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CONSTRUCTION OF CARTESIAN CLOSED TOPOLOGICAL HULLS
Jiří ADÁMEK, George E. STRECKER

Abstract: H. Herrlich and L.D. Nel proved that any category which has a cartesian closed topological extension, preserving finite products, has a smallest such extension, called the CCT hull. The present paper is devoted to a direct construction of this hull.

Key words: Topological category, cartesian closed category, initially complete category, CCT hull, power-closed sink.

Classification: 18D15, 18D99, 54A99

§ 0. Introduction. There are two essential properties of a concrete category for it to be "topologically adequate":

(a) Initial **completeness**, i.e., the existence of an initial structure for each structured source. An initially complete category has e.g. all limits and colimits constructed on the level of sets, and a lot of other convenient properties.

(b) Cartesian closedness, i.e., the existence of well-behaved function spaces.

The category of topological spaces fails in (b); the category of compactly generated Hausdorff spaces fails in (a) and the category of compact Hausdorff spaces fails in both. We are interested in extensions of a concrete category into a category with one, or both, of the properties (a) and (b).

Systematic methods for extending concrete categories to initially complete ones have been studied by many authors; see, e.g., [AHS_{1,2}; An₁; He_{2,3}; HS₂; Hu; R]. For example, in [AHS₁] it is shown that, whenever a category \mathcal{K} has any initial completion, then it has a smallest one, the so-called Mac-Neille completion, which can be described as the category of "closed" sinks in \mathcal{K} .

Cartesian closed topological categories have also been studied extensively; see, e.g., [AK; An₂; B; HN; M; N; W]. H. Herrlich and L.D. Nel proved in [HN] that whenever a category \mathcal{K} has any cartesian closed topological (CCT) extension, preserving finite products, then it has a smallest one, the so-called CCT hull of \mathcal{K} . The existence of a CCT hull is characterized in [AK] by the condition "strictly small-fibred", explained below.

In the present paper we introduce the notion of power-closed sinks, and, analogously to the case of Mac Neille completion, we prove that for any strictly small-fibred category its CCT hull is the category of power-closed sinks.

§ 1. The definition of a CCT hull

1.1. General assumptions. Throughout the paper, we deal with concrete categories (over sets); i.e., pairs (\mathcal{K}, \parallel) consisting of a category \mathcal{K} and a faithful, amnesic functor $\parallel : \mathcal{K} \rightarrow \text{Set}$. (Amnesicity means that any isomorphism f in \mathcal{K} , such that f is the identity map, is itself an identity morphism.) We use the same symbol for a morphism $f: A \rightarrow B$ in \mathcal{K} and its underlying map $f: |A| \rightarrow |B|$. Finally, we assume

that \mathcal{K} has at most one void object, i.e., an object A with $|A| = \emptyset$.

1.2. A structured map from a set X is a pair (A, a) , where A is an object of \mathcal{K} and $a: X \rightarrow |A|$ is a map; we denote it by

$$X \xrightarrow{a} |A|.$$

A family (possibly large) of structured maps from X is called a structured source on X . Let \mathcal{A} be a structured source on X ; then an object C with $|C| = X$ is the initial lift of the source \mathcal{A} if

- (i) $a: C \rightarrow A$ is a morphism for each $X \xrightarrow{a} |A|$ in \mathcal{A} ;
- (ii) given an object C' and a map $f: |C'| \rightarrow X$ such that $a \cdot f: C' \rightarrow A$ is a morphism for each $a \in \mathcal{A}$ then also $f: C' \rightarrow C$ is a morphism.

A concrete category is initially complete if each structured source has an initial lift.

1.3. Dually, a structured map into a set X is a map $a: |A| \rightarrow X$; we denote it by (A, a) or $|A| \xrightarrow{a} X$. A family of structured maps into X is a structured sink on X . The final lift of a structured sink \mathcal{A} is an object C with $|C| = X$ such that, given an object C' and a map $f: X \rightarrow |C'|$, then $f: C \rightarrow C'$ is a morphism iff each $f \cdot a: A \rightarrow C'$, $a \in \mathcal{A}$, is a morphism. Initial completeness is equivalent to each structured sink having a final lift.

1.4. Recall from [He₁] that a concrete category \mathcal{K} is topological if it is

- (i) initially complete;
- (ii) small-fibred, i.e., for every set X the collection

of all objects A with $|A| = X$ is a (small) set;

(iii) has constant morphisms, i.e., every constant map $f: |A| \rightarrow |B|$ is a \mathcal{K} -morphism $f: A \rightarrow B$.

