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# ORTHOGONAL PARTITIONS AND COVERING OF GRAPHS

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### INTRODUCTION

We investigate two notions related to the notion of chromatic number. For a given graph G one may define  $\overline{\omega}(G)$  as the minimum number of stable sets (i.e. subsets of the vertex set with no two vertices joined by an edge) such that every pair of non-adjacent vertices is contained in at least one of them, and dim G as the minimal d such that G is contained in a product of d complete graphs as an induced subgraph.

Determination of both  $\overline{\omega}$  and dim is an NP-complete problem. (See [15], [10].) In [8] and [12], dim of some special graphs was investigated. The main purpose of this paper is to prove some analogous statements concerning the notion of  $\overline{\omega}$  and extend some results of [8] and [12].

The first two sections are of technical character. In § 1 we present some remarks concerning the notion of orthogonal array (which is in our sense more general than that defined in  $\lceil 5 \rceil$ ).

In §2 we introduce the notion of orthogonal q-partitions and investigate the problem of maximal number of pairwise orthogonal q-partitions on a given set. For q = 2 we are able to give an exact formula for this number. For  $q \ge 3$  we give only some bounds.

In §4 we show that  $\overline{\omega}$  of a product of two graphs with large  $\overline{\omega}$  is large, too. To prove an analogous statement concerning chromatic number is an open (and probably difficult) problem.

In last paragraphs using the results of the first two sections we give a formula for  $\overline{\omega}(nK_2)$  and some bounds for  $\overline{\omega}(nK_q)$ , dim  $(nK_q)$  and dim  $\binom{n}{q}$ .

We state now three well known theorems which will be used throughout our paper.

**Theorem 0.1.** (see [5]) Let q be a power of a prime. Then there exist q - 1 pairwise orthogonal latin squares of order q.

**Theorem 0.2.** (see [2]) Let  $A_1, A_2, ..., A_n$  be a system of subsets of 1, 2, ..., t such that:

(i)  $A_i \notin A_j;$ (ii)  $A_i \cap A_j \neq \emptyset;$ (iii)  $|A_i| \leq \frac{1}{2}t.$ 

Then

$$n \leq \left( \frac{t-1}{\left\lfloor \frac{t}{2} \right\rfloor - 1} \right).$$

Theorem 0.3. For every integer n there exists a prime p such that

 $n \leq p \leq 2n$ .

#### 1. ORTHOGONAL ARRAYS

Let q be a positive integer. We say that two vectors  $\mathbf{u} = (u_1, u_2, ..., u_{t_1})$  and  $\mathbf{v} = (v_1, v_2, ..., v_{t_2})$  are orthogonal if for every k, k' = 1, 2, ..., q there exists j = 1, 2, ..., t so that  $u_j = k$  and  $v_j = k'$ .

For two vectors  $\mathbf{u} = (u_1, u_2, ..., u_t)$  and  $\mathbf{v} = (v_1, v_2, ..., v_t)$  we put  $(\mathbf{u}, \mathbf{v}) = (u_1, u_2, ..., u_{t_1}, v_1, v_2, ..., v_{t_2})$ .

Let  $\pi$  be a permutation of the set  $\{1, 2, ..., q\}$ . We identify  $\pi$  with the vector  $(\pi(1), \pi(2), ..., \pi(q))$ . If  $\mathbf{u} = (u_i), u_i \in \{1, 2, ..., q\}$  is another vector then we put

$$\pi(\mathbf{u}) = (\pi(u_1), \pi(u_2), ..., \pi(u_t)).$$

Obviously the following holds

$$\pi(\boldsymbol{u},\,\boldsymbol{v}) = (\pi(\boldsymbol{u}),\,\pi(\boldsymbol{v}))$$

We say that a vector w has the permutation property if

$$\mathbf{w} = (\pi_1, \pi_2, \dots, \pi_r)$$

where  $\pi_s$ , s = 1, 2, ..., r, are permutations of  $\{1, 2, ..., q\}$  for some q and r.

