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## TOLERANCES OF POLYUNARY ALGEBRAS

BOHDAN ZELINKA

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**Abstract.** A polyunary algebra is an ordered pair  $\mathfrak{U} = (A, \mathcal{S})$ , where  $A$  is a non-empty set and  $\mathcal{S}$  is a monoid consisting of unary operations on  $A$ . A tolerance on a polyunary algebra or on a monoid is defined similarly as a congruence, only the requirement of transitivity is omitted. In the paper the properties of tolerances on  $\mathfrak{U} = (A, \mathcal{S})$  and on  $\mathcal{S}$  are studied.

**Key words.** Polyunary algebra, monoid, tolerance, stable relation, congruence.

**Ms Classification.** 08 A 60, 08 A 30.

In [2] tolerances on monounary algebras, i.e. algebras with one unary operation, were studied. Here we shall investigate tolerances on polyunary algebras.

A polyunary algebra  $\mathfrak{U}$  is an ordered pair  $(A, \mathcal{F})$ , where  $A$  is a non-empty set (the support of  $\mathfrak{U}$ ) and  $\mathcal{F}$  is a non-empty set of unary operations on  $\mathfrak{U}$ , i.e. of mappings of  $A$  into  $A$ . If we speak about polyunary algebras, we suppose usually that  $\mathcal{F}$  contains at least two elements (otherwise the algebra would be monounary).

The operations from  $\mathcal{F}$  will be called fundamental operations on  $\mathfrak{U}$ . By their superpositions new operations are obtained; this superposition is associative. Thus the set  $\mathcal{F}$  generates a semigroup whose elements are unary operations acting on  $A$ . To the operations from this semigroup we may add the identity mapping  $e$  on  $A$ ; in this way a monoid  $\mathcal{S}$  is obtained. For our purposes it is reasonable to consider whole the monoid  $\mathcal{S}$  instead of the set  $\mathcal{F}$  of fundamental operations. Thus we shall consider a polyunary algebra  $\mathfrak{U}$  as an ordered pair  $(A, \mathcal{S})$ , where  $A$  is a non-empty set and  $\mathcal{S}$  is a monoid consisting of unary operations acting on  $A$ ; the operation on  $\mathcal{S}$  is the superposition of mappings.

We shall consider tolerances on  $\mathfrak{U}$  and on  $\mathcal{S}$ .

Let  $\mathfrak{U} = (A, \mathcal{S})$  be a polyunary algebra, let  $T$  be a reflexive and symmetric binary relation on  $A$ . If  $T$  has the property that  $(x, y) \in T$  implies  $(f(x), f(y)) \in T$  for any  $x \in A, y \in A, f \in \mathcal{S}$ , then  $T$  is called a tolerance on  $\mathfrak{U}$ .

Now let  $T$  be a reflexive and symmetric binary relation on  $\mathcal{S}$ . If  $T$  has the property that  $(x_1, x_2) \in T, (y_1, y_2) \in T$  imply  $(x_1 y_1, x_2 y_2) \in T$  for any elements  $x_1, x_2, y_1, y_2$  of  $\mathcal{S}$ , then  $T$  is called a tolerance on  $\mathcal{S}$ .

All tolerances on  $\mathfrak{A}$  form the lattice  $LT(\mathfrak{A})$ , all tolerances on  $\mathcal{S}$  form the lattice  $LT(\mathcal{S})$  (both with respect to set inclusion). In both  $LT(\mathfrak{A})$  and  $LT(\mathcal{S})$  the meet is the set intersection. In  $LT(\mathfrak{A})$  the join is the set union; this can be proved in the same way as in [1] for monounary algebras. In  $LT(\mathcal{S})$  the join is not equal to the set union in general.

First we shall investigate a simple case—the monogeneous polyunary algebras. A polyunary algebra  $\mathfrak{A} = (A, \mathcal{S})$  is called monogeneous, if there exists an element  $a \in A$  with the property that for each  $x \in A$  there exists an operation  $f_x \in \mathcal{S}$  such that  $x = f_x(a)$ . The element  $a$  will be called a generator of  $a$ .

We shall use also the concept of the left stable binary relation on a monoid. A binary relation  $R$  on a monoid  $\mathcal{S}$  is called left stable, if  $(x, y) \in R$  implies  $(zx, zy) \in R$  for any elements  $x, y, z$  of  $\mathcal{S}$ . The binary relations on  $\mathcal{S}$  which are reflexive, symmetric and left stable form a lattice which will be denoted by  $LS(\mathcal{S})$ .

