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# IDEALS OF NONCOMMUTATIVE $D R \ell$-MONOIDS 

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Abstract. In this paper, we introduce the concept of an ideal of a noncommutative dually residuated lattice ordered monoid and we show that congruence relations and certain ideals are in a one-to-one correspondence.

Keywords: dually residuated lattice ordered monoid, ideal, normal ideal
MSC 2000: 06F05, 06D35

## 1. Introduction

Commutative $D R \ell$-monoids (called $D R \ell$-semigroups) were introduced by K.L. N. Swamy in [11] as a common generalization of commutative $\ell$-groups and Brouwerian algebras. A noncommutative extension of $D R \ell$-semigroups is mentioned in [12], but the present definition, due to [8], is more general. In fact, Swamy's noncommutative $D R \ell$-semigroup was considered as an algebra $(A,+, 0, \vee, \wedge,-)$, where "-" agrees with " $\downarrow$ ".

Definition. An algebra $\mathfrak{A}=(A,+, 0, \vee, \wedge, \rightharpoonup, \leftharpoondown)$ is a dually residuated lattice ordered monoid, or simply a $D R \ell$-monoid, iff
(1) $(A,+, 0, \vee, \wedge)$ is an $\ell$-monoid, that is, $(A,+, 0)$ is a monoid, $(A, \vee, \wedge)$ is a lattice and, for any $x, y, s, t \in A$, the following distributive laws are satisfied:

$$
\begin{aligned}
& s+(x \vee y)+t=(s+x+t) \vee(s+y+t), \\
& s+(x \wedge y)+t=(s+x+t) \wedge(s+y+t)
\end{aligned}
$$

(2) for any $x, y \in A, x \rightharpoonup y$ is the least $s \in A$ such that $s+y \geqslant x$, and $x \leftharpoondown y$ is the least $t \in A$ such that $y+t \geqslant x$;
(3) $\mathfrak{A}$ fulfils the identities

$$
\begin{gathered}
((x \rightharpoonup y) \vee 0)+y \leqslant x \vee y, \quad y+((x \leftharpoondown y) \vee 0) \leqslant x \vee y, \\
x \rightharpoonup x \geqslant 0, \quad x \leftharpoondown x \geqslant 0 .
\end{gathered}
$$

Note that the condition (2) is equivalent to the following system of identities (see [10]):

$$
\begin{aligned}
(x \rightharpoonup y)+y \geqslant x, & y+(x \leftharpoondown y) \geqslant x, \\
x \rightharpoonup y \leqslant(x \vee z) \rightharpoonup y, & x \leftharpoondown y \leqslant(x \vee z) \leftharpoondown y, \\
(x+y) \rightharpoonup y \leqslant x, & (y+x) \leftharpoondown y \leqslant x .
\end{aligned}
$$

Also, Swamy indroduced the notion of an ideal of a commutative $D R \ell$-monoid as a nonempty subset closed under " + " containing with any $x$ also all $y$ such that $y * 0 \leqslant x * 0$ (where $a * b=(a-b) \vee(b-a)$ is the symmetric difference of $a$ and $b$ ). In addition, ideals and congruence relations are in a one-to-one correspondence; for any ideal $I$ of a commutative $D R \ell$-monoid $\mathfrak{A}$, the corresponding congruence relation $\Theta(I)$ is defined by $\langle x, y\rangle \in \Theta(I)$ iff $x * y \in I$.

We generalize the notion of an ideal and, in an attempt to describe congruence kernels of noncommutative $D R \ell$-monoids, we introduce normal ideals which in the case that a $D R \ell$-monoid is an $\ell$-group coincide with $\ell$-ideals.

The concepts of distance functions and normal ideals are motivated by $G M V$ algebras (pseudo $M V$-algebras) which are included among $D R \ell$-monoids (see [10]).

Recall that $G M V$-algebras were introduced by J. Rachůnek in [10] (and independently by G. Georgescu and A. Iorgulescu in [4] under the name pseudo $M V$ algebras) to be a noncommutative generalization of $M V$-algebras. As shown in [10], if $(A, \oplus, \neg, \sim, 0,1)$ is a $G M V$-algebra with the additional binary operation " $\odot$ " defined by $x \odot y=\sim(\neg x \oplus \neg y)$ and if we put $x \vee y=(\neg x \odot y) \oplus x, x \wedge y=(\neg x \oplus y) \odot x$, $x \rightharpoonup y=\neg y \odot x$, and $x \leftharpoondown y=x \odot \sim y$, then $(A, \oplus, 0, \vee, \wedge, \rightharpoonup, \leftharpoondown)$ is a bounded $D R \ell$-monoid whose greatest element is 1 .

## 2. Distance functions, absolute value

Definition. Let $\mathfrak{A}$ be a $D R \ell$-monoid. We define the distance functions by

$$
\begin{aligned}
d_{1}(x, y) & :=(x \rightharpoonup y) \vee(y \rightharpoonup x), \\
d_{2}(x, y) & :=(x \leftharpoondown y) \vee(y \leftharpoondown x),
\end{aligned}
$$

for any $x, y \in A$.

Further, for each $x \in A,|x|:=d_{1}(x, 0)$ is the absolute value of $x$, and $x^{+}:=x \vee 0$ is the positive part of $x$.

Before stating some results concerning the above notions, it is useful to mention basic properties of $D R \ell$-monoids.

Lemma 1 [8, Lemmas 1.1.7, 1.1.5, 1.1.8, 1.1.12]. In any $D R \ell$-monoid we have
(1) $x \vee y=(x \rightharpoonup y)^{+}+y=y+(x \leftharpoondown y)^{+}$;
(2) $x \rightharpoonup x=x \leftharpoondown x=0$;
(3) $x \geqslant y \Longrightarrow x \rightharpoonup z \geqslant y \rightharpoonup z, x \leftharpoondown z \geqslant y \leftharpoondown z, z \rightharpoonup x \leqslant z \rightharpoonup y$, and $z \leftharpoondown x \leqslant z \leftharpoondown y ;$
(4) $x \rightharpoonup(y+z)=(x \rightharpoonup z) \rightharpoonup y, x \leftharpoondown(y+z)=(x \leftharpoondown y) \leftharpoondown z$.