The orthogonal array OA(q, n, t) is an  $n \times t$  matrix  $(a_{ij})$  where  $a_{ij} \in \{1, 2, ..., q\}$  for every i, j, i = 1, 2, ..., n, j = 1, 2, ..., t, such that every two row-vectors are orthogonal. The orthogonal array (OA) has the permutation property if each row of the corresponding matrix has this property.

Let  $t = q^2$ . Then the existence of **OA** $(q, n, q^2)$  is equivalent with the existence of n - 2 pairwise orthogonal latin squares of order q. (See e.g. [1, 5].) In [3], the maximal number of pairwise orthogonal latin squares is estimated. For a given q there exist at most q - 1 pairwise orthogonal latin squares, and if q is a power of a prime then such a system exists.

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**Theorem 1.1.** Let q be a power of a prime. Then

(i) There exists  $OA(q, q + 1, q^2)$ .

(ii) There exists  $OA(q, q, q^2)$  which has the permutation property.

Proof. By Theorem 0.1 and the relation between **OA**'s we have (i). It remains to prove (ii).

Take  $OA(q, q + 1, q^2)$ . In each row there are exactly q entries equal to k for each k = 1, 2, ..., q. We can suppose that the last row is of the form (1, 1, ..., 1, 2, 2, ..., ..., 2, ..., q, q, ..., q). (This follows from the obvious fact that if we rearrange columns of OA we obtain again OA.) The first q rows obviously form an OA with the permutation property. So for i = 1, 2, ..., q the *i*-th row is of the form  $(\pi_{i1}, \pi_{i2}, ..., \pi_{iq})$ .

The following theorem as well as its proof is a simple generalization of [9] (see [5], Theorem 13.2.1).

**Theorem 1.2.** The existence of  $OA(q, n_1, t_1)$  and  $OA(q, n_2, t_2)$  implies the existence of  $OA(q_1q_2, n, t_1t_2)$ .

**Theorem 1.3.** The existence of  $OA(q, n_1, t_1)$  and  $OA(q, n_2, t_2)$  implies the existence of  $OA(q, n_1n_2, t_1 + t_2)$ .

Moreover, if the first two **OA**'s have the permutation property then the last one has the property as well.

Proof. Denote the row-vectors of  $OA(q, n_1, t_1)$  and  $OA(q, n_2, t_2)$  by  $u_1, u_2, ..., u_{t_1}$  and  $v_1, v_2, v_{..., t_2}$ , respectively. Then  $w_{i,i'} = (u_i, v_{i'})$ ,  $i = 1, 2, ..., t_1$ ,  $i' = 1, 2, ..., t_2$  are pairwise orthogonal vectors and hence  $OA(q, n_1n_2, t_1 + t_2)$  exists. Moreover, if  $u_i, v_{i'}$  have the permutation property, then  $w_{i,i'}$  has the permutation property, too.

The following theorem gives a better estimate for the case that we start the construction with q **OA**'s.

**Theorem 1.4.** Let q be a power of a prime. Let  $OA(q, n_j, t_j)$ , j = 1, 2, ..., q, exist. Then there exists  $OA(q, qn_1n_2 ... n_q, t_1 + ... + t_q)$ . Moreover, if the first q OA's have the permutation property then the last one has the property as well.

Proof. Denote by  $\mathbf{v}_{i_j}^j$  the  $i_j$ -th row-vector of  $OA(q, n_j, t_j)$ . Take  $OA(q, q, q^2)$  with the permutation property (the existence follows from Theorem 1.1) and denote its *i*-th row by  $\pi_i = (\pi_{i1}, \pi_{i2}, ..., \pi_{iq})$ . The vectors

$$\mathbf{v}_{i_0,i_1,i_2,...,i_q} = (\pi_{i_01}(\mathbf{u}_{i_1}^1), \pi_{i_02}(\mathbf{u}_{i_2}^2), ..., \pi_{i_0q}(\mathbf{u}_{i_q}^q))$$

for  $i_0 \in \{1, 2, ..., q\}$  and  $i_j \in \{1, 2, ..., n_j\}$ , for each  $j \in \{1, 2, ..., q\}$ , form **OA** $(q, qn_1n_2 ... n_q, t_1 + t_2 + ... + t_q)$ . If the vectors  $\mathbf{u}_{i_j}^j$ ,  $j \in \{1, 2, ..., q\}$ , have the permutation property then the vectors  $\mathbf{v}_{i_0,i_1,...,i_q}$  have the property as well.

