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# MEASURES OF VECTOR INFORMATION WITH THE BRANCHING PROPERTY

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It is known that, for a large class of monoids (S, \*), all solutions of the functional equation d(s, t) + d(s \* t, u) = d(s, t \* u) + d(t, u)

for  $\varDelta: S^2 \rightarrow G$  (with (G, +) any divisible abelian group) have representations of the form

$$f(s, t) = f(s) + f(t) - f(s * t) + \psi(s, t),$$

where  $\psi$  is antisymmetric and bi-additive. We show that this class is closed under the formation of direct products. This result is then used to characterize branching measures of vector information on strings of monoid elements.

#### 1. INTRODUCTION

An entropy in the "classical" (probabilistic) sense is a sequence  $(I_n)$  of mappings from the set of all *n*-ary complete probability distributions into the real numbers. More generally,  $I_n$  can be a function of several probability distributions, as is the case for directed divergence or inaccuracy. Discussions, properties, and characterizations of such information measures can be found in the book [2] by Aczél and Daróczy.

Measures of information in a different sense have been proposed and studied by the author [3], [4]. In [3], a measure of information is a function of a string of elements from a monoid. In [4], the information measure is a function of pairs of strings (quantities and "attractions") and appears as a utility function. In the present setting, a measure of information will be a function of *m*-vectors of strings.

The main results are contained in Sections 2 and 4. In Section 3, a fundamental functional equation is derived. Its general solution is found in Section 4 and used in Section 5 to prove the main result of Section 2.

### 2. BRANCHING MEASURES OF VECTOR INFORMATION

A measure of m-vector information (m = 1, 2, ...) is a sequence  $\mu_n: \sum_{j=1}^m S_j^n \to G$ (n = 3, 4, ...), where (G, +) is a divisible abelian group and each  $S_j$  is a commutative monoid (with identity  $e_j$ ) from a certain class **S** defined below. For notational convenience, we write the argument of  $\mu_n$  as if it belonged to  $(\sum_{j=1}^m S_j)^n$  instead of  $\sum_{j=1}^m S_j^n$ . The essential property for our measures is that of branching. A measure  $\mu_n$  of

The essential property for our measures is that of branching. A measure  $\mu_n$  of *m*-vector information is said to be *branching* if there are maps  $\Delta_{ni} : \underset{j=1}{\overset{m}{X}} S_j^2 \to G$ , for all i = 1, 2, ..., n - 1 (n = 3, 4, ...), such that

(2.1) 
$$\mu_n(\mathbf{v}_1, \mathbf{v}_2, \dots, \mathbf{v}_n) =$$

 $= \mu_n(\mathbf{v}_1, \ldots, \mathbf{v}_{i-1}, \mathbf{v}_i \circ \mathbf{v}_{i+1}, \mathbf{e}, \mathbf{v}_{i+2}, \ldots, \mathbf{v}_n) + \Delta_{ni}(\mathbf{v}_i, \mathbf{v}_{i+1}),$ 

for all  $\mathbf{v}_k \in \sum_{j=1}^m S_j$  (k = 1, 2, ..., n), where  $\mathbf{e} = (e_1, e_2, ..., e_m)$  and  $\mathbf{u} \circ \mathbf{w} = (u_1 w_1, u_2 w_2, ..., u_m w_m)$  for all  $(\mathbf{u}, \mathbf{w}) \in \sum_{j=1}^m S_j^2$ . For the sake of readability, the operations of all  $S_j$  are designated simply by juxtaposition, as this leads to no confusion here.

All monoids  $S_i$  are from the class **S** defined as follows.

