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COMMENTATIONES MATHEMATICAE UNIVERSITATIS CAROLINAE

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GRAPHS WITH PRESCRIBED MAXIMAL SUBGRAPHS AND CRITICAL

CHROMATIC GRAPHS

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<u>Abstract</u>: It is proved that k-chromatic-critical graphs of large order contain large subgraphs of a certain structure. One of these results is that each large k-chromatic-critical graph contains a large odd circuit. A more general result is that if a large 2-connected graph G contains sub-. graphs of a certain structure of order N but not of order > Nthen G contains at least two disjoint isomorphic subgraphs not linked by an edge which are "isomorphically" connected to the rest G - H₁ - H₂ by edges. A so-called p-reduction is studied for such graphs.

Key words: Critical chromatic graph, subgraph, p-reduction.

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1. <u>Graphs with prescribed maximal subgraphs</u>. We consider undirected finite graphs without loops and multiple edges. If we handle with infinite graphs we say this explicitly. Further definitions are used as in L151. We say that a path t is a topological edge in a graph G, iff all inner vertices have in G the valency 2 and the two endvertices have a valency \geq 3. The class K of finite graphs is said to have property E if for each graph G \in K it holds: If t is an arbitrary topological edge only the endvertices of which are contained in G then G + t contains a subgraph G \in K with t \subseteq G'.

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Now we will present a few of such classes. A graph G is said to be contractible to a graph G' iff there exists a homomorphism φ from G onto G' with the properties:

1) For each vertex X' of G' it holds: The spanning subgraph of the vertex set $\mathcal{G}^{-1}(X')$ in G is a tree.

2) For each pair $\{X', Y'\}$ of different vertices of G' it holds:

a) $\varphi^{-1}(X')$ and $\varphi^{-1}(Y')$ are joined by at most one edge. b) X' and Y' are joined by an edge in G' iff $\varphi^{-1}(X')$ and $\varphi^{-1}(Y')$ are joined by an edge.

That means G' can be obtained from G by consecutive contractions of edges not contained in triangles.

A prismgraph consists of two circuits, which have at most one common vertex and which are united by three vertex-disjoint paths; if the two circuits have a common vertex then one path has length O.

<u>Theorem 1</u>: The class W of all paths , the class C of all circuits, the class O of all odd circuits, the class P of all prismgraphs, the class $\langle r, S \rangle$ of all 2-connected graphs contractible to a complete r-graph ($v \ge 4$) and the class $V(s_1, s_2, s_3, s_4, \ldots, s_p)$ have property E.

Each graph of the latter class can be formed as follows: We start with a set $\{H_j\}$ of $s_1 + s_2 + s_3 + s_4 + \dots + s_p$ pairwise disjoint graphs; s_1 of which are isolated vertices, s_2 , s_3 , s_4 ,..., s_p are graphs of C, P, $\langle 4, S \rangle$,..., $\langle p, S \rangle$, respectively. Then we consecutively link two components by a topological edge until we have obtained a connected graph. $V(s_1)$

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is the set of all trees with at most s_1 endvertices. $V(s_1,s_2)$ is the set of all Husimi trees with at most s_1 vertices of valency 1 and exactly s_2 circuits.

<u>Proof of Theorem 1</u>: Let $G \in W \cup C \cup C \cup P \cup ...$ and let V_1 , V_2 be the two different endvertices of a topological edge w with $G \cap w = \{V_1, V_2\}$.

a) Obviously, W and C have E.

b) Let $G \in O$. The two circuit arcs of G between V_1 and V_2 har' ve different parity. Let t denote the one which has the same parity as w. Then $G + W - t \in O$ with $w \subseteq G + w - t$.

c) For the class P the proof can easily be obtained by distinguishing some cases.

d) Let $G \in \langle r, S \rangle$ and let G' denote a complete r-graph. We have only to consider two cases:

1) There exists one vertex X' of G' and a homomorphism φ of G onto G' such that $V_1, V_2 \in \varphi^{-1}(X')$. The spanned subgraph U of $\varphi^{-1}(X')$ is a tree, therefore U + w contains exactly one circuit \overline{C} .