**Definition.** Let  $\mathscr{A} = (A_1, A_2, ..., A_q), \mathscr{B} = (B_1, B_2, ..., B_q)$  be two partitions into q parts (q-partitions) of a set with t elements. We say that  $\mathscr{A}$  and  $\mathscr{B}$  are orthogonal if  $A_i \cap B_{i'} \neq 0$  for every  $i, i' \in \{1, 2, ..., q\}$ .

We denote by f(t, q) the maximal size of a system of pairwise orthogonal q-partitions of the t-element set.

Theorem 2.1.

$$f(t,2) = \left( \begin{bmatrix} t - 1 \\ \lfloor \frac{t}{2} \end{bmatrix} - 1 \right).$$

Proof. Let  $\mathscr{A}_1, \mathscr{A}_2, ..., \mathscr{A}_n$  be a system of pairwise orthogonal 2-partitions. For every i = 1, ..., n choose the partition class  $A_i$  with cardinality  $\leq t/2$ . The sets  $A_1, A_2, ..., A_n$  have the following properties:

- (i) they form an antichain in the partial order of subsets of  $\{1, ..., t\}$ ;
- (ii)  $A_i \cap A_j \neq 0$  for every  $i \neq j$ ;
- (iii)  $|A_i| \leq \frac{t}{2}$  for every *i*.

Hence, by Erdös Ko Rado theorem [2], we get

$$f(t,2) \leq \begin{pmatrix} t-1\\ \begin{bmatrix} t\\ 2 \end{bmatrix} -1 \end{pmatrix}.$$

To prove  $\geq$  consider the 2-partitions of  $\{1, ..., t\}$  defined as follows: To every subset with cardinality  $\left[\frac{t}{2}\right]$  containing 1, assign a 2-partition consisting of this set and its complement.

For q > 2 we are far from being able to find an exact formula for f(t, q). We give here some estimates.

Theorem 2.2.

$$f\left(\begin{pmatrix} q\\2 \end{pmatrix}t, q\right) \ge f(t, 2).$$

Proof. Suppose that on the set  $X = \{1, 2, ..., t\}$  there exists a system of *n* pairwise orthogonal 2-partitions. Consider  $\binom{q}{2}$  copies  $X^{kj}$   $(1 \le k < j \le q)$  of the set X and on each set  $X^{kj}$  a system of pairwise orthogonal 2-partitions  $\mathscr{A}_i^{kj} = (\mathcal{A}_i^{kj}, \mathcal{B}_i^{kj}), i =$ = 1, ..., n. For every *i* define a *q*-partition  $\mathscr{C}_i = (C_{i1}, C_{i2}, ..., C_{iq})$  of the set Y =  $= \bigcup_{k \le j} X^{kj}$  by formula

$$C_{ik} = \bigcup_{j=1}^{k=1} A_i^{jk} \cup \bigcup_{j=k+1}^{q} B_i^{kj}, \quad k = 1, ..., q.$$

It is easy to verify that  $\mathscr{C}_i$ , i = 1, ..., n, is a system of pairwise orthogonal q-partitions.

**Theorem 2.3.** OA(q, n, t) exists if and only if  $f(t, q) \ge n$ .

Proof. The proof follows from the fact that each q-partition  $\mathscr{A}$  of the set  $\{1, 2, ..., t\}$  is in a 1-1 correspondence with the vector  $\mathbf{u} = (u_1, u_2, ..., u_t)$  defined as follows:  $u_j = k$  iff  $j \in A_k$  where  $A_k$  is the k-th class of  $\mathscr{A}$ .