1.5. A concrete category is said to have concrete finite products if it is finitely productive and its forgetful functor preserves finite products; (equivalently, if the product of any finite family of objects A_i , $i \in I$, is the initial lift of the source of projections

$$(X \xrightarrow{\pi_i} |A_i|)_{I}, \text{ where } X = \prod_{i \in I} |A_i| \text{ in Set}.$$

Particularly, each topological category has concrete finite products.

1.6. Let A and B be objects of a category \mathcal{K} with finite concrete products. Their (canonical) power-object

$$B^A$$

is an object on the set of all morphisms from A to B :

$$|B^A| = \text{hom}(A, B)$$

with the following universal property. Given an object D and a map $f: |D| \rightarrow \text{hom}(A, B)$ then

$$f: D \rightarrow B^A \text{ is a } \mathcal{K}\text{-morphism}$$

iff

$$\hat{f}: D \times A \rightarrow B \text{ is a } \mathcal{K}\text{-morphism,}$$

where \hat{f} is the map defined by: $\hat{f}(d, a) = [f(d)](a)$. This notion has been introduced by P. Antoine [An₂].

1.7. A cartesian closed topological (shortly CCT) category is a topological category such that each pair of its ob-

jects has a power-object. Equivalently, a CCT category is a topological category \mathcal{K} such that, for each object A , the functor

$$A \times - : \mathcal{K} \longrightarrow \mathcal{K}$$

has a right adjoint (namely, the functor $X \mapsto X^A$). This equivalence, and other important properties of CCT categories, are proved in [He₁].

1.8. Given a concrete category \mathcal{K} , we are interested in its finitely productive CCT extensions. (I.e., in CCT categories \mathcal{L} , containing \mathcal{K} as a full, concrete subcategory⁺, closed under finite products.) Note that if \mathcal{K} has such an extension then it has

- (i) concrete finite products,
- (ii) constant morphisms.

This follows from the fact that each CCT category has both. (These two conditions are not sufficient.) Even a topological category can fail to have such an extension, as proved in [AK].

On the other hand, if \mathcal{K} has a finitely productive CCT extension then it has a smallest one, called the CCT hull of \mathcal{K} . It can be characterized as a finitely productive CCT extension contained in each such extension. Also, it can be characterized internally as follows.

1.9. Definition [HN]. Let \mathcal{K} be a concrete category with finite concrete products and constant morphisms. Its CCT hull is a CCT category \mathcal{L} , in which \mathcal{K} is a full, concrete subcategory

 +) A concrete subcategory of a concrete category \mathcal{L} is a subcategory \mathcal{K} with an underlying functor arising as a restriction of that of \mathcal{L} .

tegrary such that

- (i) \mathcal{K} is finally dense ⁺) in \mathcal{L} .
- (ii) The power-objects of \mathcal{K} -objects are initially dense ⁺) in \mathcal{L} .

1.10. Remarks. (a) It is easy to check that, since \mathcal{K} is finally dense in \mathcal{L} , all initial lifts in \mathcal{K} are also initial lifts in \mathcal{L} . Particularly, since \mathcal{L} has finite concrete products, it follows that \mathcal{K} is closed to finite products in \mathcal{L} .

(b) In [HN] the condition (i) is stated in a seemingly stronger way: each \mathcal{L} -object L is a final lift of an episink $(|A_i| \xrightarrow{\alpha_i} X)$ of \mathcal{K} -objects, i.e., a sink such that $X = \cup a_i(|A_i|)$. Since \mathcal{K} has constant morphisms, these two conditions are equivalent: enlarging a given sink by arbitrary constant structured maps does not change the final lift.

(c) The **least** initial completion of \mathcal{K} , called the Mac Neille completion, is characterized analogously: it is an initially complete category \mathcal{L} in which \mathcal{K} is a full, concrete subcategory which is both finally and initially dense. See [AHS₁].