Lemma 2. Suppose that all joins and meets on the left-hand side exist. Then the following is valid:
(1) $x+\bigwedge_{\lambda \in \Lambda} y_{\lambda}=\bigwedge_{\lambda \in \Lambda}\left(x+y_{\lambda}\right), \bigwedge_{\lambda \in \Lambda} y_{\lambda}+x=\bigwedge_{\lambda \in \Lambda}\left(y_{\lambda}+x\right)$;
(2) $x \rightharpoonup \bigwedge_{\lambda \in \Lambda} y_{\lambda}=\bigvee_{\lambda \in \Lambda}\left(x \rightharpoonup y_{\lambda}\right), x \leftharpoondown \bigwedge_{\lambda \in \Lambda} y_{\lambda}=\bigvee_{\lambda \in \Lambda}\left(x \leftharpoondown y_{\lambda}\right)$;
(3) $\bigvee_{\lambda \in \Lambda} x_{\lambda} \rightharpoonup y=\bigvee_{\lambda \in \Lambda}\left(x_{\lambda} \rightharpoonup y\right), \bigvee_{\lambda \in \Lambda} x_{\lambda} \leftharpoondown y=\bigvee_{\lambda \in \Lambda}\left(x_{\lambda} \leftharpoondown y\right)$;
(4) $x \vee \bigwedge_{\lambda \in \Lambda} y_{\lambda}=\bigwedge_{\lambda \in \Lambda}\left(x \vee y_{\lambda}\right)$.

Remark. (2) and (3) extend [8, Lemma 1.1.9] for the arbitrary existing joins and meets, respectively.

Proof. (1) From $y_{\lambda} \geqslant \bigwedge_{\lambda \in \Lambda} y_{\lambda}$ it follows that $x+y_{\lambda} \geqslant x+\bigwedge_{\lambda \in \Lambda} y_{\lambda}$, for any $\lambda \in \Lambda$. Conversely, if there is $z \in A$ with $x+y_{\lambda} \geqslant z$, for all $\lambda \in \Lambda$, then $y_{\lambda} \geqslant z \leftharpoondown x$, for all $\lambda \in \Lambda$, and so $\bigwedge_{\lambda \in \Lambda} y_{\lambda} \geqslant z \leftharpoondown x$, whence $x+\bigwedge_{\lambda \in \Lambda} y_{\lambda} \geqslant z$, proving the first identity in (1). The rest of (1), and (2) and (3) have a similar proof.
(4) Obviously, $x \vee \bigwedge_{\lambda \in \Lambda} y_{\lambda} \leqslant x \vee y_{\lambda}$ for all $\lambda \in \Lambda$. Choose $z \in A$ such that $z \leqslant x \vee y_{\lambda}=\left(x \rightharpoonup y_{\lambda}\right)^{+}+y_{\lambda}$ for each $\lambda \in \Lambda$. Then $y_{\lambda} \geqslant z \leftharpoondown\left(x \rightharpoonup y_{\lambda}\right)^{+}$, for all $\lambda \in \Lambda$, and therefore $\bigwedge_{\lambda \in \Lambda} y_{\lambda} \geqslant z \leftharpoondown\left(x \rightharpoonup \bigwedge_{\lambda \in \Lambda} y_{\lambda}\right)^{+}$which gives $z \leqslant(x \rightharpoonup$ $\left.\bigwedge_{\lambda \in \Lambda} y_{\lambda}\right)^{+}+\bigwedge_{\lambda \in \Lambda} y_{\lambda}=x \vee \bigwedge_{\lambda \in \Lambda} y_{\lambda}$.

Corollary 3 [8, Theorem 1.1.23]. For any $D R \ell$-monoid $\mathfrak{A}$, the lattice $\mathfrak{L}(\mathfrak{A})=$ $(A, \vee, \wedge)$ is distributive.

Lemma 4 [8, Lemma 1.1.11]. For all $x, y$ of any $D R \ell$-monoid, it holds

$$
\begin{aligned}
& (x \rightharpoonup y) \vee(y \rightharpoonup x)=(x \vee y) \rightharpoonup(x \wedge y), \\
& (x \leftharpoondown y) \vee(y \leftharpoondown x)=(x \vee y) \leftharpoondown(x \wedge y) .
\end{aligned}
$$

Proof. Using Lemma 2, (2) and (3), we obtain $(x \vee y) \rightharpoonup(x \wedge y)=(x \rightharpoonup$ $y) \vee(y \rightharpoonup x) \vee 0$. However, $(x \rightharpoonup y) \vee(y \rightharpoonup x) \geqslant(x \rightharpoonup(x \vee y)) \vee(y \rightharpoonup(x \vee y))=$ $(x \vee y) \rightharpoonup(x \vee y)=0$, again by Lemma 2 .

Lemma 5 [8, Lemma 1.1.15]. If $x \geqslant y \geqslant z$ then

$$
(x \rightharpoonup y)+(y \rightharpoonup z)=x \rightharpoonup z \quad \text { and } \quad(y \leftharpoondown z)+(x \leftharpoondown y)=x \leftharpoondown z .
$$

Proof. If $y \geqslant z$ then $y \rightharpoonup z \geqslant 0$ and $y=y \vee z=(y \rightharpoonup z)^{+}+z=(y \rightharpoonup z)+z$. Hence $x \rightharpoonup y=x \rightharpoonup((y \rightharpoonup z)+z)=(x \rightharpoonup z) \rightharpoonup(y \rightharpoonup z)$. Similarly, $x \geqslant y$ entails $x \rightharpoonup z \geqslant y \rightharpoonup z$ which yields $x \rightharpoonup z=((x \rightharpoonup z) \rightharpoonup(y \rightharpoonup z))+(y \rightharpoonup z)$. Summarizing, $x \rightharpoonup z=(x \rightharpoonup y)+(y \rightharpoonup z)$.

Lemma 6 [8, Lemmas 1.1.5, 1.1.13]. The following holds in any $D R \ell$-monoid:
(1) $0 \rightharpoonup x=0 \leftharpoondown x$,
(2) $(x \rightharpoonup y)+(y \rightharpoonup z) \geqslant x \rightharpoonup z$,
(3) $(y \leftharpoondown z)+(x \leftharpoondown y) \geqslant x \leftharpoondown z$.

Proof. (1) From $(x+(0 \rightharpoonup x))+x=x+((0 \rightharpoonup x)+x) \geqslant x+0=x$ it follows that $x+(0 \rightharpoonup x) \geqslant x \rightharpoonup x=0$, whence $0 \rightharpoonup x \geqslant 0 \leftharpoondown x$. Similarly, $0 \leftharpoondown x \geqslant 0 \rightharpoonup x$.
(2) and similarly $(3)(x \rightharpoonup y)+(y \rightharpoonup z)+z \geqslant(x \rightharpoonup y)+y \geqslant x$ implies $(x \rightharpoonup$ $y)+(y \rightharpoonup z) \geqslant x \rightharpoonup z$.