**Definition 2.1.** A commutative monoid (S, \*) is said to belong to class **S** if all solutions  $\Delta : S^2 \to G((G, +)$  any divisible abelian group) of the functional equation

(2.2) 
$$\Delta(s,t) + \Delta(s*t,u) = \Delta(s,t*u) + \Delta(t,u),$$

for all  $(s, t, u) \in S^3$ , have a representation

(2.3) 
$$\Delta(s,t) = \delta(s) + \delta(t) - \delta(s*t) + \psi(s,t)$$

for some map  $\delta: S \to G$  and a map  $\psi: S^2 \to G$  which is *antisymmetric* 

(2.4) 
$$\psi(s,t) = -\psi(t,s), \quad \forall (s,t) \in S^2,$$

and bi-additive. (Additivity in the first variable, for example, means that

$$\psi(st, u) = \psi(s, u) + \psi(t, u), \quad \forall (s, t, u) \in S^3.$$

Equations (2.2) and (2.1) have been studied extensively and on many different domains. On branching measures of information, see [3] and [4] for results on strings of (m=) 1- and 2-vectors. On equation (2.2), see [8], [5], [6], [1], [7], [3], [4].

The following summarizes the above results on (2.2).

**Theorem 2.2.** A commutative monoid  $(S, \cdot)$  belongs to class **S** if  $(S, \cdot)$  is any

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of the following: idempotent, a monoid with zero, a thread, a group, the set of positive elements of an ordered group, a cancellative w-thread, a near-thread.

**Definition 2.3.** A w-thread (or thread in the wider sense) is a connected, totally ordered topological semigroup. A thread is a w-thread with a greatest and a least element, both of which are idempotent. Finally, a *near-thread* is a semigroup obtained by removing the zero from a w-thread  $(S^*, \cdot)$  which has a zero as least element and the property  $S^*$ .  $S^* = S^*$  (global idempotence).

The main result, which is proved in Section 5, is the following.

**Theorem 2.4.** The measure  $\mu_n : \underset{j=1}{\overset{\mathbf{M}}{\underset{j=1}{\overset{m}{\underset{j=1}{\overset{m}{\atop}}}}} S_j^n \to G$  (n = 3, 4, ...) of *m*-vector information has the (2.1) branching property, if and only if it admits a representation

$$(2.5) \qquad \mu_n(\mathbf{v}_1,\ldots,\mathbf{v}_n) = \varphi_{n0}(\mathbf{v}_1\circ\ldots\circ\mathbf{v}_n) + \sum_{i=1}^n \varphi_{ni}(\mathbf{v}_i) + \sum_{i=1}^{n-1} \sum_{k=i+1}^n \psi_n(\mathbf{v}_i,\mathbf{v}_k),$$

for all  $(\mathbf{v}_1, ..., \mathbf{v}_n) \in \sum_{j=1}^m S_j^n$ , for some maps  $\varphi_{ni} : \sum_{j=1}^m S_j \to G$  (i = 0, 1, ..., n) and  $\psi_n : \sum_{j=1}^m S_j^2 \to G$ , where  $\psi_n$  is antisymmetric and bi-additive in the *m*-tuples  $\mathbf{v}_i, \mathbf{v}_k$ . That is,  $\psi_n$  satisfies

(2.6) 
$$\psi_n(\boldsymbol{u}, \boldsymbol{w}) = -\psi_r(\boldsymbol{w}, \boldsymbol{u}), \quad \forall (\boldsymbol{u}, \boldsymbol{w}) \in \sum_{j=1}^m S_j^2$$
$$\psi_n(\boldsymbol{u} \circ \boldsymbol{v}, \boldsymbol{w}) = -\psi_n(\boldsymbol{u}, \boldsymbol{w}) + \psi_n(\boldsymbol{v}, \boldsymbol{w}), \quad \forall (\boldsymbol{u}, \boldsymbol{v}, \boldsymbol{w}) \in \sum_{j=1}^m S_j^3.$$

We can easily obtain the following consequence.

