## Časopis pro pěstování matematiky

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Časopis pro pěstování matematiky, Vol. 101 (1976), No. 2, 199--202
Persistent URL: http://dml.cz/dmlcz/117895

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# ON THE MINIMUM DEGREE AND EDGE-CONNECTIVITY OF A GRAPH 

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(Received July 25, 1975)

Let $G$ be a graph. We denote by $V(G)$ and $E(G)$ the vertex set of $G$ and the edge set of $G$, respectively. If $v \in V(G)$, then we denote by $\operatorname{deg}_{G} v$ the degree of $v$ in $G$. Moreover, we denote by $\delta(G)$ and $\Delta(G)$ the minimum degree of $G$ and the maximum degree of $G$, respectively. If $U$ is a nonempty subset of $V(G)$, then we denote by $\langle U\rangle_{G}$ the graph $G^{\prime}$ such that $V\left(G^{\prime}\right)=U$ and

$$
E\left(G^{\prime}\right)=\{e \in E(G) ; e \text { is incident with no vertex in } V(G)-U\} .
$$

Let $G$ be a nontrivial connected graph. We say that a set $S \subseteq E(G)$ is.a cut set of $G$, if the graph $G-S$ is disconnected. A cut set $S$ with $|S|=n$ is referred to as an $n$-cut set. We denote by $\chi_{1}(G)$ the minimum integer $k$ such that there is a $k$-cut set of $G$; the integer $\varkappa_{1}(G)$ is called the edge-connectivity of $G$. (The terms not defined here can be found in M. Behzad and G. Chartrand [1].)

It is obvious that for any nontrivial connected graph $G, \chi_{1}(G) \leqq \delta(G)$. A sufficient condition for $\varkappa_{1}(G)=\delta(G)$ is due to D. R. Lick [3]; note that Lick's result is an analogue of R. Halin's theorem on the vertex-connectivity [2]. In the present note it will be shown that an analysis of a nontrivial connected graph from the point of view of its edge-connectivity can lead to an upper bound for the minimum degree. In fact, we obtain an upper bound for a more general characteristic: if $G$ is a graph and $U$ is a nonempty subset of $V(G)$, then we denote

$$
\operatorname{deg}_{G} U=\min \left\{\operatorname{deg}_{G} u ; u \in U\right\}
$$

Obviously, $\delta(G)=\operatorname{deg}_{G} V(G)$.
Let $G$ be a nontrivial connected graph, and let $U$ be a nonempty subset of $V(G)$. We denote $f_{G}(U)=\Delta\left(G_{0}\right)$, where $G_{0}$ is the spanning subgraph of $\langle U\rangle_{G}$ with the property that $e \in E\left(G_{0}\right)$ if and only if $e \in E\left(\langle U\rangle_{G}\right)$ and $x_{1}(G-e)=x_{1}(G)$. We denote by $h_{G}(U)$ the minimum integer $i$ such that there is an $i$-cut set $R_{0}$ of $G$ with
the property that for at least one component $F_{0}$ of the graph $G-R_{0}$ it holds that $V\left(F_{0}\right) \subseteq U$. Obviously, $x_{1}(G) \leqq h_{G}(U)$. It is clear that if $U_{1}$ and $U_{2}$ are subsets of $V(G)$ such that $\emptyset \neq U_{1} \subseteq U_{2}$, then $f_{G}\left(U_{1}\right) \leqq f_{G}\left(U_{2}\right)$ and $h_{G}\left(U_{2}\right) \leqq h_{G}\left(U_{1}\right)$. Denote $f_{G}=f_{G}(V(G))$. Clearly, $h_{G}(V(G))=x_{1}(G)$.

The following theorem is the main result of this note:

Theorem. Let $G$ be a nontrivial connected graph, and let $U$ be a nonempty subset of $V(G)$. Then

$$
h_{G}(U) \leqq \operatorname{deg}_{G} U \leqq \max \left(f_{G}(U), h_{G}(U)\right)
$$

Proof. It is obvious that for each $u \in U$, the set of edges incident with $u$ in $G$ is a cut set of $G$. Therefore, $h_{G}(U) \leqq \operatorname{deg}_{G} U$.