1.11. Definition. Let \mathcal{K} be a concrete category with finite concrete products. Two structured maps $|A| \xrightarrow{\alpha} X$ and $|A'| \xrightarrow{\alpha'} X$ are said to be productively equivalent, in symbols

$$(A, \alpha) \approx_X (A', \alpha'),$$

if for an arbitrary map

+) A class \mathcal{C} of objects of a concrete category \mathcal{L} is initially dense if each \mathcal{L} -object is the initial lift of a source of structured maps into \mathcal{C} -objects. Dually: finally dense.

$$h: X \times |B| \rightarrow |C|, \text{ where } B, C \in \mathcal{K},$$

we have;

$$h \circ (a \times 1_{|B|}): A \times B \rightarrow C \text{ is a } \mathcal{K}\text{-morphism}$$

iff

$$h \circ (a' \times 1_{|B|}): A' \times B \rightarrow C \text{ is a } \mathcal{K}\text{-morphism.}$$

1.12. Definition. A concrete category is said to be strictly small-fibred if it has finite concrete products and for each set X the productive equivalence \approx_X is small (i.e., it has a small set of representatives).

1.13. Theorem [AK]. Let \mathcal{K} be a concrete category with finite concrete products and with constant morphisms. Then \mathcal{K} has a CCT hull iff \mathcal{K} is strictly small-fibred.

Remark. "Usual" topological categories (and all of their full subcategories) are strictly small-fibred. Nevertheless, a topological category is constructed in [AK] which fails to be strictly small-fibred.

§ 2. The category of power-closed sinks

2.1. Throughout this section, \mathcal{K} denotes a fixed concrete category with finite concrete products and constant morphisms.

Given objects P and Q of \mathcal{K} , for each map $f: X \rightarrow \text{hom}(P, Q)$ we define a map $\hat{f}: X \times |P| \rightarrow |Q|$ by

$$\hat{f}(x, p) = [f(x)](p) \text{ for } x \in X, p \in |P|.$$

2.2. Convention. Let $\mathcal{A} = (|A_i| \xrightarrow{a_i} X)_{i \in I}$ be a structured sink. Denote by \mathcal{A}^\downarrow the source of all (non-structured!) maps $p: X \rightarrow \text{hom}(P, Q)$, where $P, Q \in \mathcal{K}$, such that for each

$i \in I$

$\widehat{p \cdot a_i}: A_i \times P \rightarrow Q$ is a morphism in \mathcal{K} .

2.3. Definition. Let \mathcal{A} be a structured sink on a set X . The power-closure of \mathcal{A} is the sink $\overline{\mathcal{A}}$ of all structured maps $|A| \xrightarrow{a} X$ with the following property: for each $p: X \rightarrow \text{hom}(P, Q)$ in \mathcal{A}^\downarrow ,

$\widehat{p \cdot a}: A \times P \rightarrow Q$ is a morphism in \mathcal{K} .

If $\overline{\mathcal{A}} = \mathcal{A}$, we call \mathcal{A} a power-closed sink.

Note that each structured sink \mathcal{A} fulfils $\mathcal{A} \subset \overline{\mathcal{A}} = \overline{\overline{\mathcal{A}}}$. The fact that \mathcal{A} is power-closed means that \mathcal{A} is "determined" by the source \mathcal{A}^\downarrow .

2.4. Example. For each object B of \mathcal{K} denote by B^0 the following structured sink on the set $X = |B|$:

$B^0 = \{ |A| \xrightarrow{a} X \mid a: A \rightarrow B \text{ is a morphism in } \mathcal{K} \}$;

then B^0 is power-closed.

Proof. Let $|A| \xrightarrow{a} X$ belong to B^0 ; we shall prove that $a: A \rightarrow B$ is a morphism in \mathcal{K} .

We use the map $p: X \rightarrow \text{hom}(B, B)$, assigning to each $x \in X$ the morphism $p(x): B \rightarrow B$, which is constant with the value x . Note that $p \in (B^0)^\downarrow$ since for each morphism $a_0: A_0 \rightarrow B$ the map $\widehat{p \cdot a_0}: |A_0 \times B| \rightarrow |B|$ is defined by

$$\widehat{p \cdot a_0}(t, x) = [p(a_0(t))](x) = a_0(t).$$

Thus,

$$\widehat{p \cdot a_0} = \pi_1 \cdot (a_0 \times 1_B): A_0 \times B \rightarrow B,$$

where $\pi_1: B \times B \rightarrow B$ is the first projection. Therefore, $\widehat{p \cdot a_0}: A_0 \times B \rightarrow B$ is a morphism in \mathcal{K} ; i.e., an element of B^0 .

