  $y \in M_1$ . Evidently  $M_1 \subseteq \bigcup_{y \in M_1} M(y)$ . As  $\mathfrak{M}$  is a  $\tau$ -covering, this implies  $\bigcap_{y \in M_1} M(y) \subseteq M_1$ . But  $x \in M(y)$  for each  $y \in M_1$ , therefore  $x \in \bigcap_{y \in M_1} M(y)$  and also  $x \in M_1$ , which is a contradiction.

**Theorem 2.** *Let  $G$  be an undirected graph with the vertex set  $V$ , let  $\mathfrak{M}$  be the system of all maximal independent sets of  $G$ . Then  $\mathfrak{M}$  is a  $\tau$ -covering of  $V$ . Conversely, if a set  $V$  and its  $\tau$ -covering  $\mathfrak{M}$  are given, then there exists a graph  $G$  such that  $V$  is its vertex set and  $\mathfrak{M}$  is the system of all maximal independent sets of  $G$ .*

*Proof.* Let  $\bar{G}$  be the complement of  $G$ . Then each maximal independent set of  $\bar{G}$  is the vertex set of a certain maximal clique in  $G$  and vice versa. Thus the assertion follows from Theorem 1.

*Remark.* The word “maximal” means here always “maximal with respect to the set inclusion” and not “with the maximal number of elements”.

As mentioned at the beginning, the  $\tau$ -coverings are of importance not only in these problems, but in all problems of finding the maximal subsets of a set with the property that any two elements of such a subset are in a certain symmetric binary relation. A particular case of a  $\tau$ -covering is a partition.

#### Reference

- [1] *I. Chajda, J. Niederle, B. Zelinka:* On existence conditions for compatible tolerances. Czech. Math. J. 26 (1976), 304–311.

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