Juraj Bosák; Pál Erdös; Alexander Rosa Decompositions of Complete Graphs into Factors with Diameter Two

Matematický časopis, Vol. 21 (1971), No. 1, 14--28

Persistent URL: http://dml.cz/dmlcz/126805

Terms of use:

© Mathematical Institute of the Slovak Academy of Sciences, 1971

Institute of Mathematics of the Academy of Sciences of the Czech Republic provides access to digitized documents strictly for personal use. Each copy of any part of this document must contain these *Terms of use*.



This paper has been digitized, optimized for electronic delivery and stamped with digital signature within the project *DML-CZ: The Czech Digital Mathematics Library* http://project.dml.cz

DECOMPOSITIONS OF COMPLETE GRAPHS INTO FACTORS WITH DIAMETER TWO

JURAJ BOSÁK, Bratislava, PÁL ERDÖS, Budapest (Hungary) and ALEXANDER ROSA, Hamilton (Canada)

In the present paper the question is studied from three points of view whether to any natural number $k \ge 2$ there exists a complete graph decomposable into k factors with diameters two. The affirmative answer to this question is given and some estimations for the minimal possible number of vertices of such a complete graph are deduced. As a corollary it follows that given k diameters d_1, d_2, \ldots, d_k (where $k \ge 3$ and $d_i \ge 2$ for $i = 1, 2, 3, \ldots, k$), there always exists a finite complete graph decomposable into k factors with diameters d_1, d_2, \ldots, d_k . Thus Problem 1 from [1] is solved.

1

In this paper we deal only with nonoriented graphs. By a *factor* of a graph G we mean any subgraph of G containing all the vertices of G. By a *diameter* of G we understand the supremum of the set of all distances between the pairs of vertices of G (e.g. a disconnected graph has the diameter ∞). The symbol $\langle n \rangle$ denotes the complete graph with n vertices.

Let k be a natural number. By a decomposition of a graph G into k factors we mean a finite system $\{\varphi_1, \varphi_2, \ldots, \varphi_k\}$ of factors of G such that every edge of G belongs to exactly one of the factors $\varphi_1, \varphi_2, \ldots, \varphi_k$. The symbol $F_k(d_1, d_2, \ldots, d_k)$ denotes the smallest natural number n such that the complete graph $\langle n \rangle$ can be decomposed into k factors with diameters d_1, d_2, \ldots, d_k ; if such an n does not exists, we put $F_k(d_1, d_2, \ldots, d_k) = \infty$. Further, put $f_k(d) = F_k(d, d, \ldots, d)$. The main aim of the present paper is to find estimations for $f_k(2)$. From [1] it follows that $f_2(2) = 5$, $12 \leq f_3(2) \leq 13$.

Theorem 1. For any integer $k \ge 3$ we have:

$$4k-1 \leqslant f_k(2) \leqslant \binom{6k-7}{2k-2}.$$

Proof. To prove the upper estimation it suffices to decompose the graph

$$G = \left\langle \begin{pmatrix} 6k - 7\\ 2k - 2 \end{pmatrix} \right
angle$$

into k factors with diameters two. The vertices of G can be represented by (2k-2)-tuples formed from elements $1, 2, 3, \ldots, 6k-7$. The *i*th factor $(i \quad 1, 2, \ldots, k)$ consists of all edges joining (2k-2)-tuples with just i-1 common elements. The remaining edges can be added to any factor. It is easy to prove that all the factors have diameter two.

Suppose that for some $k \ge 4$ we have $f_k(2) \le 4k - 2$. Then, according to Theorem 1 of [1], $\langle 4k - 2 \rangle$ is decomposable into k factors $\varphi_1, \varphi_2, \ldots, \varphi_k$ with diameter two. Put n = 4k - 2. None of the factors φ_i $(i = 1, 2, \ldots, k)$ may have a vertex of degree n - 1 (otherwise the other factors are not con-1 ected), therefore, by [4], φ_i has at least 2n - 5 edges. The number of all edges of n is

$$\binom{n}{2} \geqslant k(2n-5),$$

whence it follows that

(1)

$$n^2 + 10k \ge 4kn + n.$$

But

$$n^2 + 10k = 16k^2 - 6k + 4,$$

 $4kn + n = 16k^2 - 4k - 2,$

thus for $k \ge 4$ we have $n^2 + 10k < 4kn + n$, which contradicts (1). For k = 3 our assertion follows from [1], Theorem 7.