Since $a \in B^0$ and $p \in (B^0)^\downarrow$, by the definition of power-

closure $\widehat{p \cdot a}: A \times B \rightarrow B$ is a morphism of \mathcal{K} . It suffices to exhibit a morphism $r: A \rightarrow A \times B$ such that $a = \widehat{p \cdot a} \cdot r$. For this choose an arbitrary morphism $r_0: A \rightarrow B$. This is possible since if $X \neq \emptyset$ then r_0 can be any constant map; if $X = \emptyset$ then $|A| = \emptyset$ (because we have a map $a: |A| \rightarrow X$) and, by the standing hypothesis of 1.1, it follows that $A = B$ and $r_0 = 1_B$. Now, let $r: A \rightarrow A \times B$ have components 1_A and r_0 . Then, for each $t \in |A|$,

$$\widehat{p \cdot a} \cdot r(t) = \widehat{p \cdot a}(t, r_0(t)) = [p \cdot a(t)](r_0(t)) = a(t).$$

Thus,

$$a = \widehat{p \cdot a} \cdot r: A \rightarrow B$$

is a morphism of \mathcal{K} , which was to be proved.

2.5. Remark. Each power-closed sink \mathcal{A} on a set X is a sieve in the terminology of P. Antoine [An₂]; i.e.,

(i) \mathcal{A} is closed under composition (in the sense that for each $|A| \xrightarrow{a} X$ in \mathcal{A} and each morphism $f: A' \rightarrow A$ we have $|A'| \xrightarrow{a \cdot f} X$ in \mathcal{A});

(ii) \mathcal{A} contains all constant structured maps into X .

Proof. (ii) Let $|A| \xrightarrow{a} X$ be constant with a value $x_0 \in X$. For each map $p: X \rightarrow \text{hom}(P, Q)$ the map $\widehat{p \cdot a}: |A \times P| \rightarrow |Q|$ is defined by

$$\widehat{p \cdot a}(t, x) = [p(a(t))](x) = [p(x_0)](x).$$

Thus $\widehat{p \cdot a}$ is the composition of the second projection $A \times P \rightarrow P$ and the morphism $p(x_0): P \rightarrow Q$. Thus, $\widehat{p \cdot a}$ is a morphism in \mathcal{K} .

Since this holds in particular for each $p \in \mathcal{A}^\downarrow$, we see that $a \in \overline{\mathcal{A}} = \mathcal{A}$.

(i) Let $p \in \mathcal{A}^\downarrow$ be arbitrary. We know that $\widehat{p \cdot a}: A \times$

$\times P \rightarrow Q$ is a morphism and we are to prove that so is

$$\widehat{p \cdot a \cdot f}: \widehat{A' \times P} \rightarrow Q.$$

This is a consequence of the fact that $f \times 1_P: A' \times P \rightarrow A \times P$ is a morphism and the following lemma.

2.6. Lemma. Given objects P and Q and maps

$$Y \xrightarrow{f} X \xrightarrow{g} \text{hom}(P, Q)$$

then

$$\widehat{g \cdot f} = \widehat{g} \cdot (f \times 1_{|P|}): Y \times |P| \rightarrow |Q|.$$

Proof. For arbitrary $y \in Y$, $p \in |P|$ we have

$$\widehat{g \cdot f}(y, p) = [g(f(y))](p)$$

as well as

$$\widehat{g} \cdot (f \times 1_{|P|})(y, p) = \widehat{g}(f(y), p) = [g(f(y))](p).$$

2.7. Proposition. Let \mathcal{K} be a strictly small-fibred category. Then for each set X the conglomerate of all power-closed sinks on X is small.

Proof. Since the equivalence \approx_X of Definition 1.11 has a set of representatives (say, of cardinality α), it suffices to prove that each power-closed sink \mathcal{A} is closed under this equivalence, i.e.,

$|A| \xrightarrow{a} X$ in \mathcal{A} implies $|A'| \xrightarrow{a'} X$ is in \mathcal{A} whenever

$$(A, a) \approx_X (A', a').$$

Then power-closed sinks can be indexed by sets of representatives of \approx_X ; hence, the number of all power-closed sinks on the set X cannot exceed 2^α .