Applying (2) and (3), we immediately get

Lemma 7. In every $D R \ell$-monoid we have
(1) $y \rightharpoonup x \geqslant(z \rightharpoonup x) \leftharpoondown(z \rightharpoonup y)$,
(2) $y \leftharpoondown x \geqslant(z \leftharpoondown x) \rightharpoonup(z \leftharpoondown y)$,
(3) $y \rightharpoonup x \geqslant(y \rightharpoonup z) \rightharpoonup(x \rightharpoonup z)$,
(4) $y \leftharpoondown x \geqslant(y \leftharpoondown z) \leftharpoondown(x \leftharpoondown z)$.

Proposition 8. The distance functions have the following properties:
(1) $d_{1}(x, y)=d_{1}(y, x)$,
(2) $d_{2}(x, y)=d_{2}(y, x)$,
(3) $d_{1}(x, 0)=d_{2}(x, 0)$,
(4) $d_{1}(x, y)=(x \rightharpoonup y)^{+}+(y \rightharpoonup x)^{+}$,
(5) $d_{2}(x, y)=(y \leftharpoondown x)^{+}+(x \leftharpoondown y)^{+}$,
(6) $d_{1}(x, y) \geqslant 0$ with $d_{1}(x, y)=0$ iff $x=y$,
(7) $d_{2}(x, y) \geqslant 0$ with $d_{2}(x, y)=0$ iff $x=y$,
(8) $d_{1}(x, y) \leqslant d_{1}(x, z)+d_{1}(z, y)+d_{1}(x, z)$,
(9) $d_{1}(x, y) \leqslant d_{1}(z, y)+d_{1}(x, z)+d_{1}(z, y)$,
(10) $d_{2}(x, y) \leqslant d_{2}(x, z)+d_{2}(z, y)+d_{2}(x, z)$,
(11) $d_{2}(x, y) \leqslant d_{2}(z, y)+d_{2}(x, z)+d_{2}(z, y)$,
(12) $d_{1}(x, y) \vee d_{1}(s, t) \geqslant d_{1}(x \vee s, y \vee t), d_{1}(x \wedge s, y \wedge t)$,
(13) $d_{2}(x, y) \vee d_{2}(s, t) \geqslant d_{2}(x \vee s, y \vee t), d_{2}(x \wedge s, y \wedge t)$,
(14) $d_{2}(z \rightharpoonup x, z \rightharpoonup y) \leqslant d_{1}(x, y)$,
(15) $d_{1}(z \leftharpoondown x, z \leftharpoondown y) \leqslant d_{2}(x, y)$,
(16) $d_{1}(x \rightharpoonup z, y \rightharpoonup z) \leqslant d_{1}(x, y)$,
(17) $d_{2}(x \leftharpoondown z, y \leftharpoondown z) \leqslant d_{2}(x, y)$.

Proof. Obviously, (1) and (2) hold; (3) follows by Lemma 6 (1). To check (4), and similarly (5), we compute

$$
\begin{array}{rlrl}
d_{1}(x, y) & =(x \rightharpoonup y) \vee(y \rightharpoonup x)=(x \vee y) \rightharpoonup(x \wedge y) & & \text { by Lemma } 4 \\
& =[(x \vee y) \rightharpoonup y]+[y \rightharpoonup(x \wedge y)] & & \text { by Lemma } 5 \\
& =[(x \rightharpoonup y) \vee(y \rightharpoonup y)]+[(y \rightharpoonup x) \vee(y \rightharpoonup y)] & & \text { by Lemma } 2 \\
& =[(x \rightharpoonup y) \vee 0]+[(y \rightharpoonup x) \vee 0] & & \\
& =(x \rightharpoonup y)^{+}+(y \rightharpoonup x)^{+} . &
\end{array}
$$

Further, (6) follows from (4) and (7) from (5), respectively, since

$$
d_{1}(x, y)=(x \rightharpoonup y)^{+}+(y \rightharpoonup x)^{+} \geqslant 0 .
$$

It is clear that $x=y$ entails $d_{1}(x, y)=0$. Conversely, if

$$
d_{1}(x, y)=(x \rightharpoonup y)^{+}+(y \rightharpoonup x)^{+}=0
$$

then $(x \rightharpoonup y)^{+}=(y \rightharpoonup x)^{+}=0$. Hence $x \rightharpoonup y \leqslant 0$ and $y \rightharpoonup x \leqslant 0$, and so $x \leqslant y$ and $y \leqslant x$, thus $x=y$.

Now, we will prove (8) (similarly (9), (10) and (11)):

$$
\begin{aligned}
d_{1}(x, z)+ & d_{1}(z, y)+d_{1}(x, z) \\
= & {[(x \rightharpoonup z) \vee(z \rightharpoonup x)]+[(z \rightharpoonup y) \vee(y \rightharpoonup z)]+[(x \rightharpoonup z) \vee(z \rightharpoonup x)] } \\
= & {[(x \rightharpoonup z)+(z \rightharpoonup y)+(x \rightharpoonup z)] \vee[(x \rightharpoonup z)+(z \rightharpoonup y)+(z \rightharpoonup x)] } \\
& \vee[(x \rightharpoonup z)+(y \rightharpoonup z)+(x \rightharpoonup z)] \vee[(x \rightharpoonup z)+(y \rightharpoonup z)+(z \rightharpoonup x)] \\
& \vee[(z \rightharpoonup x)+(z \rightharpoonup y)+(x \rightharpoonup z)] \vee[(z \rightharpoonup x)+(z \rightharpoonup y)+(z \rightharpoonup x)] \\
& \vee[(z \rightharpoonup x)+(y \rightharpoonup z)+(x \rightharpoonup z)] \vee[(z \rightharpoonup x)+(y \rightharpoonup z)+(z \rightharpoonup x)] \\
\geqslant & {[(x \rightharpoonup z)+(z \rightharpoonup y)+(x \rightharpoonup z)] \vee[(x \rightharpoonup z)+(z \rightharpoonup y)+(z \rightharpoonup x)] } \\
& \vee[(x \rightharpoonup z)+(y \rightharpoonup z)+(z \rightharpoonup x)] \vee[(z \rightharpoonup x)+(y \rightharpoonup z)+(z \rightharpoonup x)] \\
= & {[((x \rightharpoonup z)+(z \rightharpoonup y))+((x \rightharpoonup z) \vee(z \rightharpoonup x))] } \\
& \vee[((x \rightharpoonup z) \vee(z \rightharpoonup x))+((y \rightharpoonup z)+(z \rightharpoonup y))] \\
& (u \operatorname{sing}(x \rightharpoonup z) \vee(z \rightharpoonup x) \geqslant 0, \text { by }(4)) \\
\geqslant & {[(x \rightharpoonup z)+(z \rightharpoonup y)] \vee[(y \rightharpoonup z)+(z \rightharpoonup x)] } \\
\geqslant & (x \rightharpoonup y) \vee(y \rightharpoonup x)=d_{1}(x, y) .
\end{aligned}
$$