**Theorem 1.** *Let  $\mathfrak{A} = (A, \mathcal{S})$  be a monogeneous polyunary algebra, let  $\mathcal{S}$  be a commutative monoid. Then there exists a one-to-one mapping  $\varphi$  of  $A$  onto  $\mathcal{S}$  such that  $\varphi(x) = f_x$  for each  $x \in A$ .*

*Proof.* For each  $x \in A$  there exists at least one operation from  $\mathcal{S}$  such that  $x$  is the image of  $a$  in this operation; namely the operation  $f_x$ . Let  $f, g$  be two operations from  $\mathcal{S}$  such that  $f(a) = g(a)$ . Let  $b$  be an arbitrary element of  $A$ ; as  $A$  is monogeneous, there exists  $h \in \mathcal{S}$  such that  $h(a) = b$ . Then we have:

$$f(b) = fh(a) = hf(a) = hg(a) = gh(a) = g(b).$$

As  $b$  was chosen arbitrarily, this implies  $f = g$ . Hence to each  $x \in A$  there exists exactly one operation  $f_x \in \mathcal{S}$  such that  $f_x(a) = x$  and the mapping  $\varphi$  is one-to-one.

**Theorem 2.** *Let  $\mathfrak{A} = (A, \mathcal{S})$  be a monogeneous polyunary algebra, let  $\mathcal{S}$  be a commutative monoid. Then  $LT(\mathfrak{A}) \cong LS(\mathcal{S})$ .*

*Proof.* Let  $\varphi$  be the mapping defined in Theorem 1. Let  $\bar{\varphi}$  be the mapping defined so that  $\bar{\varphi}(T)$  for  $T \in LT(\mathfrak{A})$  is the set of all pairs  $(\varphi(x), \varphi(y))$  for such pairs  $x, y$  that  $(x, y) \in T$ . Let  $(x, y) \in T$ . We have  $x = f_x(a), y = f_y(a), f_x = \varphi(x), f_y = \varphi(y)$ . As  $T$  is a tolerance on  $\mathfrak{A}$ , it contains all pairs  $(g(x), g(y))$  for  $g \in \mathcal{S}$ . We have  $(g(x), g(y)) = (gf_x(a), gf_y(a))$ . Hence for each  $g \in \mathcal{S}$  the relation  $\bar{\varphi}(T)$  contains the pair  $(gf_x, gf_y)$ ; as  $x, y$  were chosen arbitrarily, we see that  $\bar{\varphi}(T) \in LS(\mathcal{S})$  and  $\bar{\varphi}$  is a mapping of  $LT(\mathfrak{A})$  onto  $LS(\mathcal{S})$ . On the other hand, let  $U \in LS(\mathcal{S}), (h_1, h_2) \in U$ . As  $U$  is a left stable relation on  $\mathcal{S}$ , it contains all pairs  $(gh_1, gh_2)$  for  $g \in \mathcal{S}$ . Let  $u = h_1(a), v = h_2(a)$ . Then  $h_1 = \varphi(u), h_2 = \varphi(v), gh_1 = \varphi(g(u)), gh_2 = \varphi(g(v))$  and thus  $(gh_1, gh_2)$  is the image of  $(g(u), g(v))$  in the mapping  $\varphi$ . Hence  $\varphi^{-1}(U)$  with  $(u, v) = (h_1(a), h_2(a))$  contains also  $(g(u), g(v)) = (gh_1(a), gh_2(a))$ ; as  $u, v, g$  were chosen arbitrarily,  $\varphi^{-1}(U) \in LT(\mathfrak{A})$  and  $\varphi^{-1}$  is a mapping of  $LS(\mathcal{S})$  into  $LT(\mathfrak{A})$ . Combining these two results, we obtain that  $\bar{\varphi}$  is a one-to-one mapping of  $LT(\mathfrak{A})$  onto  $LS(\mathcal{S})$ . It is

easy to see that  $\bar{\varphi}$  preserves the lattice operations and thus it is an isomorphism of  $LT(\mathfrak{A})$  onto  $LS(\mathcal{S})$ .

If a left stable binary relation on a monoid  $\mathcal{S}$  is an equivalence, we shall call it a left stable equivalence. If  $\mathcal{S}$  is commutative, we may omit the word "left" and, according to [1], any stable equivalence is a congruence; hence  $LS(\mathcal{S}) = Con(\mathcal{S})$ .

**Theorem 3.** *Let  $\mathfrak{A} = (A, \mathcal{S})$  be a monogeneous polyunary algebra with the generator  $a$ . Then there exists a left stable equivalence  $R$  on  $S$  such that  $f(a) = g(a)$  if and only if  $(f, g) \in R$ .*

*Proof.* Let  $R$  be the relation on  $\mathcal{S}$  such that  $(f, g) \in R$  if and only if  $f(a) = g(a)$ ; it is evidently an equivalence. Let  $h \in \mathcal{S}$ . If  $(f, g) \in R$ , then  $f(a) = g(a)$  and thus also  $hf(a) = hg(a)$  and  $(hf, hg) \in R$ . Hence  $R$  is left stable, which was to be proved.