Let e be an edge of \overline{C} not contained in w. We delete in G + w the topological edge t of G + w containing e and we have a new tree U + w - t. It is easily to be seen that with φ (U + w - t) =_{def} X' we have G + w - t $\epsilon < r, S >$.

2) There exist two different adjacent vertices X' and Y' of G' and a homomorphism φ of G onto G' such that $V_1 \in \varphi^{-1}(X')$ and $V_2 \in \varphi^{-1}(Y')$. Then there is in G an edge e connecting $\varphi^{-1}(X')$ and $\varphi^{-1}(Y')$. Let t denote the topological edge of G + w containing e. Now it can easily be seen that $w \subseteq G + w - t$ and G + w - t is an element of $\langle r, S \rangle$. Thus

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the proof of d) is complete.

e) Let $G \in V(s_1, \ldots, s_p)$. If $w \cap H_j = \{V_1, V_2\}$ then the validity of the Theorem can easily be derived from a) ... d). If for all j it holds $||w \cap H_j|| \le 1$ then G + w has a circuit containing w and a topological edge t with $t \cap w \le \{V_1, V_2\}$ and for all j it holds $||t \cap H_j|| \le 1$. Then $w \le G + w - t \in V(s_1, \ldots, s_p)$. Q.e.d.

<u>Remark</u>: The class of all circuits containing a certain vertex X also has property E but the class of all circuits containing two certain vertices X and Y has not property E.

<u>Theorem 2</u>: a) Let K denote a class of finite graphs with property E and let N be a positive integer. Let G be a 2-connected (finite or infinite) graph which contains a graph $H \in K$ of order N but which does not contain an element of K of order > N. Then the length 1 of a maximal circuit L of G is $1 \le N^2$.

b) If K = 0 then $l \leq 2(N - 1)$.

In a) the bound is not best possible. That in b) the bound is best possible is shown by the graph which consists of two vertices of valency 3 which are linked by an edge and by two topological edges of length N - $\mathcal{L}(N \text{ odd})$.

By Theorem 2b in each 2-connected nonbipartite graph G it holds $\ell^* \neq \ell \neq 2(\ell^* - 1)$ where ℓ, ℓ^* denote the maximal circuit length and the maximal odd circuit length, respectively (provided ℓ^* exists).

A similar assertion for the maximal even circuit length does not hold. G.A. Dirac proved in [2] that in 2-connected finite graphs it holds that $\mathcal{L} - 1 \leq \overline{\mathcal{I}} \leq \mathcal{L}^2$ where $\overline{\mathcal{I}}$ denotes the maximal

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path length of G. It can be shown that $\ell - 1 \neq \overline{\ell} \neq \ell^2/4$ (see G.A. Dirac [2], see also [15] and [13] . In the following λ (H) and o(H) denote the number of edges or vertices of H,

respectively.

<u>Proof of Theorem 2</u>: a) We distinguish two cases. 1) Let $|| L \cap H || \leq 1$. Then there exist two disjoint paths w_1 , w_2 connecting L and H (possibly $\lambda(w_1) = 0$). Let $X_i =_{def} w_i \cap H$ and $Y_i =_{def} w_i \cap L$. Let L_1 , L_2 denote the two circuit arcs of L' between Y_1 and Y_2 . $p_i = w_1 + L_1 + w_2$ (i = 1,2) are two paths with $p_i \cap H = \{X_1, X_2\}$.

By Theorem 1 it follows that $H + p_i$ contains a subgraph H_i with $p_i \in H_i \in K$. Therefore

$$\begin{split} \lambda(\mathbf{w}_1) & \star \lambda(\mathbf{L}_1) + \lambda(\mathbf{w}_2) = \lambda(\mathbf{p}_1) \leq o(\mathbf{H}_1) \leq o(\mathbf{H}) = \mathbf{N}. \\ \text{Hence } \lambda(\mathbf{L}_1) \leq \mathbf{N} - 1 \text{ and } \mathcal{I} = \lambda(\mathbf{L}) \leq 2(\mathbf{N} - 1). \end{split}$$

2) Let $||L \cap H|| \ge 2$. Then L can be split up in arcs $L_1, L_2, ...$..., L_q such that for each i it holds: Either $L_1 \cap H = L_1$ or $L_1 \cap H = \{P_1^1, P_2^1\}$ where P_1^1, P_2^1 are the two endvertices of L_1 . In both cases we have $\mathcal{A}(L_1) \ge o(H) =$ = N (in the case $L_1 \cap H = \{P_1^1, P_2^1\}$ see 1)). Because $q \le N$ it follows $\mathcal{A}(L) \le N^2$.

b) Let $H \in O$. If L is and odd circuit $\mathcal{A}(L) = o(H) = N$. Now let L be an even circuit.