**Corollary 2.5.** A measure  $\mu_n$  of *m*-vector information is (2.1) branching and symmetric in its arguments  $\mathbf{v}_i$  (i = 1, ..., n), if and only if there exist maps  $\tilde{\varphi}_n, \varphi_n$ :  $: \bigotimes_{j=1}^m S_j \to G$  such that

(2.7) 
$$\mu_n(\mathbf{v}_1,\ldots,\mathbf{v}_n) = \tilde{\varphi}_n(\mathbf{v}_1\circ\ldots\circ\mathbf{v}_n) + \sum_{i=1}^n \varphi_n(\mathbf{v}_i).$$

Proof. Assume that  $\mu_n$  has the (2.1) branching property.  $\mu_n$  has representation (2.5) by Theorem 2.4. Interchanging  $\mathbf{v}_k$  with  $\mathbf{v}_{k+1}$   $(1 \le k \le n-1)$ , by the symmetry hypothesis, we get

(2.8) 
$$\varphi_{nk}(\mathbf{v}_{k}) + \varphi_{n,k+1}(\mathbf{v}_{k+1}) + \psi_{n}(\mathbf{v}_{k}, \mathbf{v}_{k+1}) = \\ = \varphi_{nk}(\mathbf{v}_{k+1}) + \varphi_{n,k+1}(\mathbf{v}_{k}) + \psi_{n}(\mathbf{v}_{k+1}, \mathbf{v}_{k}),$$

 $= \varphi_{nk}(\mathbf{v}_{k+1}) + \varphi_{n,k+1}(\mathbf{v}_k) + \psi_n(\mathbf{v}_{k+1}, \mathbf{v}_k),$ for all  $(\mathbf{v}_k, \mathbf{v}_{k+1}) \in \sum_{j=1}^m S_j^2$ , from (2.5). Letting  $\mathbf{v}_{k+1} = \mathbf{e}$  and using  $\psi_n(\mathbf{e}, \mathbf{v}_k) = \psi_n(\mathbf{v}_k, \mathbf{e}) =$ 

= 0 (which follow from bi-additivity), we have

(2.9) 
$$\varphi_{n,k+1}(\boldsymbol{u}) = \varphi_{nk}(\boldsymbol{u}) + c_{nk}, \quad \forall \boldsymbol{u} \in \sum_{j=1}^{m} S_j$$

Now (2.8) yields

$$\psi_n(\mathbf{v}_k, \mathbf{v}_{k+1}) = \psi_n(\mathbf{v}_{k+1}, \mathbf{v}_k), \quad \forall (\mathbf{v}_k, \mathbf{v}_{k+1}) \in \bigotimes_{j=1}^m S_j^2.$$

By the (2.6) antisymmetry of  $\psi_n$ , therefore, (2.10) $\psi_n = 0$ .

Thus. by (2.9) and (2.10), (2.5) becomes (2.7), where  $\varphi_n$  and  $\tilde{\varphi}_n$  are defined by  $\varphi_n := \varphi_{n1}, \tilde{\varphi}_n := \varphi_{n0} + \sum_{k=1}^{n-1} (n-k) c_{nk}.$ The converse is easy to check.

# 3. DERIVATION OF THE FUNDAMENTAL EQUATION

We begin by using equation (2.1) in two different ways for a given string ( $\mathbf{v}_1, \mathbf{v}_2, \dots$ ...,  $\mathbf{v}_n$ )  $\in X_{j=1}^m S_j^n$ . On one hand, apply (2.1) for i = k + 1 ( $1 \le k \le n - 2$ ), then for i = k; on the other hand, apply (2.1) first for i = k, then for i = k + 1, then for i = k again. Comparing the two results, we find that