Theorems 1.1-1.4 and Theorem 2.3 imply immediately

**Theorem 2.4.** Let  $t, t_1, t_2, n, n_1, n_2, q, q_1, q_2$  be positive integers. Then

(i)  $f(t_1, q_1) \ge n, f(t_2, q_2) \ge n$  implies  $f(t_1t_2, q_1q_2) \ge n$ ;

(ii)  $f(t_1, q) \ge n_1, f(t_2, q) \ge n_2$  implies  $f(t_1 + t_2, q) \ge n_1 n_2$ .

**Theorem 2.5.** Let q be a power of a prime. Then

(i)  $f(q^2, q) = q + 1;$ 

(ii)  $f(t_j, q) \ge n_j$  for j = 1, ..., q implies  $f(t_1 + ... + t_q, q) \ge qn_1n_2 ... n_q$ .

An easy computation gives the following

Corollary 2.6. For q a power of a prime and t of the form

$$t = \alpha_j q^j + \alpha_{j-1} q^{j-1} + \ldots + \alpha_1 q + \alpha_0$$

it is

$$f(t, q) \ge \prod_{i=2}^{j} \left[ (q + 1)^{q^{i-2}} q^{(q^{i-1}-1)/(q-1)} \right]^{\alpha_i}$$

and hence e.g.  $f(t, q) \ge (q + 1)^{[t/q^2]}$ .

**Remark 2.7.** For q = 3,4 Theorem 2.2 gives a better result than Corollary 2.6. Nevertheles, we have not been able to generalize the construction given in 2.2 to improve the estimates given in 2.6 for q > 3.

**Theorem 2.8.** Let t, p, q be positive integers,  $p \leq q$ . Then

$$f(t, q) \leq f(t, p)/f(q, p).$$

Particularly,  $f(t, q) \leq f(t, q - 1)$ .

Proof. Let  $\mathfrak{S}$  and  $\mathfrak{T}$  be systems of orthogonal q-partitions and p-partitions on the t-set and q-set, respectively. To every pair  $\mathscr{A}, \mathscr{B}, \mathscr{A} = (A_1, ..., A_q), \mathscr{B} = (B_1, ..., B_p)$ ,

 $\mathscr{A} \in \mathfrak{S}, \mathscr{B} \in \mathfrak{T}$ , assign  $\mathscr{C}(\mathscr{A}, \mathscr{B}) = (C_1, ..., C_p)$  defined as follows:

$$C_i = \bigcup \{A_j | j \in B_i\}, \quad i = 1, ..., p.$$

 $\{\mathscr{C}(\mathscr{A}, \mathscr{B}) | \mathscr{A} \in \mathfrak{S}, \mathscr{B} \in \mathfrak{T}\}$  obviously forms a system of pairwise orthogonal *p*-partitions.

Corollary 2.9.

$$f(t,q) \ge (2q)^{t/4q^2}$$

for every integers  $t, q, t \ge 8q^3$ .

Proof. Immediately follows from 2.6 and 0.3.

**Proposition 2.10.** Let t > q be positive integers. Then

$$f(t,q) \leq \left( \begin{bmatrix} t-1\\ \lfloor \frac{t}{q} - 1 \end{bmatrix} \right).$$

Proof. Let  $\mathfrak{S}$  be a maximal system of pairwise orthogonal q-partitions of a set  $\{1, ..., t\}$ . From every  $\mathscr{A}_j \in \mathfrak{S}, j = 1, ..., f(t, q)$  choose one of its partition classes  $A^j \in \mathscr{A}_j$  with  $|A^j| \leq t/q$ . The system  $A^j$  has the following properties:

 $|A^j| \leq t/q$ ,  $A^i \notin A^j$  and  $A^i \cap A^j \neq 0$  for every  $i \neq j$ .