Let us verify (12):

$$
\begin{aligned}
d_{1}(x, y) \vee d_{1}(s, t)= & (x \rightharpoonup y) \vee(y \rightharpoonup x) \vee(s \rightharpoonup t) \vee(t \rightharpoonup s) \\
= & (x \rightharpoonup y) \vee(s \rightharpoonup t) \vee(y \rightharpoonup x) \vee(t \rightharpoonup s) \\
\geqslant & {[x \rightharpoonup(y \vee t)] \vee[s \rightharpoonup(y \vee t)] \vee[y \rightharpoonup(x \vee s)] \vee[t \rightharpoonup(x \vee s)] } \\
& (\text { by Lemma } 2) \\
= & {[(x \vee s) \rightharpoonup(y \vee t)] \vee[(y \vee t) \rightharpoonup(x \vee s)]=d_{1}(x \vee s, y \vee t) . }
\end{aligned}
$$

The rest of (12) and (13) is analogous. Finally, (14)-(17) are consequences of Lemma 7.

Proposition 9. The following holds in any $D R \ell$-monoid:
(1) $|x| \geqslant 0$ with $|x|=0$ iff $x=0$,
(2) $|x|=x$ iff $x \geqslant 0$,
(3) $|x+y| \leqslant|x|+|y|+|x|,|x+y| \leqslant|y|+|x|+|y|$,
(4) $|x \vee y| \leqslant|x| \vee|y|$.

Proof. (1) follows immediately by Proposition 8 (6), (7); (4) is a consequence of Proposition 8 (12).
(2) If $x \geqslant 0$ then $x \geqslant 0 \geqslant 0 \rightharpoonup x$, whence $|x|=x \vee(0 \rightharpoonup x)=x$. Obviously, $x=|x|$ entails $x \geqslant 0$.
(3) Since

$$
\begin{aligned}
d_{1}(x+y, y) & =[(x+y) \rightharpoonup y] \vee[y \rightharpoonup(x+y)] \\
& =[(x+y) \rightharpoonup y] \vee[(y \rightharpoonup y) \rightharpoonup x] \\
& =[(x+y) \rightharpoonup y] \vee(0 \rightharpoonup x) \\
& \leqslant x \vee(0 \rightharpoonup x)=|x|,
\end{aligned}
$$

it follows that

$$
|x+y|=d_{1}(x+y, 0) \leqslant d_{1}(x+y, y)+d_{1}(y, 0)+d_{1}(x+y, y) \leqslant|x|+|y|+|x|
$$

## 3. Ideals

Definition. Let $\mathfrak{A}$ be a $D R \ell$-monoid. A subset $I \subseteq A$ is said to be an ideal of $\mathfrak{A}$ if the following conditions are fulfilled:
(I1) $0 \in I$;
(I2) if $x, y \in I$ then $x+y \in I$;
(I3) if $x \in I, y \in A$ and $|y| \leqslant|x|$ then $y \in I$.
It can be easily seen that the intersection of any family of ideals of $\mathfrak{A}$ is still an ideal. For any $M \subseteq A$, the smallest ideal containing $M$, i.e., the intersection of all ideals $I$ such that $M \subseteq I$, is called the ideal generated by $M$. It will be denoted by $I(M)$.

Proposition 10. Let $\mathfrak{A}$ be a $D R \ell$-monoid. Then for any $\emptyset \neq M \subseteq A$, for any $a \in A$, and for any ideal $J$ we have
(1) $I(M)=\left\{x \in A ;|x| \leqslant\left|a_{1}\right|+\ldots+\left|a_{n}\right|\right.$ for some $\left.a_{1}, \ldots, a_{n} \in M, n \geqslant 1\right\}$;
(2) $I(a)=\{x \in A ;|x| \leqslant n|a|$ for some $n \geqslant 1\}$;
(3) $I(J \cup\{a\})=\left\{x \in A ;|x| \leqslant \sum_{i=1}^{k}\left(a_{i}+n_{i}|a|\right)\right.$, for some $a_{1}, \ldots, a_{k} \in J, n_{1}, \ldots, n_{k} \geqslant$ $0, k \geqslant 1\}$.
Proof. (1) Suppose that $x, y \in I(M)$, i.e., $|x| \leqslant\left|a_{1}\right|+\ldots+\left|a_{n}\right|,|y| \leqslant$ $\left|b_{1}\right|+\ldots+\left|b_{m}\right|$ for some $a_{1}, \ldots, a_{n}, b_{1}, \ldots, b_{m} \in M$ and $n, m \geqslant 1$. Then

$$
\begin{aligned}
|x+y| & \leqslant|x|+|y|+|x| \\
& \leqslant\left|a_{1}\right|+\ldots+\left|a_{n}\right|+\left|b_{1}\right|+\ldots+\left|b_{m}\right|+\left|a_{1}\right|+\ldots+\left|a_{n}\right| .
\end{aligned}
$$

Hence $x+y \in I(M)$. It is easy to see that $|y| \leqslant|x|, x \in I(M)$, implies $y \in I(M)$. Thus $I(M)$ is an ideal. Finally, if $I$ is an ideal such that $M \subseteq I$ then $I(M) \subseteq I$.
(2) and (3) follow by (1); note only that $a_{i} \in J$ iff $\left|a_{i}\right| \in J$ since $J$ is an ideal.

Lemma 11. For each $0 \leqslant x, y, z \in A$, it holds $x \wedge(y+z) \leqslant(x \wedge y)+(x \wedge z)$.
Proof. We compute $(x \wedge y)+(x \wedge z)=(x+x) \wedge(x+z) \wedge(y+x) \wedge(y+z) \geqslant$ $x \wedge x \wedge x \wedge(y+z)=x \wedge(y+z)$.

Proposition 12. If $\mathfrak{A}$ is a $D R \ell$-monoid then for all $x, y \in A$ we have

$$
I(x) \cap I(y)=I(|x| \wedge|y|) \quad \text { and } \quad I(x) \vee I(y)=I(|x| \vee|y|)=I(|x|+|y|)
$$

Proof. Since $\|x\|=|x|$ it is obvious that $I(x)=I(|x|)$. Further, $|x| \wedge|y| \leqslant$ $|x|,|y|$ implies $|x| \wedge|y| \in I(x) \cap I(y)$. Thus $I(|x| \wedge|y|) \subseteq I(x) \cap I(y)$. Conversely, $z \in I(x) \cap I(y)$ iff $|z| \leqslant n|x|$ and $|z| \leqslant m|y|$ for some $n, m \in \mathbb{N}$. Hence $|z| \leqslant$ $n|x| \wedge m|y| \leqslant n m(|x| \wedge|y|)$, by Lemma 11. Therefore, $z \in I(|x| \wedge|y|)$, and so $I(x) \cap I(y) \subseteq I(|x| \wedge|y|)$.