(3.1) 
$$\Delta_{n,k+1}(\mathbf{v}_{k+1}, \mathbf{v}_{k+2}) + \Delta_{nk}(\mathbf{v}_{k}, \mathbf{v}_{k+1} \circ \mathbf{v}_{k+2}) =$$
$$= \Delta_{nk}(\mathbf{v}_{k}, \mathbf{v}_{k+1}) + \Delta_{n,k+1}(\mathbf{e}, \mathbf{v}_{k+2}) + \Delta_{nk}(\mathbf{v}_{k} \circ \mathbf{v}_{k+1}, \mathbf{v}_{k+2}),$$

for any  $(\mathbf{v}_k, \mathbf{v}_{k+1}, \mathbf{v}_{k+2}) \in \sum_{j=1}^m S_j^3$ . With  $\mathbf{v}_k = \mathbf{e}$ , (3.1) becomes

(3.2) 
$$\Delta_{n,k+1}(\mathbf{v}_{k+1},\mathbf{v}_{k+2}) = \Delta_{nk}(\mathbf{v}_{k+1},\mathbf{v}_{k+2}) + \Delta_{nk}(\mathbf{e},\mathbf{v}_{k+1}) + \Delta_{n,k+1}(\mathbf{e},\mathbf{v}_{k+2}) - \Delta_{nk}(\mathbf{e},\mathbf{v}_{k+1}\circ\mathbf{v}_{k+2}).$$

From (3.1), using (3.2), we now get (with  $\boldsymbol{u} = \boldsymbol{v}_k, \boldsymbol{v} = \boldsymbol{v}_{k+1}, \boldsymbol{w} = \boldsymbol{v}_{k+2}$ )

(3.3) 
$$\Delta_{nk}(\mathbf{v}, \mathbf{w}) - \Delta_{nk}(\mathbf{e}, \mathbf{v} \circ \mathbf{w}) + \Delta_{nk}(\mathbf{u}, \mathbf{v} \circ \mathbf{w}) =$$
$$= \Delta_{nk}(\mathbf{u}, \mathbf{v}) - \Delta_{nk}(\mathbf{e}, \mathbf{v}) + \Delta_{nk}(\mathbf{u} \circ \mathbf{v}, \mathbf{w}),$$

for all  $(u, v, w) \in X_{j=1}^{m} S_{j}^{3}$  and all k = 1, 2, ..., n - 1. (To get (3.3) for k = n - 1, use (3.1) and (3.2) for k = n - 2 and solve (3.2) for  $\Delta_{n,n-2}(\mathbf{v}_{n-1}, \mathbf{v}_n)$  instead of  $\Delta_{n,n-1}(\mathbf{v}_{n-1}, \mathbf{v}_n)$ ; then substitute for  $\Delta_{n,n-2}$  terms in (3.1).

Fix 
$$n \ge 2$$
 and  $k = 1$  (temporarily) and define  $F : \underset{j=1}{\overset{m}{\times}} S_j^2 \to G$  by  
 $F(\mathbf{u}, \mathbf{v}) := \Delta_{n1}(\mathbf{u}, \mathbf{v}) - \Delta_{n1}(\mathbf{e}, \mathbf{v}), \quad \forall (\mathbf{u}, \mathbf{v}) \in \underset{j=1}{\overset{m}{\times}} S_j^2.$ 

Then, by (3.3), F satisfies

(3.4) 
$$F(\boldsymbol{u},\boldsymbol{v}) + F(\boldsymbol{u}\circ\boldsymbol{v},\boldsymbol{w}) = F(\boldsymbol{u},\boldsymbol{v}\circ\boldsymbol{w}) + F(\boldsymbol{v},\boldsymbol{w})$$

for all  $(\boldsymbol{u}, \boldsymbol{v}, \boldsymbol{w}) \in \sum_{j=1}^{m} S_{j}^{3}$ , and  $\Delta_{n1}$  is given by

(3.5) 
$$\Delta_{n1}(\boldsymbol{u},\boldsymbol{v}) = F(\boldsymbol{u},\boldsymbol{v}) + \Delta_{n1}(\boldsymbol{e},\boldsymbol{v}), \quad \forall (\boldsymbol{u},\boldsymbol{v}) \in \overset{\sim}{\underset{j=1}{\overset{\sim}{\longrightarrow}}} S_{j}^{2},$$

for an arbitrary map  $\Delta_{n1}(\mathbf{e}, \cdot) : \underset{j=1}{\overset{m}{\times}} S_j \to G.$ 

Thus we have proved the following.