We shall prove the inequality $\operatorname{deg}_{G} U \leqq \max \left(f_{G}(U), h_{G}(U)\right)$. Clearly, there is a nonempty subset $U_{0}$ of $U$ such that $h_{G}\left(U_{0}\right)=h_{G}(U)$ and that for each nonempty subset $U^{\prime}$ of $U,\left|U^{\prime}\right|<\left|U_{0}\right|$ implies $h_{G}\left(U^{\prime}\right)>h_{G}(U)$. Obviously, $U_{0} \neq V(G)$. We denote by $F$ the graph $\left\langle U_{0}\right\rangle_{G}$. It is obvious that there is an $h_{G}(U)$-cut set $R$ of $G$ such that $F$ is a component of $G-R$. It is easy to see that $E(F) \cap R=\emptyset$. Denote $n=\left|U_{0}\right|$. Obviously,

$$
\begin{equation*}
\operatorname{deg}_{G} U \leqq \operatorname{deg}_{G} U_{0} \leqq \Delta(F)+\left[\frac{h_{G}(U)}{n}\right] \leqq n-1+\frac{h_{G}(U)}{n} . \tag{1}
\end{equation*}
$$

(Note that if $x$ is a real number, then $[x]$ denotes the maximum integer $j$ such that $j \leqq x$.)

Let $n \leqq h_{G}(U)$. If $h_{G}(U)<\operatorname{deg}_{G} U$, then it follows from (1) that $h_{G}(U)$. $(n-1)<$ $<n(n-1)$, and thus $h_{G}(U)<n$, which is a contradiction. Hence $\operatorname{deg}_{G} U \leqq h_{G}(U) \leqq$ $\leqq \max \left(f_{G}(U), h_{G}(U)\right)$.

Let $h_{G}(U)<n$. From (1) it follows that $\operatorname{deg}_{G} U \leqq \Delta(F)$. We distinguish two cases:
(I) For each $e \in E(F), x_{1}(G-e)=x_{1}(G)$. Then $f_{G}(U) \geqq \Delta(F)$. Therefore, $\operatorname{deg}_{G} U \leqq \Delta(F) \leqq f_{G}(U) \leqq \max \left(f_{G}(U), h_{G}(U)\right)$.
(II) There exists $e \in E(F)$ such that $x_{1}(G-e) \neq x_{1}(G)$. Then there exists a $x_{1}(G)$-cut set $S$ of $G$ such that $e \in S$. Obviously, the graph $G-S$ has precisely two components, say $G_{1}$ and $G_{2}$, and $E\left(G_{1}\right) \cap S=\emptyset=E\left(G_{2}\right) \cap S$. We denote by $H$ the graph $G-U_{0}$. It is easy to see that $E(H) \cap R=\emptyset$. Next, we denote $V_{11}=U_{0} \cap$ $\cap V\left(G_{1}\right), \quad V_{12}=U_{0} \cap V\left(G_{2}\right), \quad V_{21}=V(H) \cap V\left(G_{1}\right), \quad$ and $\quad V_{22}=V(H) \cap V\left(G_{2}\right)$. Finally, we denote by

$$
E_{1}, \ldots, E_{5} \text {, and } E_{6} .
$$

the set of all $e \in R \cup S$ with the property that $e$ is incident

| with | $V_{11}$ and $V_{21}$, |
| :--- | :--- | :--- |
| with | $V_{12}$ and $V_{22}$, |
| with | $V_{11}$ and $V_{12}$, |
| with | $V_{21}$ and $V_{22}$, |
| with | $V_{11}$ and $V_{22}$, |
| with | $V_{12}$ and $V_{21}$, |