Let \mathcal{A} be power-closed and let $(A, a) \approx_X (A', a')$. Assuming that $|A| \xrightarrow{a} X$ is in \mathcal{A} , we are to show that $|A'| \xrightarrow{a'} X$ is in \mathcal{A} . Given $p: X \rightarrow \text{hom}(P, Q)$ in \mathcal{A}^\downarrow , we know that

$\widehat{p \cdot a}: A \times P \rightarrow Q$ is a \mathcal{K} -morphism. By Lemma 2.6, this means that

$\widehat{p \cdot (a \times 1_{|P|})}: A \times P \rightarrow Q$ is a \mathcal{K} -morphism.

Since $(A, a) \approx_X (A', a')$, we have

$\widehat{p \cdot (a' \times 1_{|P|})} = \widehat{p \cdot a'}: A' \times P \rightarrow Q$ is a \mathcal{K} -morphism.

Thus, since \mathcal{A} is power-closed, (A', a') must belong to it.

2.8. Corollary. For each strictly small-fibred category, the conglomerate of all power-closed sinks is legitimate, i.e., is isomorphic to a class (in the Bernays-Gödel terminology).

Even if all power-closed sinks form a legitimate conglomerate, we cannot, strictly speaking, work with the "class of all power-closed sinks". (Since a power-closed sink, which is itself a proper class, cannot be a member of a class.) Nevertheless, we shall disregard this difficulty which is, evidently, only formal: instead of the "class" of power-closed sinks, considered below, we would formally work with an arbitrary class isomorphic to it.

2.9. Definition. Let \mathcal{X} be a strictly small-fibred category. Then its category of power-closed sinks is the following concrete category, denoted by $\text{PCS}(\mathcal{X})$:

Objects are all power-closed sinks.

Morphisms from a sink $\mathcal{A} = (A_i \xrightarrow{a_i} X)_I$ to a sink $\mathcal{B} = (B_j \xrightarrow{b_j} Y)_J$ are maps $f: X \rightarrow Y$ such that for each $i \in I$ there exists $j \in J$ with $A_i = B_j$ and $f \cdot a_i = b_j$;

the forgetful functor is defined by $|(A_i \xrightarrow{a_i} X)| = X$.



It is easy to see that $\text{PCS}(\mathcal{K})$ is indeed a correctly defined concrete category (up to the tolerance mentioned in 2.8).

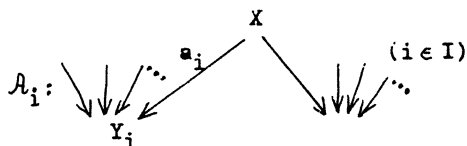
Identifying \mathcal{K} -objects B with the power-closed sinks B^0 of 2.4, the category \mathcal{K} becomes a full, concrete subcategory of $\text{PCS}(\mathcal{K})$. Indeed, each \mathcal{K} -morphism $f: B \rightarrow C$ is clearly a sink-morphism $f: B^0 \rightarrow C^0$. Conversely, if $f: B^0 \rightarrow C^0$ is a sink-morphism then

$$|B| \xrightarrow{1_B} |B| \text{ in } B^0 \text{ implies } |B| \xrightarrow{f \cdot 1_B} |C| \text{ in } C^0,$$

hence, $f: B \rightarrow C$ is a \mathcal{K} -morphism.

2.10. Proposition. Let \mathcal{K} be a strictly small-fibred category with constant morphisms. Then $\text{PCS}(\mathcal{K})$ is a topological category.

Proof. (i) $\text{PCS}(\mathcal{K})$ is initially complete. To prove this, consider a structured source $(X \xrightarrow{a_i} |A_i|)_I$



to power-closed sinks A_i . Define a structured sink \mathcal{C} on X to consist of precisely those structured maps $|C| \xrightarrow{c} X$ which satisfy, for each $i \in I$:

$$|C| \xrightarrow{a_i \cdot c} Y_i \text{ is in } A_i.$$

A) \mathcal{C} is power-closed. Let $|D| \xrightarrow{d} X$ be a structured map in \mathcal{C} . We are to show that $(D, d) \in \mathcal{C}$.