It is easy to see that $I(x) \vee I(y) \subseteq I(|x| \vee|y|) \subseteq I(|x|+|y|)$. Suppose that $J$ is an ideal such that $I(x), I(y) \subseteq J$ and $z \in I(|x|+|y|)$. Then $|z| \leqslant n(|x|+|y|)$ for some $n \in \mathbb{N}$. But $|x|,|y| \in J$, thus $|x|+|y| \in J$ and $z \in J$. This yields $I(x) \vee I(y)=I(|x| \vee|y|)=I(|x|+|y|)$.

Theorem 13. If $\mathfrak{A}$ is a $D R \ell$-monoid then any ideal $I$ is a convex subalgebra in $\mathfrak{A}$. Conversely, if $C$ is a convex subalgebra of $\mathfrak{A}$ such that, for each $x \in A,|x| \in C$ iff $x \in C$, then $C$ is an ideal of $\mathfrak{A}$.

Proof. If $x, y \in I$ then, by Proposition 8,

$$
\left|d_{1}(x, y)\right|=d_{1}(x, y) \leqslant d_{1}(0, y)+d_{1}(x, 0)+d_{1}(0, y)=|y|+|x|+|y| \in I .
$$

Hence $d_{1}(x, y) \in I$. Further,

$$
|x \rightharpoonup y|=(x \rightharpoonup y) \vee(0 \rightharpoonup(x \rightharpoonup y)) \leqslant(x \rightharpoonup y) \vee(y \rightharpoonup x)=d_{1}(x, y) \in I
$$

since $y \rightharpoonup x \geqslant 0 \rightharpoonup(x \rightharpoonup y)$. Thus $x \rightharpoonup y \in I$. Similarly, $d_{2}(x, y) \in I,|x \leftharpoondown y| \leqslant$ $d_{2}(x, y) \in I$ and hence $x \leftharpoondown y \in I$.

To prove that $I$ is a convex subset, suppose $a, b \in I$ and $a \wedge b \leqslant x \leqslant a \vee b$ for some $x \in A$. Then

$$
\begin{gathered}
|x|=x \vee(0 \rightharpoonup x) \leqslant(a \vee b) \vee(0 \rightharpoonup(a \wedge b))=a \vee b \vee(0 \rightharpoonup a) \vee(0 \rightharpoonup b) \\
=a \vee(0 \rightharpoonup a) \vee b \vee(0 \rightharpoonup b)=|a| \vee|b| \leqslant|a|+|b| \in I .
\end{gathered}
$$

Hence $x \in I$.
The proof of the second statement is straightforward.

As argued at the beginning of this section, it is obvious that the set of all ideals of an arbitrary $D R \ell$-monoid, ordered by set inclusion, is a complete lattice.

Theorem 14. For any $D R \ell$-monoid $\mathfrak{A}$, the lattice $\operatorname{Id}(\mathfrak{A})$ of all its ideals is algebraic and Brouwerian.

Proof. It suffices to show that $\operatorname{Id}(\mathfrak{A})$ is distributive and algebraic. (It is wellknown that every algebraic distributive lattice satisfies the join infinite distributive identity and any such a lattice is Brouwerian.)

Let $I, J, K \in \operatorname{Id}(\mathfrak{A})$ and suppose that $x \in I \cap(J \vee K)$. Then $|x| \leqslant a_{1}+\ldots+a_{n}$, for some $0 \leqslant a_{1}, \ldots, a_{n} \in J \cup K$. Hence $|x|=|x| \wedge\left(a_{1}+\ldots+a_{n}\right) \leqslant\left(|x| \wedge a_{1}\right)+\ldots+$ $\left(|x| \wedge a_{n}\right)$. But $|x| \wedge a_{i} \in(I \cap J) \cup(I \cap K) \subseteq(I \cap J) \vee(I \cap K)$, for all $i=1, \ldots, n$, and so $I \cap(J \vee K) \subseteq(I \cap J) \vee(I \cap K)$, proving the distributivity of $\operatorname{Id}(\mathfrak{A})$.

Let $\emptyset \neq M \subseteq A$. For any $x \in A, x \in I(M)$ iff there are $a_{1}, \ldots, a_{n} \in M$ such that $|x| \leqslant\left|a_{1}\right|+\ldots+\left|a_{n}\right|$. Hence $x \in I\left(\left\{a_{1}, \ldots, a_{n}\right\}\right)$ and therefore

$$
I(M)=\bigcup\left\{I(X) ; X \subseteq M,|X|<\aleph_{0}\right\}
$$

Thus $M \mapsto I(M)$ is an algebraic closure operator and, consequently, $\operatorname{Id}(\mathfrak{A})$ is an algebraic lattice.

The following result describes relative pseudocomplements in the lattice $\operatorname{Id}(\mathfrak{A})$.

Proposition 15. For any ideals $J, K$ of $\mathfrak{A}$, the relative pseudocomplement of $J$ with respect to $K$ is given by

$$
J * K=\{x \in A ;|x| \wedge|a| \in K \text { for any } a \in J\} .
$$

Proof. Let us denote by $H$ the set on the right-hand side. First, we will prove that $H$ is an ideal. (I1) $0 \in H$, because $|0| \wedge|a|=0 \in K$ for all $a \in J$. (I2) If $x, y \in H$ then, for each $a \in J$,

$$
|x+y| \wedge|a| \leqslant(|x|+|y|+|x|) \wedge|a| \leqslant(|x| \wedge|a|)+(|y| \wedge|a|)+(|x| \wedge|a|) \in K
$$

so that $x+y \in H$. (I3) If $x \in H$ and $|y| \leqslant|x|$ then $|y| \wedge|a| \leqslant|x| \wedge|a| \in K$, for any $a \in J$, whence $y \in H$.

Now, we have to prove that $H=J * K$. If $x \in J \cap H$ then $|x| \wedge|x| \in K$, thus $x \in K$ and therefore $J \cap H \subseteq K$. In addition, from

$$
J * K=\bigvee\{I \in \operatorname{Id}(\mathfrak{A}) ; I \cap J \subseteq K\}
$$

it follows that $H \subseteq J * K$. Conversely, if $x \in J * K$ then, for each $a \in J,|x| \wedge|a| \in$ $J \cap(J * K) \subseteq K$ since $|x| \wedge|a| \leqslant|a| \in J$ and $|x| \wedge|a| \leqslant|x| \in J * K$. Hence $x \in H$. So $H=J * K$.