**Lemma 3.1.** If  $\mu_n$  (n = 3, 4, ...) is a (2.1) branching measure of *m*-vector information, then the branching functions  $\Delta_{ni}$  (i = 2, 3, ..., n - 1) can be obtained recursively from  $\Delta_{n1}$  and arbitrary one-place functions  $\Delta_{ni}$  ( $\mathbf{e}, \cdot$ ) through (3.2). Moreover, m

$$\Delta_{n1}$$
 is given by (3.5) for an arbitrary solution,  $F : \underset{j=1}{\mathsf{X}} S_j^2 \to G$ , of (3.4).

Our immediate goal is, therefore, to solve (3.4).

## 4. SOLUTION OF THE FUNDAMENTAL EQUATION

The principal tool to be used in solving equation (3.4) is the following.

**Lemma 4.1.** If  $(X, \oplus)$  is a commutative monoid, and if  $(S, *) \in S$ , then a map  $F : (X \times S)^2 \to G$  (with (G, +) a divisible abelian group) satisfies

$$(4.1) F(x, r; y, s) + F(x \oplus y, r * s; z, t) =$$

$$= F(x, r; y \oplus z, s * t) + F(y, s; z, t)$$

for all  $(x, y, z) \in X^3$  and all  $(r, s, t) \in S^3$ , if and only if there exist maps  $\varphi : X \times S \to G$ ,  $\psi : (X \times S)^2 \to G$ ,  $\tilde{F} : X^2 \to G$  such that

(4.2) 
$$F(x, r; y, s) = \tilde{F}(x, y) + \varphi(x, r) + \varphi(y, s) - \varphi(x \oplus y, r * s) + \psi(x, r; y, s),$$

for all  $(x, y) \in X^2$ ,  $(r, s) \in S^2$ , where  $\psi$  is (2.6) antisymmetric and bi-additive, and  $\tilde{F}$  satisfies

$$(4.3) \qquad \tilde{F}(x, y) + \tilde{F}(x \oplus y, z) = \tilde{F}(x, y \oplus z) + \tilde{F}(y, z), \quad \forall (x, y, z) \in X^3.$$

Proof. Lemma 4.1 in [3].

With Lemma 4.1, we obtain the following result.

**Theorem 4.2.** Let  $S_j \in S$ , j = 1, 2, ..., m; let (G, +) be a divisible abelian group.

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A map  $F : \underset{m}{\overset{m}{X}} S_j^2 \to G$  satisfies (3.4) for all  $(\boldsymbol{u}, \boldsymbol{v}, \boldsymbol{w}) \in \underset{j=1}{\overset{m}{X}} S_j^3$ , if and only if there exist a map  $\varphi : \underset{j=1}{\overset{m}{X}} S_j \to G$  and an (2.6) antisymmetric bi-additive map  $\psi : \underset{j=1}{\overset{m}{X}} S_j^2 \to G$  for which

(4.4) 
$$F(\boldsymbol{u},\boldsymbol{v}) = \varphi(\boldsymbol{u}) + \varphi(\boldsymbol{v}) - \varphi(\boldsymbol{u} \circ \boldsymbol{v}) + \psi(\boldsymbol{u},\boldsymbol{v}),$$

for all  $(\boldsymbol{u}, \boldsymbol{v}) \in \overset{m}{\underset{j=1}{\mathsf{X}}} S_{j}^{2}$ .

Proof. Let F satisfy (3.4). In this direction, the proof is by induction on m. For m = 1, it is simply a restatement of Definition 2.1.