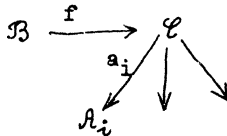
$$\begin{array}{ccc} |D| & \xrightarrow{d} & X \\ & & \downarrow a_i \\ & & Y_i \xrightarrow{\eta} \text{hom}(P, \mathcal{C}). \end{array}$$

For each $i \in I$ the structured map $|D| \xrightarrow{a_i \cdot d} Y_i$ is in $\overline{\mathcal{A}}_i$.
 (Indeed, let $p: Y_i \rightarrow \text{hom}(P, Q)$ be an element of \mathcal{A}_i^\downarrow ; then
 $p \cdot a_i: X \rightarrow \text{hom}(P, Q)$ is easily seen to be an element of \mathcal{C}^\downarrow .
 Hence, by the hypothesis on d ,

$$\widehat{p \cdot a_i \cdot d}: D \times P \rightarrow Q \text{ is a } \mathcal{K}\text{-morphism.}$$

Since \mathcal{A}_i is a power-closed sink, it follows that $a_i \cdot d$ belongs to it (for each $i \in I$); in other words, d is an element of \mathcal{C} .

B) \mathcal{C} is the initial lift of the given source. Let \mathcal{B} be a power-closed sink on a set Z and let $f: Z \rightarrow X$ be



a map such that $a_i \cdot f: \mathcal{B} \rightarrow \mathcal{A}_i$ is a sink-morphism for each $i \in I$. We are to show that then $f: \mathcal{B} \rightarrow \mathcal{C}$ is a sink-morphism. Indeed, given $|B| \xrightarrow{b} Z$ in \mathcal{B} we know that

$$|B| \xrightarrow{a_i \cdot f \cdot b} Y_i \text{ is in } \mathcal{A}_i$$

for each $i \in I$. This means that $(B, f \cdot b)$ is in \mathcal{C} .

(ii) $\text{PCS}(\mathcal{K})$ is small-fibred; see Proposition 2.7.

(iii) $\text{PCS}(\mathcal{K})$ has constant morphisms. This follows immediately from part (ii) of Lemma 2.5.

2.11. Proposition. For each strictly small-fibred category with constant morphisms \mathcal{K} , the category $\text{PCS}'(\mathcal{K})$ is cartesian closed. The power-object of sinks \mathcal{A} and \mathcal{B} (with $|\mathcal{A}| = X$; $|\mathcal{B}| = Y$) is the sink $\mathcal{B}^{\mathcal{A}}$ of all structured maps $|C| \xrightarrow{c} \text{hom}(\mathcal{A}, \mathcal{B})$ such that

$|A| \xrightarrow{a} X$ in \mathcal{A} implies $|C \times A| \xrightarrow{\hat{c} \cdot (1_C \times a)} Y$ in \mathcal{B} .

Proof. A) The above sink $\mathcal{B}^{\mathcal{A}}$ is power-closed. To prove this, consider an arbitrary structured map

$|C_0| \xrightarrow{c_0} \text{hom}(\mathcal{A}, \mathcal{B})$ in $\overline{\mathcal{B}^{\mathcal{A}}}$. We shall prove then (C_0, c_0) is an element of $\mathcal{B}^{\mathcal{A}}$. Thus, given $|A_0| \xrightarrow{a_0} X$ in \mathcal{A} we are to verify that the structured map

$$|C_0 \times A_0| \xrightarrow{1 \times a_0} |C_0| \times X \xrightarrow{\hat{c}_0} Y$$

is in \mathcal{B} . Now, \mathcal{B} is power-closed; therefore it suffices to show that $\hat{c}_0 \cdot (1 \times a_0) \in \overline{\mathcal{B}}$. Hence, for each $p_0: Y \rightarrow \text{hom}(P, Q)$ in \mathcal{B}^\downarrow we shall prove that

$$\widehat{p_0 \cdot \hat{c}_0 \cdot (1 \times a_0)}: C_0 \times A_0 \times P \rightarrow Q \text{ is a } \mathcal{K}\text{-morphism,}$$

thus concluding the proof of A).

First, we define a map

$$p: \text{hom}(\mathcal{A}, \mathcal{B}) \rightarrow \text{hom}(A_0 \times P, Q),$$

for which we shall verify that it is an element of $(\mathcal{B}^{\mathcal{A}})^\downarrow$.