Thus the Erdös Ko Rado theorem yields

$$|\mathfrak{S}| \leq \begin{pmatrix} t-1\\ \left\lfloor \frac{t}{q} - 1 \right\rfloor \end{pmatrix}.$$

3. 
$$\overline{\omega}(nK_a)$$

We shall now give estimates of the number  $\overline{\omega}$ , defined in the introduction, for the graphs  $nK_{q'}$  (The graph  $nK_q$  consists of *n* disjoint copies of the complete graph  $K_{q'}$ ) The notion of  $\overline{\omega}$  was in fact studied in a "complementary form" in several papers, e.g. [4], [15]. The number  $\omega(G)$  is defined in [6] as the minimal size of a set S such that G is isomorphic to an intersection graph of a system of subsets of the set S.

**Proposition 3.1.**  $\overline{\omega}(G) = \omega(\overline{G})$  where  $\overline{G}$  denotes the complement of G. Hence  $\overline{\omega}(G)$  equals to the minimal size of a set S such that to every vertex of G one can assign a subset of S so that two vertices are adjacent iff the corresponding subsets are disjoint. (This assignment is called a disjoint representation of G.)

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The proof immediately follows from Proposition 1, Chapter 17, 4 in 0.

# Theorem 3.2.

- (i)  $\overline{\omega}(nK_q) \leq t$  iff OA(q, n, t) exists;
- (ii)  $\overline{\omega}(nK_q) = \text{Min}\{t \mid \text{there exists at least } n \text{ orthogonal } q\text{-partitions of the set} \{1, 2, ..., t\}\}.$

Proof. We shall write  $nK_q = K_q^1 + \ldots + K_q^n$  where  $V(K_q^i) = \{x_1^i, \ldots, x_q^i\}$ . Let  $A_1, \ldots, A_t$  be a system of stable sets so that every pair of nonadjacent vertices of  $nK_q$  is contained in at least one of  $A_i$ 's. Without loss of generality we may suppose  $A_j$  to be maximal, i.e.  $|A_j| = n$  for each j. To every  $j \in \{1, \ldots, t\}$  assign a column vector  $(\mathbf{u}_{ij})_{1 \le i \le n}$  so that

$$(*) u_{ij} = k \quad \text{iff} \quad x_k^i \in A_j \,.$$

The above defined system of column vectors forms an OA(q, n, t). On the other hand, if an OA(q, n, t) exists then the formula (\*) defines a system of  $A_j$ 's with the required properties.

(ii) follows from (i) and Theorem 2.3.

The proof of the following theorem is easy and therefore it is omitted.

**Theorem 3.3.** Let  $A_1, \ldots, A_{q,r}$  be a system of stable sets of  $nK_q$  such that every pair of nonadjacent vertices is contained in at least one of them. Then the sets  $A_{(s-1)q+1}, A_{(s-1)q+2}, \ldots, A_{s,q}$  form a partition of  $V(nK_q)$  for every  $s = 1, \ldots, r$ iff the corresponding **OA** (constructed in the proof of the above theorem) has the permutation property.

Theorem 3.2 and the results of § 2 imply

Corollary 3.4.

(i) 
$$\overline{\omega}(nK_2) = \operatorname{Min}\left\{ t / \left( \begin{bmatrix} t - 1 \\ \frac{t}{2} \end{bmatrix} - 1 \right) \ge n \right\},$$
  
(ii)  $\frac{\log n}{1 - \frac{1}{2}} \le \overline{\omega}(nK_q) \le \frac{q^2 \log n}{\log q} (1 + q)$ 

(ii) 
$$\frac{\log n}{\frac{1}{q}\log q + \left(1 - \frac{1}{q}\right)\log \frac{q}{q-1}} \leq \overline{\omega}(nK_q) \leq \frac{q^2 \log n}{\log q} (1 + o(1)) \text{ for } q \geq 3$$

#### 4. PRODUCT

By a categorical product of two graphs G, H we mean the following graph  $G \times H$ :  $V(G \times H) = V(G) \times V(H);$  $E(G \times H) = \{\{\langle x_1, y_1 \rangle, \langle x_2, y_2 \rangle\} | \{x_1, x_2\} \in E(G), \{y_1, y_2\} \in E(H)\}.$ 

# Theorem 4.1.

$$\overline{\omega}(G) + \overline{\omega}(H) \ge \overline{\omega}(G \times H) > \frac{1}{2} \left[ \log_2 \overline{\omega}(G) + \log_2 \overline{\omega}(H) \right].$$

Proof. Let  $C_1, C_2, ..., C_l$  and  $D_1, D_2, ..., D_k$  be systems of stable sets covering all pairs of nonadjacent vertices of G and H, respectively. Then  $C_i \times V(H)$ ,  $V(G) \times D_j$ , i = 1, ..., l, j = 1, ..., k, is a system covering all pairs of nonadjacent vertices of  $G \times H$ . This proves the first inequality.