Assume the truth of the theorem for m = n. Let  $\tilde{\mathbf{v}}_i$  denote  $(v_1, v_2, ..., v_i) =$  the first *i* coordinates of  $\mathbf{v} \in \bigotimes_{j=1}^m S_j$ , for i = 1, 2, ..., m - 1. Also, let  $\tilde{\mathbf{u}}_i \oplus \tilde{\mathbf{v}}_i$  denote  $(u_1v_1, u_2v_2, ..., u_iv_i) =$  the restriction of  $\mathbf{u} \circ \mathbf{v}$  to its first *i* coordinates (i = 1, 2, ..., m - 1), for any  $(\mathbf{u}, \mathbf{v}) \in \bigotimes_{j=1}^m S_j^2$ . Then, for m = n + 1, (3.4) can be written

$$(3.4) F(\tilde{\boldsymbol{u}}_n, u_{n+1}; \tilde{\boldsymbol{v}}_n, v_{n+1}) + F(\tilde{\boldsymbol{u}}_n \oplus \tilde{\boldsymbol{v}}_n, u_{n+1}v_{n+1}; \tilde{\boldsymbol{w}}_n, w_{n+1}) = \\ = F(\tilde{\boldsymbol{u}}_n, u_{n+1}; \tilde{\boldsymbol{v}}_n \oplus \tilde{\boldsymbol{w}}_n, v_{n+1}w_{n+1}) + F(\tilde{\boldsymbol{v}}_n, v_{n+1}; \tilde{\boldsymbol{w}}_n, w_{n+1})$$

Now, applying Lemma 4.1 with  $X = \bigvee_{j=1}^{n} S_j$  and  $S = S_{n+1}$ , we get

$$(4.5) \quad F(\tilde{\boldsymbol{u}}_n, \boldsymbol{u}_{n+1}; \tilde{\boldsymbol{v}}_n, \boldsymbol{v}_{n+1}) = \tilde{F}(\tilde{\boldsymbol{u}}_n, \tilde{\boldsymbol{v}}_n) + \hat{\boldsymbol{\varphi}}(\boldsymbol{u}) + \hat{\boldsymbol{\varphi}}(\boldsymbol{v}) - \hat{\boldsymbol{\varphi}}(\boldsymbol{u} \circ \boldsymbol{v}) + \hat{\boldsymbol{\psi}}(\boldsymbol{u}, \boldsymbol{v}),$$

where  $\tilde{F}$  satisfies (4.3) and  $\hat{\psi}$  is (2.6) antisymmetric and bi-additive. But, by the induction hypothesis,  $\tilde{F}$  has a representation of the form

(4.6) 
$$\widetilde{F}(\widetilde{\boldsymbol{u}}_n, \widetilde{\boldsymbol{v}}_n) = \widetilde{\varphi}(\widetilde{\boldsymbol{u}}_n) + \widetilde{\varphi}(\widetilde{\boldsymbol{v}}_n) - \widetilde{\varphi}(\widetilde{\boldsymbol{u}}_n \oplus \widetilde{\boldsymbol{v}}_n) + \widetilde{\psi}(\widetilde{\boldsymbol{u}}_n, \widetilde{\boldsymbol{v}}_n)$$

where  $\tilde{\psi}$  is (2.6) antisymmetric and bi-additive. Combining (4.6) with (4.5) and defining  $\varphi(\mathbf{u}) := \hat{\varphi}(\mathbf{u}) + \tilde{\varphi}(\tilde{\mathbf{u}}_n), \psi(\mathbf{u}, \mathbf{v}) := \hat{\psi}(\mathbf{u}, \mathbf{v}) + \tilde{\psi}(\tilde{\mathbf{u}}_n, \tilde{\mathbf{v}}_n)$  for all  $(\mathbf{u}, \mathbf{v}) \in \sum_{j=1}^{n+1} S_j^2$ , we have (4.4) for m = n + 1 (with  $\psi$  (2.6) antisymmetric and bi-additive), as required.