Let $h \in \text{hom}(\mathcal{A}, \mathcal{B})$ be any sink-morphism; then $|A_0| \xrightarrow{a_0} X$ in \mathcal{A} implies

$$|A_0| \xrightarrow{h \cdot a_0} Y \text{ is in } \mathcal{B}.$$

Since $p_0 \in \mathcal{B}^\downarrow$, it follows that

$$\widehat{p_0 \cdot h \cdot a_0}: A_0 \times P \rightarrow Q \text{ is a } \mathcal{K}\text{-morphism.}$$

Put

$$p(h) = \widehat{p_0 \cdot h \cdot a_0} \text{ for each } h \in \text{hom}(\mathcal{A}, \mathcal{B}).$$

We prove that p is in $(\mathcal{B}^{\mathcal{A}})^\downarrow$, i.e., that given

$$|C| \xrightarrow{c} \text{hom}(\mathcal{A}, \mathcal{B}) \text{ in } \mathcal{B}^{\mathcal{A}}$$

then

$\widehat{p \cdot c}: C \times (A_0 \times P) \rightarrow Q$ is a \mathcal{K} -morphism.

Since c is in $\mathcal{B}^{\mathcal{A}}$ and a_0 is in \mathcal{A} , we conclude that

$$\widehat{c \cdot (1 \times a_0)}: |C \times A_0| \rightarrow Y \text{ is in } \mathcal{B}.$$

Thus, $p_0 \in \mathcal{B}^\downarrow$ implies that

$$\widehat{p_0 \cdot \widehat{c \cdot (1 \times a_0)}}: C \times A_0 \times P \rightarrow Q \text{ is a } \mathcal{K}\text{-morphism}$$

and it suffices to show that

$$\widehat{p \cdot c} = \widehat{p_0 \cdot \widehat{c \cdot (1 \times a_0)}}.$$

Indeed, for arbitrary points $z \in |C|$, $x \in |A_0|$ and $t \in |P|$ we have

$$\begin{aligned} \widehat{p \cdot c}(z, x, t) &= [p(c(z))](x, t) = [\widehat{p_0 \cdot c(z) \cdot a_0}](x, t) = \\ &= [(p_0 \cdot c(z))(a_0(x))](t) \end{aligned}$$

as well as

$$\begin{aligned} \widehat{p_0 \cdot \widehat{c \cdot (1 \times a_0)}}(z, x, t) &= [(p_0 \cdot \widehat{c \cdot (1 \times a_0)})](z, x)(t) = \\ &= [(p_0 \cdot \widehat{c})(z, a_0(x))](t) = [(p_0 \cdot c(z))(a_0(x))](t). \end{aligned}$$

Since p is in $(\mathcal{B}^{\mathcal{A}})^\downarrow$ and $c_0 \in \overline{\mathcal{B}^{\mathcal{A}}}$ we conclude that

$$\widehat{p \cdot c_0}: C \times A_0 \times P \rightarrow Q \text{ is a } \mathcal{K}\text{-morphism.}$$

This concludes the proof since, as above, $\widehat{p \cdot c_0} = \widehat{p_0 \cdot \widehat{c_0 \cdot (1 \times a_0)}}$.

B) The category PCS (\mathcal{K}) has finite concrete products, since it is topological. It remains to be shown that the sinks $\mathcal{B}^{\mathcal{A}}$ have the required universal property. Indeed, let \mathcal{A} , \mathcal{B} and \mathcal{C} be power-closed sinks with underlying sets X , Y and Z respectively; let $f: Z \rightarrow \text{hom}(\mathcal{A}, \mathcal{B})$ be an arbitrary map. Then

$f: \mathcal{C} \rightarrow \mathcal{B}^{\mathcal{A}}$ is a sink-morphism

iff for arbitrary $|C| \xrightarrow{c} Z$ in \mathcal{C} and $|A| \xrightarrow{a} X$ in \mathcal{A} we have:

$$\widehat{f \cdot c} \cdot (1_C \times a): |C \times A| \rightarrow Y \text{ is in } \mathcal{B} .$$

On the other hand, the product $\mathcal{C} \times \mathcal{A}$ consists of all structured maps

$$|D| \xrightarrow{(c, a)} Z \times X$$

where (c, a) is the map with components $|D| \xrightarrow{c} Z$ in \mathcal{C} and $|D| \xrightarrow{a} X$ in \mathcal{A} . Hence

$\widehat{f}: \mathcal{C} \times \mathcal{A} \rightarrow \mathcal{B}$ is a sink-morphism iff for arbitrary $|C| \xrightarrow{c} Z$ in \mathcal{C} and $|A| \xrightarrow{a} X$ in \mathcal{A} we have:

$$\widehat{f} \cdot (c \times a): |C \times A| \rightarrow Y \text{ is in } \mathcal{B} .$$

But by Lemma 2.6,

$$\widehat{f} \cdot (c \times a) = \widehat{f \cdot c} \cdot (1_C \times a),$$

hence the two conditions on f coincide.