Remark. The upper estimation given in Theorem 1 is too high. Therefore we later present some methods enabling to improve it, namely for a "small" kin the second part of this article, and for a "great" k in the third part.

Lemma 1. Let $k \ge 2$, $2 - d_1 \le d_2 \le d_3 \le \ldots \le d_k < \infty$. We have: $F_k(d_1, d_2, \ldots, d_k) \le f_k(2) + d_1 + d_2 + \ldots + d_k - 2k$.

Proof. From Theorem 1 it follows that $f_k(2)$ is a natural number. If d_1

 $d_2 \quad \ldots = d_k = 2$, the assertion of the lemma is evident. Thus we can suppose that there exists an integer $i \ (1 \le i \le k-1)$ such that $d_1 = d_2$

 \dots d_i $2 < d_{i+1} \leq \dots \leq d_k$. Let us construct a decomposition of the graph

 $G = \langle f_k(2) + d_1 + d_2 + \ldots + d_k - 2k \rangle$

into k factors with diameters d_1, d_2, \ldots, d_k .

The vertex set of G consists (as we may suppose) of vertices $u_1, u_2, u_3, \ldots, u_{j_k(2)}$ and of vertices $v_{j,1}, v_{j,2}, v_{j,3}, \ldots, v_{j,d_j-2}$ $(i+1 \leq j \leq k)$. Obviously, the total number of vertices is $f_k(2) + d_1 + d_2 + \ldots + d_k - 2k$. The complete subgraph of G generated by the vertices $u_1, u_2, u_3, \ldots, u_{j_k(2)}$ according to the definition of $f_k(2)$ can be decomposed into k factors $\varphi_1, \varphi_2, \ldots, \varphi_k$ with diameter two. Define a decomposition of G into factors φ'_m (m = 1, 2, ..., k) thus: Into φ'_m there belong (i) all the edges of φ_m ; (ii) all the edges $u_s v_{j,t}$ $(1 < s \leq f_k(2), i+1 \leq j \leq k, 1 \leq t \leq d_j - 2)$ such that the edge $u_s u_1$ belongs to φ_m and $j \neq m$; (iii) all the edges of the path $u_1 v_{m,1} v_{m,2} \ldots v_{m,d_{m-2}}$ (if $m \geq i + 1$). All the remaining edges are placed into φ'_1 .

It is easy to show that q'_m has diameter d_m (m = 1, 2, ..., k). The lemma follows.

Lemma 2. Let
$$k \ge 3$$
, $2 \le d_1 \le d_2 \le ... \le d_k < \infty$. Then we have:
 $F_k(d_1, d_2, ..., d_k) \le \binom{6k - 7}{2k - 2} + d_1 + d_2 + ... + d_k - 2k.$

Proof. Distinguish two cases:

I. $d_1 = 2$. Then the assertion follows from Lemma 1 and Theorem 1.

II. $d_1 > 2$. By [1], Theorem 4, we have:

$$F_k(d_1, d_2, ..., d_k) \leq d_1 + d_2 + ... + d_k - k_k$$

Since for any $k \ge 2$ we have

$$k \leqslant \binom{6k-7}{2k-2}$$

the lemma follows.

Corollary. Let $k \ge 3$, $2 \le d_1 \le d_2 \le \ldots \le d_k \le \infty$. Then $F_k(d_1, d_2, \ldots, d_k)$ is a natural number.

Proof. If $d_k < \infty$, our assertion follows from Lemma 2. If $d_2 = \infty$, the assertion follows from [1], Theorem 3. Therefore we may suppose that $d_2 < \infty$, $d_k = \infty$, i. e. there is an integer i $(2 \le i \le k-1)$ such that $2 \le d_1 \le d_2 \le \le \ldots \le d_i < \infty = d_{i+1} = d_{i-2} = \ldots = d_k$.

If $i \ge 3$, according to Lemma 2, $F_i(d_1, d_2, ..., d_i)$ is a natural number. Therefore the finite complete graph

$$G = \langle F_i(d_1, d_2, \ldots, d_i) \rangle$$

is decomposable into *i* factors with diameters d_1, d_2, \ldots, d_i . If we add k - i null factors (i. e., factors without edges), we obtain a decomposition of G into k factors with diameters $d_1, d_2, \ldots, d_i, d_{i-1}, \ldots, d_k$.