For the proof of the second inequality we shall need the following three propositions.

Let G be a graph. Define an equivalence  $\sim$  on the set of its vertices as follows:  $x \sim y$  iff x and y have the same neighbourhoods. From each equivalence class choose one point and consider the subgraph induced on this set. Denote this graph by m(G).

**Proposition 4.2.**  $|V(m(G))| \ge \sqrt{(2\overline{\omega}(G))}$  for every graph G.

Proof. It follows from the well known theorem of Erdős, Goodman and Posa [4] which states that  $\omega(G) \leq |V(G)|^2/4$ , and from the fact that  $\overline{\omega}(m(G)) = \overline{\omega}(G)$ .

**Proposition 4.3.** 

$$m(G) \times m(H) = m(G \times H)$$

Proof. It suffices to realize that a neighbourhood of a vertex  $\langle x, y \rangle$  of the product is the product of neighbourhoods of vertices x and y.

**Proposition 4.4.**  $\overline{\omega}(F) \ge \log_2 |V(F)|$  for F = m(F).

Proof. Consider the disjoint representation  $\mathscr{S}$  of the graph F. From F = m(F) it follows that  $\mathscr{S}$  is a system of distinct sets and hence  $|V(F)| \leq 2^{\overline{\omega}(F)}$ .

Now we can prove the second inequality of 4.1:

$$\overline{\omega}(G \times H) = \overline{\omega}(\mathsf{m}(G \times H)) \ge \log_2(|V(\mathsf{m}(G \times H))|) =$$
  
=  $\log_2(V(\mathsf{m}(G)) \cdot V(\mathsf{m}(H))) = \log_2|V(\mathsf{m}(G))| + \log_2|V(\mathsf{m}(H))| >$   
>  $\frac{1}{2}[\log_2 \overline{\omega}(G) + \log_2 \overline{\omega}(H))].$ 

Remark 4.5. We have just proved:

If we denote  $f(r) = \min \{\overline{\omega}(G \times H) / \overline{\omega}(G) = \overline{\omega}(H) = r\}$  then  $f(r) \to \infty$  with  $r \to \infty$ .

An analogous statement for the chromatic number is not known (see [14]).

# 5. DIMENSION

The notion of dimension of graphs, defined in the introduction, was introduced in  $\lceil 11 \rceil$  and  $\lceil 8 \rceil$ . For a survey of recent results concerning this notion see  $\lceil 10 \rceil$ .

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The following proposition is shown in [11].

**Proposition 5.1.** The dimension of G equals the minimal number of equivalences  $E_1, \ldots, E_n$  such that

1)  $E(G) = \bigcup_{i=1}^{n} E_i;$ 2)  $\bigcap_{i=1}^{n} E_i = 0.$ 

In [8], dimension of the graphs  $nK_2$  was studied and it was shown that dim  $nK_2 = \log_2^+ n + 1$ . We give here some estimates of dim  $nK_q$  for  $q \ge 3$ .

**Proposition 5.2.** If there exists OA(q, n, qr) with the permutation property then dim  $nK_q \leq r$ .

Proof follows immediately from Theorem 3.3 and Proposition 5.1.

# Theorem 5.3.

- (i) dim  $(qK_q) = q$  for q a power of a prime;
- (ii) dim  $(n_1n_2K_q) \leq \dim(n_1K_q) + \dim(n_2K_q);$
- (iii) dim  $(nK_{q_1q_2}) \leq \dim nK_{q_1}$ . dim  $nK_{q_2}$ .