The pseudocomplement of an ideal $I$ is $I^{*}:=I *\{0\}$.

Corollary 16. $I^{*}=\{x \in A ;|x| \wedge|a|=0$ for each $a \in I\}$.
Let $\mathfrak{A}$ be a $D R \ell$-monoid and $I \in \operatorname{Id}(\mathfrak{A})$. Let us define two binary relations on $A$ by

$$
\begin{aligned}
& \langle x, y\rangle \in \Theta_{1}(I) \Longleftrightarrow d_{1}(x, y) \in I, \\
& \langle x, y\rangle \in \Theta_{2}(I) \Longleftrightarrow d_{2}(x, y) \in I,
\end{aligned}
$$

for each $x, y \in A$.

Lemma 17. For any ideal $I, \Theta_{1}(I)$ and $\Theta_{2}(I)$ are equivalence relations.
Proof. It is obvious that $\Theta_{1}(I)$ is reflexive and symmetric. The transitivity follows from Proposition 8. Indeed, if $\langle x, y\rangle,\langle y, z\rangle \in \Theta_{1}(I)$ then $d_{1}(x, z) \leqslant d_{1}(x, y)+$ $d_{1}(y, z)+d_{1}(x, y) \in I$, hence $d_{1}(x, z) \in I$. Similarly for $\Theta_{2}(I)$.

Theorem 18. For any ideal $I$ of $\mathfrak{A}$, the relations $\Theta_{1}(I)$ and $\Theta_{2}(I)$ are congruence relations on the lattice $\mathfrak{L}(\mathfrak{A})$. Moreover, $I=[0]_{\Theta_{1}(I)}=[0]_{\Theta_{2}(I)}$.

Proof. Let $\langle x, y\rangle,\langle s, t\rangle \in \Theta_{1}(I)$, i.e., $d_{1}(x, y), d_{1}(s, t) \in I$. Then, by Proposition 8 ,

$$
\begin{aligned}
& d_{1}(x \vee s, y \vee t) \leqslant d_{1}(x, y) \vee d_{1}(s, t) \leqslant d_{1}(x, y)+d_{1}(s, t) \in I, \\
& d_{1}(x \wedge s, y \wedge t) \leqslant d_{1}(x, y) \vee d_{1}(s, t) \leqslant d_{1}(x, y)+d_{1}(s, t) \in I
\end{aligned}
$$

Hence $\langle x \vee s, y \vee t\rangle,\langle x \wedge s, y \wedge t\rangle \in \Theta_{1}(I)$. Similarly for $\Theta_{2}(I)$.
For each $x \in A, x \in[0]_{\Theta_{1}(I)}$ iff $\langle x, 0\rangle \in \Theta_{1}(I)$ iff $d_{1}(x, 0)=|x| \in I$ iff $x \in I$.

Theorem 19. Let $I$ be an ideal of a $D R \ell$-monoid $\mathfrak{A}$. Then $\mathfrak{L}(\mathfrak{A}) / \Theta_{1}(I)$ is a distributive lattice whose partial order relation is defined by

$$
[x]_{\Theta_{1}(I)} \leqslant[y]_{\Theta_{1}(I)} \Longleftrightarrow(x \rightharpoonup y)^{+} \in I .
$$

Similarly, $\mathfrak{L}(\mathfrak{A}) / \Theta_{2}(I)$ is a distributive lattice in which

$$
[x]_{\Theta_{2}(I)} \leqslant[y]_{\Theta_{2}(I)} \Longleftrightarrow(x \leftharpoondown y)^{+} \in I .
$$

Proof. Since $\mathfrak{L}(\mathfrak{A})$ is a distributive lattice, by Corollary 3, so is $\mathfrak{L}(\mathfrak{A}) / \Theta_{1}(I)$. Further, for each $x, y \in A,[x]_{\Theta_{1}(I)} \leqslant[y]_{\Theta_{1}(I)}$ iff $[x]_{\Theta_{1}(I)} \vee[y]_{\Theta_{1}(I)}=[x \vee y]_{\Theta_{1}(I)}=$
$[y]_{\Theta_{1}(I)}$ iff $\langle x \vee y, y\rangle \in \Theta_{1}(I)$ iff $d_{1}(x \vee y, y) \in I$ iff $(x \rightharpoonup y)^{+} \in I$. Indeed, since

$$
\begin{aligned}
d_{1}(x \vee y, y) & =[((x \vee y) \rightharpoonup y) \vee 0]+[(y \rightharpoonup(x \vee y)) \vee 0] \\
& =[(x \rightharpoonup y) \vee(y \rightharpoonup y) \vee 0]+0 \\
& =(x \rightharpoonup y) \vee 0=(x \rightharpoonup y)^{+} .
\end{aligned}
$$

The proof of the other statement is analogous.

## 4. Normal ideals

Definition. An ideal $I$ is said to be normal if it satisfies the following condition, for each $x, y \in A$ :

$$
(x \rightharpoonup y)^{+} \in I \Longleftrightarrow(x \leftharpoondown y)^{+} \in I .
$$

The set of all normal ideals of a $D R \ell$-monoid $\mathfrak{A}$ will be denoted by $N(\mathfrak{A})$. For an ideal $I$, we denote $I^{+}=\{x \in I ; x \geqslant 0\}$.

Proposition 20. For any $I \in \operatorname{Id}(\mathfrak{A})$, the following conditions are equivalent:
(i) $I \in \mathrm{~N}(\mathfrak{A})$;
(ii) $x+I^{+}=I^{+}+x$, for any $x \in A$.

Proof. (i) $\Rightarrow$ (ii) Let $x \in A, h \in I^{+}$and set $y=h+x \in I^{+}+x$. It is clear that $y \geqslant x$ and, consequently, $y=x \vee y=(y \rightharpoonup x)^{+}+x=x+(y \leftharpoondown x)^{+}$. From $h+x \geqslant y$ it follows that $h \geqslant y \rightharpoonup x \geqslant 0$, since $y \geqslant x$. Hence $(y \rightharpoonup x)^{+}=y \rightharpoonup x \in I^{+}$. But $I \in \mathrm{~N}(\mathfrak{A})$, so that $(y \leftharpoondown x)^{+} \in I^{+}$. Thus $y \in x+I^{+}$. Similarly, $x+I^{+} \subseteq I^{+}+x$.
(ii) $\Rightarrow$ (i) If $(y \rightharpoonup x)^{+} \in I$ then $x \vee y=(y \rightharpoonup x)^{+}+x=x+h$ for some $h \in I^{+}$. Therefore $y \leqslant x+h$, which yields $(y \leftharpoondown x)^{+} \leqslant((x+h) \leftharpoondown x)^{+} \leqslant h \vee 0=h \in I^{+}$. Thus $(y \leftharpoondown x)^{+} \in I$. The converse is analogous.