**Theorem 4.** *Let  $\mathfrak{A} = (A, \mathcal{S})$  be a monogeneous polyunary algebra. Then  $LT(\mathfrak{A})$  is isomorphic to a sublattice of  $LS(\mathcal{S})$ .*

*Proof.* Let  $R$  be the relation from Theorem 4. We have  $R \in LS(\mathcal{S})$ . Consider the set  $L^*$  of the products  $UR$  for  $U \in LS(\mathcal{S})$ . (We have  $(x, y) \in UR$ , if and only if there exists  $z$  such that  $(x, z) \in U$ ,  $(z, y) \in R$ ). Each element of  $L^*$  is evidently a left stable relation on  $\mathcal{S}$ ; the join and the meet of any two elements of  $L^*$  in  $LS(\mathcal{S})$  belongs to  $L^*$ . Thus  $L^*$  is a sublattice of  $LS(\mathcal{S})$ . Now we define a mapping  $\chi$  of the set of equivalence classes of  $R$  onto  $A$  such that if  $C$  is an equivalence class of  $R$ , then  $\chi(C) = f(a)$ , where  $f \in C$ ; obviously this does not depend on the choice of  $f$ . Now we define the mapping  $\bar{\chi}$  of  $L^*$  onto  $LT(A)$  in such a way that  $\bar{\chi}(U)$  for  $U \in L^*$  is the set of all pairs  $(\chi(C_1), \chi(C_2))$ , where  $C_1, C_2$  are classes of  $R$  with the property that  $(g_1, g_2) \in U$  for any  $g_1 \in C_1$  and  $g_2 \in C_2$ . The mapping  $\bar{\chi}$  is evidently an isomorphism of  $L^*$  onto  $LT(A)$ .

By the symbol  $Con(\mathfrak{A})$  (or  $Con(\mathcal{S})$ ) we denote the lattice of all congruences on  $\mathfrak{A}$  (or on  $\mathcal{S}$  respectively).

**Theorem 5.** *Let  $\mathfrak{A} = (A, \mathcal{S})$  be a monogeneous polyunary algebra, let  $\mathcal{S}$  be a commutative monoid. Then  $Con(\mathfrak{A}) \cong Con(\mathcal{S})$ .*

*Proof.* If we consider the mapping  $\bar{\varphi}$  from the proof of Theorem 2, we see that  $\bar{\varphi}$  maps transitive relations again onto transitive relations. Thus the restriction of  $\bar{\varphi}$  onto  $Con(\mathfrak{A})$  maps  $Con(\mathfrak{A})$  onto the set of stable equivalences of  $\mathcal{S}$  which is  $Con(\mathcal{S})$ .

We have described the case of monogeneous polyunary algebras. For polyunary algebras which are not monogeneous the situation is more complicated. We shall not describe the properties of  $LT(\mathfrak{A})$  in that case; we shall only point out the difference between monogeneous polyunary algebras and non-monogeneous ones. We limit our considerations to algebras with finitely many generators. Then evidently there exists a finite set  $G(\mathfrak{A})$  of independent generators of  $\mathfrak{A}$ , i.e. the subset of  $A$  with the property that to each  $x \in \mathfrak{A}$  there exists  $a \in G(\mathfrak{A})$  and  $f \in \mathcal{S}$  such that  $x = f(a)$  and

if  $a \in G(\mathfrak{A})$ ,  $b \in G(\mathfrak{A})$ ,  $a \neq b$ , then  $a \neq f(b)$  for any  $f \in \mathcal{S}$ . For each  $a \in G(\mathfrak{A})$  we denote  $A(a) = \{f(a) \mid f \in \mathcal{S}\}$  and we call this set the orbit of  $a$ .

Orbits of different generators may have a non-empty intersection. But the following theorem holds.

**Theorem 6.** *Let  $a, b$  be two different independent generators of  $\mathfrak{A} = (A, \mathcal{S})$ . If  $A(a) \cap A(b) \neq \emptyset$ , then  $A(a) \cap A(b)$  contains only elements of the form  $f(a)$ , where  $f$  does not divide the unit  $e$  of  $\mathcal{S}$ .*

*Proof.* Let  $x \in A(a) \cap A(b)$ ; then  $x = f(a) = g(b)$  for some elements  $f, g$  of  $\mathcal{S}$ . Suppose that  $f$  is a divisor of  $e$ ; then there exists an element  $f^{-1} \in \mathcal{S}$  such that  $f^{-1}f = e$ . But then  $a = f^{-1}(x) = f^{-1}g(b)$ ; as  $f^{-1}g \in \mathcal{S}$ , the generators  $a, b$  are not independent.

**Corollary.** *Let  $a, b$  be two different independent generators of  $\mathfrak{A} = (A, \mathcal{S})$ , let  $\mathcal{S}$  be a group. Then  $A(a) \cap A(b) = \emptyset$ .*

Note that any monounary algebra  $\mathfrak{A}$  can be considered as a polyunary algebra  $(A, \mathcal{S})$ , where  $\mathcal{S}$  is a monogeneous monoid. If the graph of  $\mathfrak{A}$  is a cycle or a two-way infinite directed path,  $\mathcal{S}$  may be considered as a cyclic group.

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