If i = 2, then according to Theorem 8 of [1] $F_3(d_1, d_2, d_3 = \infty)$ is a natural number. Since

$$F_k(d_1, d_2, d_3 = \infty, ..., d_k = \infty) \leqslant F_3(d_1, d_2, d_3 = \infty),$$

then $F_k(d_1, d_2, \ldots, d_k)$ is also a natural number. The corollary follows.

Remark. As the supposition $d_1 \leq d_2 \leq \ldots \leq d_k$ is not essential, the preceding corollary completely solves Problem 1 from [1], p. 53.

 $\mathbf{2}$

Let a natural number n and a set $A \subseteq \{1, 2, ..., n\}$ be given. A is called an S_n -set if each $x \in \{1, 2, ..., n\}$, $x \notin A$ can be written in at least one of the following forms

$$egin{array}{ll} x &= a + b, \ x &= a - b, \ x &= 2n + 1 - (a + b), \end{array}$$

where $a, b \in A$.

Let k be a natural number. Denote by g(k) the least natural number l such that the set $\{1, 2, ..., l\}$ can be partitioned into k disjoint S_l -sets. (If such ι natural number l does not exist, put $g(k) = \infty$.)

Lemma 3. $f_k(2) \leq 2g(k) + 1$ for any integer $k \geq 2$.

Proof. Let natural numbers m and n be given. We shall call a finite graph (without loops or multiple edges) with m labelled vertices v_1, v_2, \ldots, v_m cyclic, if it contains with each edge $v_i v_j$ $(i, j \in \{1, 2, \ldots, m\})$ the edge $v_i \ v_{j+1}$ (the indices taken modulo m) as well. By the *length* of an edge $v_i v_j$ we mean the number

$$\min\{|i-j|, m-|i-j|\}$$

Evidently, a cyclic graph contains either every or no edge of length i for each $i \in \{1, 2, ..., [m/2]\}$.

Assign to a given S_n -set A a cyclic graph with 2n + 1 vertices containing edges of length i if and only if $i \in A$ (i = 1, 2, ..., n). It is clear that thus a one-to-one correspondence between cyclic graphs with 2n + 1 labelled vertices with diameter two and S_n -sets is defined. Further, it is obvious that to different [disjoint] S_n -sets different [edge-disjoint, respectively] cyclic factors with diameter two of $\langle 2n + 1 \rangle$ are assigned. Therefore the assertion of the lemma follows immediately from the definitions of $f_k(2)$ and g(k).

Let natural numbers n, i, integers c, d and a set $A \subseteq \{1, 2, ..., n\}$ be given. Denote by $\operatorname{red}_{n}c$ the (uniquely determined) integer r such that

$$r \equiv c \pmod{2n+1},$$
$$|r| \leq n.$$

Further, put

$$egin{aligned} r^{(t)} &= |\mathrm{red}_n r^i|\,, \ c \,\circ\, d &= |\mathrm{red}_n c d|\,, \ c \,\circ\, A &= \{c \,\circ\, d\,; \ d \in A\}\,. \end{aligned}$$

Evidently, we always have

$$\begin{array}{ll} (*) & 0 \leqslant c \circ d \leqslant n, \\ c \circ A \subseteq \{0, 1, 2, \ldots, n\}. \end{array}$$

Lemma 4. If n and r are such natural numbers that the greatest common divisor (2n + 1, r) = 1 and A is an S_n -set, then $r \circ A$ is an S_n -set as well.

Proof. Choose $x \in \{1, 2, ..., n\}$. It suffices to prove that either $x \in r$ A or there exist $a, b \in A$ such that one of the equalities

$$egin{array}{ll} x=r&a+r\circ b,\ x=r\circ a-r\circ b,\ x=(2n+1)-(r\circ a+r\circ b), \end{array}$$

holds.

.

It is easy to see that there is a $y \in \{1, 2, ..., n\}$ such that $r \circ y = x$. In fact as (r, 2n + 1) = 1, the congruence

$$rz \equiv x \pmod{2n+1}$$

has a solution $z \in \{1, 2, ..., 2n\}$. If $1 \leq z \leq n$, we put y = z, and if $n + 1 \leq z \leq 2n$, we put y = 2n + 1 - z.

Since A is an S_n -set, either $y \in A$ or there exist $a, b \in A$ such that one of the following cases occurs:

$$y = a - b$$
,
 $y = a + b$,
 $y = 2n + 1 - (a + b)$.