If $||H \cap L|| \le 1$ the assertion b) of the Theorem follows from the proof al).

Now let $||H \cap L|| \ge 2$. Let P and Q be two arbitrary vertices of $H \cap L$. Because $H \in O$ and L is an even circuit the parity of one of the two circuit arcs of H between P and Q is different from the parity of the two circuit arcs of L between P and Q. From

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this it can easily be derived that there exists a circuit arc \overline{H} of H with $\|\overline{H} \cap L\| = 2$ and the parity of \overline{H} is different from the parity of the two circuit arcs L_1, L_2 of L connecting the two vertices of $\overline{H} \cap L$ in L. Hence $L_1 + \overline{H} \in O$ and

 $\lambda(L) + 2 \lambda(\overline{H}) = \lambda(L_1 + \overline{H}) + \lambda(L_2 + \overline{H}) \leq 2 N.$

With $\lambda(\bar{H}) \ge 1$ the Theorem 2 is proved.

Let p be a positive integer. Now we will define the so called p-reduction of finite and infinite graphs described in [15]. Two finite subgraphs U_1 , U_2 of a finite or infinite graph G are called independent, if they do not have a common vertex and if they are not connected by an edge. Two finite subgraphs U_1 , U_2 of G are called equivalent, if $U_1 \equiv U_2$ or if U_1 and U_2 are independent and there exists an isomorphism of the grap G - U_1 onto the graph G - U_2 such that all vertices of G - U_1 - U_2 are fixed.

Let M be a class of pairwise independent finite subgraphs of G, then the above formulated so called equivalence is an equivalence relation in M. Therefore M is divided in equivalence classes. From each equivalence class with more than p elements we delete in G so many elements of this equivalence class that in G only p elements remain. We call this deletion an "elementary p-reduction".

A sequence of a finite number of elementary p-reductions is called a p-reduction. If the obtained graph is denoted by G' then we write $G \succ G'$.

Let K be a class of finite graphs with property E. Let N be an integer.

Z(K,N) denotes the class of all 2-connected finite and

infinite graphs which contain an element of K of order N but which do not contain an element of K of order > N. For this class we can prove a finiteness condition in the following sense:

Each large graph $G \in Z(K,N)$ contains p + 1 equivalent subgraphs; that means to each positive integer p there exists a positive integer n(p,K,N) such that every $G \in Z(K,N)$ of order $\ge n(p,K,N)$ contains p + 1 equivalent subgraphs. We define

$$\alpha(C,N) = \left[\frac{N}{2}\right] + 1 \text{ and } \beta(C,N) = \left[\frac{N}{2}\right] - 1,$$

$$\alpha(O,N) = N \quad \text{and} \quad \beta(O,N) = N - 2,$$

$$\alpha(K,N) = \beta(K,N) = \left[\frac{N^2}{2}\right] + 1, \text{ if } K \neq C, K \neq 0.$$

<u>Theorem 3</u>: Let p, N be integers with $N \ge 3$ and $p \ge \infty$ (K,N). Then

a) To each G Z(K,N) there exists a p-irreducible graph G' with G > G'. The graph G' can be obtained from G by a sequence of at most $\beta(K,N)$ elementary p-reductions.

b) Every p-irreducible graph G' with G' \prec G is up to isomorphism uniquely determined, is finite and it is also G' ϵ $\epsilon Z(K,N)$.

c) Z(K,N) only contains a finite number of unisomorphi pirreducible graphs.