Lemma 21. If $J$ and $K$ are normal ideals of a $D R \ell$-monoid $\mathfrak{A}$ then

$$
J \vee K=\left\{x \in A ;|x| \leqslant a+b \text { for some } a \in J^{+}, b \in K^{+}\right\} .
$$

In addition, $J \vee K$ is a normal ideal of $\mathfrak{A}$.
Proof. Let us denote the set on the right-hand side by $M$. (I1) and (I3) are obviously satisfied. To prove (I2), let $x, y \in M$, i.e., $|x| \leqslant a+b$ and $|y| \leqslant c+d$ for some $a, c \in J^{+}$and $b, d \in K^{+}$. Then $|x+y| \leqslant|x|+|y|+|x| \leqslant a+b+c+d+a+b=a^{\prime}+b^{\prime}$ for some $a^{\prime} \in J^{+}, b^{\prime} \in K^{+}$, by Proposition 20. Consequently, $M \in \operatorname{Id}(\mathfrak{A})$. Finally, it is easy to see that any ideal $H$ such that $J, K \subseteq H$ contains $M$.

If $(x \rightharpoonup y)^{+} \in J \vee K$ then $(x \rightharpoonup y)^{+} \leqslant a+b$ for some $a \in J^{+}, b \in K^{+}$. Hence $a+b \geqslant x \rightharpoonup y$ iff $a+b+y \geqslant x$. Since $J, K \in \mathrm{~N}(\mathfrak{A})$, there exist $a^{\prime} \in J^{+}$and $b^{\prime} \in K^{+}$ such that $a+b+y=y+a^{\prime}+b^{\prime}$. Therefore $y+a^{\prime}+b^{\prime} \geqslant x$ iff $a^{\prime}+b^{\prime} \geqslant x \leftharpoondown y$, whence $a^{\prime}+b^{\prime} \geqslant(x \leftharpoondown y)^{+}$and $(x \leftharpoondown y)^{+} \in J \vee K$, proving that $J \vee K$ is a normal ideal.

Proposition 22. Let $\mathfrak{A}$ be a $D R \ell$-monoid. Then $N(\mathfrak{A})$ is a complete sublattice of $\operatorname{Id}(\mathfrak{A})$.

Proof. Let $\left\{I_{\lambda}\right\}_{\lambda \in \Lambda}$ be a family of normal ideals. Obviously, $\bigcap_{\lambda \in \Lambda} I_{\lambda}$ is a normal ideal. Let us assume $(x \rightharpoonup y)^{+} \in \bigvee_{\lambda \in \Lambda} I_{\lambda}$, for some $x, y \in A$; then $(x \rightharpoonup y)^{+} \in \bigvee_{\lambda \in \Lambda_{0}} I_{\lambda}$, for some finite subset $\Lambda_{0}$ of $\Lambda$. Hence $(x \leftharpoondown y)^{+} \in \underset{\lambda \in \Lambda_{0}}{ } I_{\lambda}$ since it is a normal ideal. Thus $(x \leftharpoondown y)^{+} \in \underset{\lambda \in \Lambda}{\bigvee} I_{\lambda}$. The converse is analogous.

Proposition 23. Let $\mathfrak{A}$ and $\mathfrak{B}$ be $D R \ell$-monoids and $\varphi: A \rightarrow B$ a homomorphism. Then $\varphi^{-1}(0)=\{x \in A ; \varphi(x)=0\}$ is a normal ideal of $\mathfrak{A}$.

Proof. Clearly, the conditions (I1) and (I2) hold. Suppose $\varphi(x)=0$ and $|y| \leqslant|x|$. Then $\varphi(|x|)=\varphi(x \vee(0 \rightharpoonup x))=\varphi(x) \vee(0 \rightharpoonup \varphi(x))=0$ and, consequently, $\varphi(|y|)=0$. Hence $\varphi(y \vee(0 \rightharpoonup y))=\varphi(y) \vee(0 \rightharpoonup \varphi(y))=0$, which gives $\varphi(y)=0$. Thus, $\varphi^{-1}(0)$ is an ideal in $\mathfrak{A}$.

Finally, $(x \rightharpoonup y)^{+} \in \varphi^{-1}(0)$ iff $\varphi((x \rightharpoonup y) \vee 0)=(\varphi(x) \rightharpoonup \varphi(y)) \vee 0=0$. Hence $0 \geqslant \varphi(x) \rightharpoonup \varphi(y)$ iff $\varphi(y) \geqslant \varphi(x)$ iff $0 \geqslant \varphi(x) \leftharpoondown \varphi(y)$. Therefore $0=(\varphi(x) \leftharpoondown$ $\varphi(y)) \vee 0=\varphi((x \leftharpoondown y) \vee 0)$, thus $(x \leftharpoondown y)^{+} \in \varphi^{-1}(0)$.

Proposition 24. If $I \in \mathrm{~N}(\mathfrak{A})$ then, for all $x, y \in A, d_{1}(x, y) \in I$ iff $d_{2}(x, y) \in I$.
Proof. If $d_{1}(x, y)=(x \rightharpoonup y)^{+}+(y \rightharpoonup x)^{+} \in I$ then $(x \rightharpoonup y)^{+},(y \rightharpoonup x)^{+} \in I$. Since $I$ is a normal ideal, this implies $(x \leftharpoondown y)^{+},(y \leftharpoondown x)^{+} \in I$. Hence $d_{2}(x, y)=$ $(x \leftharpoondown y)^{+}+(y \leftharpoondown x)^{+} \in I$.

Corollary 25. If $I$ is a normal ideal then $\Theta_{1}(I)=\Theta_{2}(I)$; it will be denoted by $\Theta(I)$.

Lemma 26. Let $I \in \mathrm{~N}(\mathfrak{A})$. If $\langle x, y\rangle \in \Theta(I)$ then, for each $z \in A$,

$$
\begin{array}{ll}
\langle x \rightharpoonup z, y \rightharpoonup z\rangle \in \Theta(I), & \langle x \leftharpoondown z, y \leftharpoondown z\rangle \in \Theta(I), \\
\langle z \rightharpoonup x, z \rightharpoonup y\rangle \in \Theta(I), & \langle z \leftharpoondown x, z \leftharpoondown y\rangle \in \Theta(I) .
\end{array}
$$

Proof. This follows from Proposition 8 (14)-(17).

Theorem 27. If I is a normal ideal of a $D R \ell$-monoid $\mathfrak{A}$ then $\Theta(I)$ is a congruence relation on $\mathfrak{A}$. In addition, $[0]_{\Theta}=I$.