Proof. According to (ii) of 1.1, there exists  $OA(q, q, q^2)$  with the permutation property and hence according to 5.2, dim  $qK_q \leq q$ . The second inequality follows from the easy fact that dim  $(K_q + x) = q$ .  $(K_q + x)$  is the disjoint union of  $K_q$  and a single vertex x.)

(ii) As  $n_1n_2K_q$  is an induced subgraph of  $n_1K_q \times n_2K_q$  we have dim  $(n_1n_2K_q) \le$  $\le \dim (n_1K_q \times n_2K_q) \le \dim n_1K_q + \dim n_2K_q.$ 

(iii) It is easily seen that dim G = r if one can assign to each vertex x of G a vector  $\mathbf{v}(x)$  with r coordinates so that x = y implies  $\mathbf{v}(x) = \mathbf{v}(y)$  and  $(x, y) \in E(G)$  iff  $v_i(x) \neq v_i(y)$  for every  $i = 1, ..., r. (v_i(x)$  denotes the *i*-th coordinate of  $\mathbf{v}(x)$ .) We shall call such an assignment an encoding.

Denote by  $x_j^i$  and  $x_k^i$  the *j*-th and the *k*-th vertex of the *i*-th copy of  $K_{q_2}$  in the graph  $nK_{q_2}$  and  $nK_{q_2}$ , respectively. Put  $nK_{q_1q_2} = K_{q_1q_2}^1 + \ldots + K_{q_1q_2}^n$ , where  $V(K_{q_1q_2}^i) = \{x_{1,1}^i, \ldots, x_{j,k}^i, \ldots, x_{q_1,q_2}\}$ . Let **u** and **v** be an encoding of  $nK_{q_1}$  and  $nK_{q_2}$ , respectively. Define a mapping **w** as follows:  $\mathbf{w}_{\alpha,\beta}(x_{j,k}^i) = \langle \mathbf{u}_{\alpha}(x_{j}^i), \mathbf{v}_{\beta}(x_{k}^i) \rangle$ . It can be verified easily that **w** is an encoding.

Theorem 5.4.

dim 
$$nK_q \leq \frac{q \log n}{\log q} (1 + o(1))$$
 for  $q \geq 3$ 

Proof follows from 1.4, 5.2 and 0.3.

Let n > q be positive integers. Define the graph  $\binom{n}{q}$  whose vertices are q-point subsets of the set  $\{1, ..., n\}$  with two vertices adjacent if and only if they are disjoint subsets.

The graphs  $\binom{n}{q}$  have interesting properties. It is not difficult to see that they are universal, i.e. every finite graph is an induced subgraph of some  $\binom{n}{q}$ . Kneser conjectured that the chromatic number of  $\binom{2n+k}{n}$  is k+2. Lovasz proved this in [7] using methods of algebraic topology. In [13] we studied  $\overline{\omega}$  of  $\binom{n}{q}$  and proved that  $\overline{\omega} \binom{n}{q} = n$  for  $q \leq \frac{1}{2}n$ .

In [12] the dimension of Kneser graphs was studied. It was proved that

(\*) 
$$\log_2^+ \log_2^+ n - o(1) \leq \dim \binom{n}{q} \leq (q-1) q^2 \log_2^+ \log_2^+ n$$
.

Using 2.9 and [12, Remark 3.7] we obtain a slight improvement of (\*):

$$\dim \binom{n}{q} \leq \frac{4(q-1) q^2}{1 + \log_2 q} \log_2^+ \log_2^+ n \,.$$

 $(\log_2^+ n \text{ denotes the smallest integer not less then } \log_2 n)$ .

# 6. PROBLEMS

There are many open problems left. We shall mention only three of them.

- 1) Is  $nK_q$  an induced subgraph of a product of dim  $nK_q$  copies of  $K_q$ ?
  - (I.e., is the condition in Prop. 5.2 also necessary?)
- 2) Find a better estimates for f(n, q)!
- 3) Let G be an arbitrary graph. Is it true that

$$\overline{\omega}(nG) \geq \overline{\omega}(nK_{\chi(G)})?$$

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