In Theorem 3a) for some graphs of Z(C,N) and Z(O,N) we really need $\beta(C,N) = \left[\frac{N}{2}\right] - 1$ or $\beta(O,N) = N - 2$ elementary p-reductions, respectively, to obtain the p-irreducible graph. For Z(C,N) we have shown this in [15]. For Z(O,N) this is proved by the following graph:

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We consider a tree T with a root X such that the distance between X and each endvertex is N - 2, that each inner vertex has a valency $\ge p + 1$ and that every two inner vertices have different valencies. To this tree we add a new vertex Y and link it with X and the endvertices of T by edges.

Let $G \circ X$ denote the graph obtained from a graph G by adding a new vertex X and linking X to each vertex of G by an edge. Let $\overline{Z}(W,N)$ be the class of all connected graphs containing a path of length N but no path of length > N. Then G \in $\in \overline{Z}(W,N)$ iff $G \circ X \in Z(C, N + 2)$. Therefore it yields the

<u>Remark</u>: The Theorem 3 is also valid for $\overline{Z}(W,N)$ with $\overline{\alpha}(W,N) = \left[\frac{N}{2}\right] + 2$ and $\overline{\beta}(W,N) = \left[\frac{N}{2}\right]$.

<u>Proof of Theorem 3</u>: In the case K = C the Theorem was proved in [15], it is not proved here again. If $K \neq C$ then from Theorem 2 it follows:

 $Z(0,N) \subseteq \bigcup_{k=N}^{2(N-1)} Z(C,i) \text{ and}$ $Z(K,N) \subseteq \bigcup_{k=N}^{N^2} Z(C,i), \text{ if } K \neq 0.$

By applying the result already known for Z(C,i) we obtain the Theorem also in the case that $K \neq C$. It remains only to show that if $G \in Z(K,N)$ and $G \vdash G'$ then G' contains a subgraph of K of order N. But this can easily be done by taking the following into consideration: If $H \in K$, $H \subseteq G$ and o(H) = N, then each p-reduction can be chosen such that no vertex of H is deleted (notice that $p \ge N$ if $K \ne C$ and $p \ge N/2 + 1$ if K = C). Q.e.d.

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2. <u>Critical chromatic graphs</u>. In 2. we only consider finite graphs. The chromatic number of a graph G is denoted by $\chi(G)$. A graph is called k-critical, iff its chromatic number is k and by the deletion of an arbitrary edge the resulting graph has the chromatic number k - 1. In this paper only k-critical graphs with $k \ge 3$, are considered.

Lemma: Let G, G' be graphs with $G \succ G'$. Then $\chi(G') = \chi(G)$.

<u>Proof</u>: It suffices to show that if U_1 and U_2 are two equivalent subgraphs of G then $\chi(G - U_2) = \chi(G)$. But this can be seen by the fact that each suitable colouring of G - $-U_2$ can be extended to a suitable colouring of G by giving the same colour to the vertices X and $\varphi(X)$ for each $X \in U_1$ whereby φ denotes an isomorphism of G - U_2 onto G - U_1 with fixed G - $U_1 - U_2$. It is well known that each critical graph is 2-connected. Z(K,N,k) denotes the set of all k-critical graphs G $\in Z(K,N)$.

<u>Theorem 4</u>: Let K be a class of graphs with property E. Then Z(K,N,k) contains only a finite number of nonisomorphic graphs.

<u>Proof</u>: The Lemma shows that all graphs $G \in Z(K,N,k)$ are l-irreducible and also p-irreducible. Because Z(K,N) only contains a finite number of p-irreducible graphs (Theorem 3) the truth of Theorem 4 follows from $Z(K,N,k) \subseteq Z(K,N)$. Q.e.d.

Theorem 4 states that each k-critical graph of large order which has an element of K as a subgraph contains also a large graph of K. If F(K,n,k) is the largest integer such that every k-critical graph of order n which has an element of K

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as a subgraph, contains a subgraph of K of order $\geq F(K,n,k)$ then

(1) $\lim_{m \to \infty} F(K,n,k) = +\infty$

Obviously, for $k \ge 3$ each k-critical graph contains a subgraph $H' \in C$ and a subgraph $H' \in O$. Thus each large k-critical graph contains a large circuit and also a large odd circuit. The first assertion was proved by J.B. Kelly and L.M. Kelly [10] in 1954, the second assertion gives an answer of case $\mathcal{H} = 3$ of the question posed by J. Nešetřil and V. Rödl at the International Colloquium on Finite and Infinite Sets held in 1973 in Keszthely in Hungary(oral communication):

<u>Problem</u>: Let \mathcal{X} , k, N be arbitrary positive integers with $\mathcal{X} < k$. Does there exist a positive integer n such that each k-critical graph G with at least n vertices contains a \mathcal{X} -critical subgraph G' with at least N vertices?