2.12. Remark. Particularly, given objects A and B in \mathcal{K} , the power-sink $(B^0)^{\mathcal{A}^0}$ consists of those structured maps $|C| \xrightarrow{c} \text{hom}(A, B)$ which fulfil:

$$\widehat{c}: C \times A \rightarrow B \text{ is a } \mathcal{K}\text{-morphism.}$$

We denote this sink by B^A .

§ 3. The description of the CCT hull

3.1. Theorem. For each strictly small-fibred category \mathcal{K} with constant morphisms the CCT hull is the category PCS (\mathcal{K}) of power-closed sinks.

Proof. We know that PCS (\mathcal{K}) is a CCT category (2.10

and 2.11) which contains \mathcal{K} as a full, concrete subcategory (2.9).

(i) \mathcal{K} is finally dense in PCS (\mathcal{K}). Indeed, each power-closed sink $\mathcal{A} = (|A_i| \xrightarrow{a_i} X)_I$ is the final lift of itself, i.e., more precisely, of the sink $(|A_i^0| \xrightarrow{a_i} X)_I$.

(ii) The power-objects of \mathcal{K} are initially dense in PCS (\mathcal{K}). Let $\mathcal{A} = (|A_i| \xrightarrow{a_i} X)_I$ be a power-closed sink with

$$\mathcal{A}^\downarrow = (X \xrightarrow{P_j} \text{hom}(P_j, Q_j))_J.$$

We shall prove that then \mathcal{A} is the initial lift of the source

$$(X \xrightarrow{P_j} |Q_j^{P_j}|)_J$$

(where the sinks $Q_j^{P_j}$ are as described in 2.12).

First, for each $j \in J$,

$$p_j: \mathcal{A} \longrightarrow Q_j^{P_j}$$

is a sink-morphism. Indeed, given $|A| \xrightarrow{a} X$ in \mathcal{A} then $p_j \in \mathcal{A}^\downarrow$ implies that

$$\widehat{p_j \cdot a}: A \times P_j \longrightarrow Q_j \text{ is a } \mathcal{K}\text{-morphism,}$$

in other words that

$$|A| \xrightarrow{p_j \cdot a} \text{hom}(P_j, Q_j) \text{ is an element of } Q_j^{P_j}.$$

Secondly, let \mathcal{B} be a power-closed sink and let $f: |\mathcal{B}| \rightarrow X$ be a map such that $p_j \cdot f: \mathcal{B} \rightarrow Q_j^{P_j}$ are sink-morphisms

$$\begin{array}{ccc} \mathcal{B} & \xrightarrow{f} & \mathcal{A} \\ & \searrow p_j & \downarrow \\ & P_j & \\ & Q_j & \dots \end{array}$$

for all $j \in J$. We are to show that also $f: \mathcal{B} \rightarrow \mathcal{A}$ is a

sink-morphism, i.e., that given $|B| \xrightarrow{b} Y$ in \mathcal{B} then $|B| \xrightarrow{f \cdot b} X$ is in $\overline{\mathcal{A}}$. (Then, of course, $f \cdot b$ is an element of \mathcal{A} .) For each p_j in \mathcal{A}^\downarrow we know that $p_j \cdot f$ is a **sink-morphism**, thus

$$p_j \cdot f \cdot b: |B| \rightarrow \text{hom}(P_j, Q_j) \text{ is in } Q_j^{P_j}.$$

This means that

$$\widehat{p_j \cdot f \cdot b}: B \times P_j \rightarrow Q_j \text{ is a } \mathcal{K}\text{-morphism.}$$

3.2. Corollary. For each category \mathcal{K} with finite concrete products and constant morphisms the following conditions are equivalent:

- (i) \mathcal{K} has a finitely productive CCT extension;
- (ii) \mathcal{K} has a CCT hull;
- (iii) for each set X the conglomerate of all power-closed sinks on X is small;
- (iv) $\text{PCS}(\mathcal{K})$ is a CCT hull of \mathcal{K} ;
- (v) \mathcal{K} is strictly small-fibred.

Let us remark that the proof of (v) \implies (ii), presented in the current paper, is much simpler than the original proof of [AK], where an extension of \mathcal{K} is constructed by a complicated transfinite induction.

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