respectively.
It is clear the sets $E_{1}, \ldots, E_{5}, E_{6}$ are mutually disjoint, $R=E_{1} \cup E_{2} \cup E_{5} \cup E_{6}$ and $S=E_{3} \cup E_{4} \cup E_{5} \cup E_{6}$. Since $E(F) \cap S \neq \emptyset$, we have $V_{11} \neq \emptyset \neq V_{12}$. Since $V(H) \neq \emptyset$, we have that either $V_{21} \neq \emptyset$ or $V_{22} \neq \emptyset$. Without loss of generality we assume that $V_{21} \neq \emptyset$. We distinguish two subcases:
(1) $V_{22}=\emptyset$. Then $E_{2}=E_{4}=E_{5}=\emptyset$. Therefore $S=E_{3} \cup E_{6}$. This implies that $h_{G}\left(V_{12}\right) \leqq \varkappa_{1}(G) \leqq h_{G}(U)$, which is a contradiction.
(2) $V_{22} \neq \emptyset$. Then both $E_{1} \cup E_{4} \cup E_{6}$ and $E_{2} \cup E_{4} \cup E_{5}$ are cut sets of $G$. Therefore, $\left|E_{1} \cup E_{4} \cup E_{6}\right| \geqq \varkappa_{1}(G)$ and $\left|E_{2} \cup E_{4} \cup E_{5}\right| \geqq \varkappa_{1}(G)$. Clearly, $E_{1} \cup$ $\cup E_{3} \cup E_{5}$ and $E_{2} \cup E_{3} \cup E_{6}$ are also cut sets of $G$. Since $h_{G}\left(V_{11}\right)>h_{G}(U)$ and $h_{G}\left(V_{12}\right)>h_{G}(U)$, we have $\left|E_{1} \cup E_{3} \cup E_{5}\right|>h_{G}(U)$ and $\left|E_{2} \cup E_{3} \cup E_{6}\right|>h_{G}(U)$. Thus $2|R|+2|S|=2 h_{G}(U)+2 \varkappa_{1}(G)<\left|E_{1} \cup E_{3} \cup E_{5}\right|+\left|E_{2} \cup E_{3} \cup E_{6}\right|+\mid E_{1} \cup$ $\cup E_{4} \cup E_{6}\left|+\left|E_{2} \cup E_{4} \cup E_{5}\right| \leqq 2\right| R|+2| S \mid$, which is a contradiction.

Hence the proof is complete.
Proofs of the following corollaries are omitted:

Corollary 1. Let $G$ be a nontrivial connected graph. Then $\varkappa_{1}(G) \leqq \delta(G) \leqq$ $\leqq \max \left(f_{G}, x_{1}(G)\right)$.

Corollary 2. Let $G$ be a nontrivial connected graph, and let $U$ be a nonempty subset of $V(G)$. If $f_{G}(U) \leqq h_{G}(U)$, then $\operatorname{deg}_{G} U=h_{G}(U)$.

Corollary 3. Let $G$ be a nontrivial connected graph, and let $n$ be a positive integer such that $n \geqq x_{1}(G)$. Then there exists a vertex $u$ of $G$ such that $\operatorname{deg}_{G} u=n$ if and only if there exists a nonempty subset $U$ of $V(G)$ such that $f_{G}(U) \leqq h_{G}(U)=n$.

Corollary 4. Let $G$ be a nontrivial connected graph. Then $\delta(G)=x_{1}(G)$ if and only if there exists a nonempty subset $U$ of $V(G)$ such that $f_{G}(U) \leqq h_{G}(U)=\varkappa_{1}(G)$.

Corollary 5. (D. R. Lick [3]). Let $G$ be a nontrivial connected graph such that for each $e \in E(G), \chi_{1}(G-e)=\chi_{1}(G)-1$. Then $\delta(G)=\chi_{1}(G)$.

Note that an upper bound for the minimum degree of a graph different from the upper bound in Corollary 1 was obtained by the author in [4].

Added in proof. The graphs $G$ fulfilling the assumption of Corollary 5 were also studied by W. Mader (Minimale n-fach kantenzusammenhängende Graphen. Math. Ann. 191 (1971), 21-28).

## References

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