The order of the magnitude of F(C,n,k) was investigated by J.B. Kelly and L.M. Kelly [10], G.A. Dirac [3] and R.C. Read [12]. T. Gallai [8] has obtained a sharpening of these results by showing that for an infinite set of different positive integers n there exist k-critical graphs of order n of maximal circuit length $\leq c_k \log n$, where c_k is an appropriate constant. From Theorem 2b) it follows that F(C,n,k) and F(0,n,k)have the same magnitude.

It also yields that the result "each large k-critical graph contains a large odd circuit" can be derived from the result of Kelly/Kelly "each large k-critical graph contains a large circuit" by means of Theorem 2b.

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Before discussing other classes K we define: If $r \ge 3$, then a topological complete r-graph consists of r branching vertices and of $\binom{n}{2}$ topological edges such that every two branching vertices are linked by exactly one topological edge. $\langle r, U \rangle$ denotes the class of all topological complete rgraph.

G.A. Dirac [1] has proved that each 4-critical graph contains a $\langle 4, U \rangle$. B. Zeidl [16] has shown that for $k \ge 4$ each k-critical graph has a $\langle 4, U \rangle$, containing a circuit of odd length.

In [4] G.A. Dirac has proved that each circuit of a 4-critical graph is contained in a $\langle 4, U \rangle$. If we apply this result to the largest circuits and to the largest odd circuits, then we obtain from (1) with respect to K = C and K = 0: For $k \ge 4$ each large k-critical graph has a large $\langle 4, U \rangle$ and also a large $\langle 4, U \rangle$ containing a circuit of odd length, respectively.

In this paper I proved the first statement again (see (1)) but I cannot reprove the second statement with the aid of Theorem 3 because the class of all graphs of $\langle 4, U \rangle$ containing an odd circuit has not property E. Because each k-critical graph has no vertex of valency $\leq k - 2$, every k-critical graph of order n has $\geq \frac{4}{2}$ (k - 1) n edges. This lower bound was improved by T. Gallai [8] and G.A. Dirac [6]. For $k \geq 6$ each k-critical graph contains at lest $\frac{5}{2}$ n edges. A result of G.A. Dirac [5] says that each simple graph of order $n \geq 5$ with at least $\frac{5}{2}$ n - 3 edges contains a graph

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obtained from a graph of the class $\langle 5, U \rangle$ by deleting one and only one topological edge. Because this graph has a special prismgraph it follows: For $k \ge 6$ each k-critical graph contains a prismgraph and with (1) each large k-critical graph contains a large prismgraph.

K. Wagner [14](H.A. Jung [9]) has proved: For every positive integer r there exists an integer k_r (an integer k_r') such that for all positive integers $k \ge k_r$ ($k \ge k_r'$) each kcritical graph contains a $\langle r, S \rangle$ (a $\langle r, U \rangle$) - also set W. Mader [11]. By (1) from this it follows: For all $k \ge k_r$ each large k-critical graph contains a large $\langle r, S \rangle$. But by our methods it cannot be shown that for all $k \ge k_r'$ each large k-critical graph contains a large $\langle r, U \rangle$ because $\langle r, U \rangle$ has not property E. We do also not know whether this assertion is true.

By definition we have

 $\bigcup_{n \to \infty} Z(K, N, k) \subseteq Z(K, N).$

Because by Theorem 3c the number of nonisomorphic graphs of Z(K,N) is finite we have that there exists a positive integer k(K,N) such that $Z(K,N,k) = \emptyset$ for all $k \ge k(K,N)$. By a result of P. Erdös and H. Hajnal we can take k(0,N) = N + 2 because they showed in [7]: Every graph which does not contain circuits of lengths 2j + 1 for all $j \ge i$ is suitable colourable by 2i colours.

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