The converse is easy to check.

Remark. Theorem 4.2 says that S is closed under the formation of direct products.

#### 5. PROOF OF THEOREM 2.4.

Let  $\mu_n$  (n = 3, 4, ...) be a (2.1) branching measure of *m*-vector information. By Lemma 3.1 and Theorem 4.2,  $\Delta_{n1}$  has a representation in the form

$$\Delta_{n1}(\boldsymbol{u},\boldsymbol{v}) = \varphi_{n1}(\boldsymbol{u}) + \varphi_{n2}(\boldsymbol{v}) - \varphi_{n1}(\boldsymbol{u} \circ \boldsymbol{v}) + \psi_n(\boldsymbol{u},\boldsymbol{v})$$

for all  $(\mathbf{u}, \mathbf{v}) \in \sum_{j=1}^{m} S_{j}^{2}$ , where  $\varphi_{n2}(\mathbf{v}) := \varphi_{n1}(\mathbf{v}) + \Delta_{n1}(\mathbf{e}, \mathbf{v})$  and  $\psi_{n}$  is (2.6) antisymmetric and bi-additive. Furthermore,  $\Delta_{n2}, \Delta_{n3}, \dots, \Delta_{n,n-1}$  are defined recursively through (3.2), giving

(5.1) 
$$\Delta_{ni}(\boldsymbol{u},\boldsymbol{v}) = \varphi_{ni}(\boldsymbol{u}) + \varphi_{n,i+1}(\boldsymbol{v}) - \varphi_{ni}(\boldsymbol{u} \circ \boldsymbol{v}) + \psi_n(\boldsymbol{u},\boldsymbol{v})$$

for all  $(u, v) \in X_{j=1}^{S_{j}}$  and all i = 1, 2, ..., n - 1, where  $\varphi_{ni}$  is defined recursively by  $\varphi_{n,i+1}(u) := \varphi_{ni}(u) + \Delta_{ni}(e, u)$ . We have also used the fact that  $\psi_{n}(e, u) = 0$ for any  $u \in X_{j}$ , which follows from additivity in the first variable.

Now, by (2.1) and (5.1), we have

$$\mu_{n}(\mathbf{v}_{1}, \mathbf{v}_{2}, \dots, \mathbf{v}_{n}) = \mu_{n}(\mathbf{v}_{1}, \dots, \mathbf{v}_{n-2}, \mathbf{v}_{n-1} \circ \mathbf{v}_{n}, \mathbf{e}) + \varphi_{n,n-1}(\mathbf{v}_{n-1}) + \varphi_{nn}(\mathbf{v}_{n}) - \varphi_{n,n-1}(\mathbf{v}_{n-1} \circ \mathbf{v}_{n}) + \psi_{n}(\mathbf{v}_{n-1}, \mathbf{v}_{n}) = \dots$$
$$\dots = \mu_{n}(\mathbf{v}_{1} \circ \mathbf{v}_{2} \circ \dots \circ \mathbf{v}_{n}, \mathbf{e}, \mathbf{e}, \dots, \mathbf{e}) + \sum_{i=1}^{n} \varphi_{ni}(\mathbf{v}_{i}) - \varphi_{n1}(\mathbf{v}_{1} \circ \mathbf{v}_{2} \circ \dots \circ \mathbf{v}_{n}) + \sum_{i=1}^{n-1} \sum_{k=i+1}^{n} \psi_{n}(\mathbf{v}_{i}, \mathbf{v}_{k}),$$

where we have also used the additivity of  $\psi_n$  in the second variable. Defining  $\varphi_{n0}(\mathbf{u}) := = \mu_n(\mathbf{u}, \mathbf{e}, \mathbf{e}, ..., \mathbf{e}) - \varphi_{n1}(\mathbf{u})$ , we obtain the asserted form (2.5).

Again, the converse is easy to check, and Theorem 2.4 is established.

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