Proof. Let $\langle x, y\rangle \in \Theta(I)$ and $\langle s, t\rangle \in \Theta(I)$. Then $(x \rightharpoonup y)^{+},(s \rightharpoonup t)^{+} \in I$. Obviously, $x \leqslant x \vee y=(x \rightharpoonup y)^{+}+y$ and $s \leqslant s \vee t=(s \rightharpoonup t)^{+}+t$. Hence, it follows that

$$
\begin{aligned}
x+s & \leqslant(x \rightharpoonup y)^{+}+y+(s \rightharpoonup t)^{+}+t \\
& =(x \rightharpoonup y)^{+}+\left(y+(s \rightharpoonup t)^{+}\right)+t \\
& =(x \rightharpoonup y)^{+}+(h+y)+t \\
& =\left((x \rightharpoonup y)^{+}+h\right)+(y+t)
\end{aligned}
$$

for some $h \in I^{+}$such that $y+(s \rightharpoonup t)^{+}=h+y$. However, $\left((x \rightharpoonup y)^{+}+h\right)+(y+t) \geqslant$ $x+s$ iff $(x \rightharpoonup y)^{+}+h \geqslant(x+s) \rightharpoonup(y+t)$. Therefore, $((x+s) \rightharpoonup(y+t))^{+} \leqslant((x \rightharpoonup$ $\left.y)^{+}+h\right)^{+}=(x \rightharpoonup y)^{+}+h \in I$. So $((x+s) \rightharpoonup(y+t))^{+} \in I$. We can similarly show that $((y+t) \rightharpoonup(x+s))^{+} \in I$. Hence, we conclude that $d_{1}(x+s, y+t)=((x+s) \rightharpoonup$ $(y+t))^{+}+((y+t) \rightharpoonup(x+s))^{+} \in I$ and $\langle x+s, y+t\rangle \in \Theta(I)$.

By Lemma 26, $\langle x \rightharpoonup s, y \rightharpoonup s\rangle \in \Theta(I)$ and $\langle y \rightharpoonup s, y \rightharpoonup t\rangle \in \Theta(I)$. This yields $\langle x \rightharpoonup s, y \rightharpoonup t\rangle \in \Theta(I)$. Similarly, $\langle x \leftharpoondown s, y \leftharpoondown t\rangle \in \Theta(I)$.

The rest follows by Theorem 18.
Theorem 28. If $\Theta$ is a congruence on $\mathfrak{A}$ then $[0]_{\Theta}=\{x \in A ;\langle x, 0\rangle \in \Theta\}$ is a normal ideal in $\mathfrak{A}$. Moreover, $\Theta=\Theta\left([0]_{\Theta}\right)$.

Proof. The first part follows by Proposition 23. Further, we claim that

$$
\begin{equation*}
\langle x, y\rangle \in \Theta \Longleftrightarrow\left\langle d_{1}(x, y), 0\right\rangle \in \Theta \tag{C}
\end{equation*}
$$

or equivalently,

$$
\langle x, y\rangle \in \Theta \Longleftrightarrow\left\langle d_{2}(x, y), 0\right\rangle \in \Theta
$$

Indeed, if $\langle x, y\rangle \in \Theta$ then $\langle x \rightharpoonup y, 0\rangle \in \Theta$ and $\langle y \rightharpoonup x, 0\rangle \in \Theta$, whence $\left\langle d_{1}(x, y), 0\right\rangle=\langle(x \rightharpoonup y) \vee(y \rightharpoonup x), 0\rangle \in \Theta$. Conversely, $\left\langle d_{1}(x, y), 0\right\rangle \in \Theta$ iff $d_{1}(x, y) \in[0]_{\Theta}$ which implies $(x \rightharpoonup y)^{+},(y \rightharpoonup x)^{+} \in[0]_{\Theta}$. This gives

$$
\begin{align*}
& x \vee y=(x \rightharpoonup y)^{+}+y \equiv 0+y=y \\
& x \vee y=(y \rightharpoonup x)^{+}+x \equiv 0+x=x
\end{align*}
$$

Thus, by the transitivity, $\langle x, y\rangle \in \Theta$.
Now, $\Theta=\Theta\left([0]_{\Theta}\right)$ is an immediate consequence of (C).

Corollary 29. In any $D R \ell$-monoid, there is a one-to-one correspondence between congruences and normal ideals.

## 5. Deductive systems

It was proved in [8] that the variety of $D R \ell$-monoids is weakly regular, that is, $[0]_{\Phi}=[0]_{\Psi}$ entails $\Phi=\Psi$, for any congruences $\Phi, \Psi$ on an arbitrary $D R \ell$-monoid. Hence it follows that congruence kernels of $D R \ell$-monoids can also be described by means of so-called deductive systems (see [6]).

Definition. Let $\mathfrak{A}$ be a $D R \ell$-monoid and $D \subseteq A$. Then $D$ is said to be a deductive system if the following conditions are fulfilled:
(D1) $0 \in D$;
(D2) if $x \in D$ and $d_{1}(x, y) \in D$ then $y \in D$;
(D3) if $x \in D$ then $d_{1}(x, 0) \in D$.
A deductive system $D$ is called compatible iff the following holds:
If $d_{1}(x, y) \in D$ and $d_{1}(s, t) \in D$, for $x, y, s, t \in A$, then $d_{1}(f(x, s), f(y, t)) \in D$, for each $f \in\{+, \vee, \wedge, \rightharpoonup, \leftharpoondown\}$.

The following result is only a special case of [ 6 , Theorems 1, 2] and it generalizes the analogous property of $G M V$-algebras ([7, Theorems 2.8, 2.9]).

Theorem 30. Let $\mathfrak{A}$ be a $D R \ell$-monoid, $D \subseteq A$. Let us define a binary relation $\Theta_{D}$ via

$$
\langle x, y\rangle \in \Theta_{D} \Longleftrightarrow d_{1}(x, y) \in D,
$$

for every $x, y \in A$. If $D$ is a compatible deductive system then $\Theta_{D}$ is a congruence on $\mathfrak{A}$ such that $[0]_{\Theta_{D}}=D$. Conversely, if $\Theta$ is a congruence relation on $\mathfrak{A}$ then $[0]_{\Theta}$ is a compatible deductive system and $\Theta_{[0]_{\Theta}}=\Theta$.

Therefore by Theorems 27 and 28 we immediately obtain

Corollary 31. If $\mathfrak{A}$ is a $D R \ell$-monoid and $D \subseteq A$ then the following conditions are equivalent:
(i) $D$ is a normal ideal;
(ii) $D$ is a compatible deductive system;
(iii) $D=[0]_{\Theta}$ for some congruence relation $\Theta$ on $\mathfrak{A}$.

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