If $y \in A$, then evidently $x = r \circ y \in r \circ A$. Let us analyze the other cases (all the following congruences are related to the modul 2n + 1).

(I) y = a - b. Obviously $\pm r \circ y \equiv ry - ra - rb$, where $ra \equiv \pm r - a$, $rb \equiv \pm r \circ b$.

By examining all 8 possibilities for choice of signs we find that one of the following 4 cases occurs (we use inequality (*)):

$$\begin{array}{ll} x - r \circ y \equiv & r \circ a + r \circ b, \text{ hence } x = r \circ a + r \circ b, \\ x & r \circ y \equiv & r \circ a - r \circ b, \text{ hence } x = r \circ a - r \circ b, \\ x & r \circ y \equiv -r \circ a + r \circ b, \text{ hence } x = r \circ b - r \circ a, \\ x & r \circ y \equiv -r \circ a - r \circ b \equiv (2n + 1) - r \circ a - r \circ b, \\ \text{ so } x = 2n + 1 - (r \circ a + r \circ b). \end{array}$$

(II) y = a + b. Evidently

$$\pm k \circ y \equiv ky = ka + kb \equiv \pm k \circ a \pm k \circ b,$$

where we again have 8 possibilities for choice of the signs. Further procedure is the same as in case (I).

(III) y = 2n + 1 - (a + b). We have: $\pm k \circ y \equiv ky = k(2n + 1) - ka$ $kb = -ka - kb \equiv \pm k \circ a \pm k \circ b$. Further we proceed as in case (I). The lemma follows.

Lemma 5. Let r, n and k be such natural numbers that

- (1) 2n = 1 is a prime number,
- (2) k divides n,
- (3) r is a primitive root of 2n + 1, (1)
- (4) $A = \{r^{(k)}, r^{(2k)}, r^{(3k)}, \dots, r^{(n)} 1\}$ is an S_n -set.

Then $g(k) \leq n$.

Proof. From (1) and (3) it follows that (r, 2n + 1) = 1 and that the numbers $r, r^2, \ldots, r^n, \ldots, r^{2n}$ represent all non-zero residue classes modulo 2n + 1. From this fact it can be easily deduced that $\{r^{(1)}, r^{(2)}, \ldots, r^{(n)}\} = \{1, 2, \ldots, n\}$. From (2) and (4) it follows that the sets $A, r \circ A, r^2 \circ A, \ldots, r^{k-1} \circ A$ are mutually disjoint. They are S_n -sets, as it follows from (4) and Lemma 4. Therefore the set $\{1, 2, \ldots, n\}$ can be decomposed into k disjoint S_n -sets, consequently $g(k) \leq n$.

Lemma 6. We have: $g(1) \leq 1$, $g(2) \leq 2$, $g(3) \leq 6$, $g(4) \leq 20$, $g(5) \leq 35$, $g(6) \leq 78$, $g(7) \leq 98$, $g(8) \leq 96$, $g(9) \leq 189$, $g(10) \leq 260$.

Proof. We use the method from Lemma 5: we look for such a multiple n of k that (1) is valid and the least primitive root r of 2n + 1 satisfies (4). With the help of tables of the least primitive roots of primes (see, e. g. [5]) we can construct the following S_n -sets A:

(1) A natural number r is called a *primitive root* of a prime number p if the numbers r, r^2 , r^3 , ..., $r^{p-1} \equiv 1$ represent all non-zero residue classes modulo p.

 $\begin{array}{l} k=1,\ n=\ 1,\ r=2,\ A=\{1\}.\\ k=2,\ n=\ 2,\ r=2,\ A=\{1\}.\\ k=3,\ n=\ 6,\ r=2,\ A=\{1,5\}.\\ k=4,\ n=20,\ r-3,\ A=\{1,4,10,16,18\}.\\ k=5,\ n=35,\ r=7,\ A=\{1,20,23,26,30,32,34\}.\\ k=6,\ n=78,\ r=5,\ A=\{1,4,14,16,27,39,46,49,56,58,64,67,75\}.\\ k=7,\ n=98,\ r=2,\ A=\{1,6,14,19,20,33,36,68,69,77,83,84,87,93\}.\\ k=8,\ n=96,\ r=5,\ A=\{1,7,9,12,16,43,49,55,63,81,84,85\}. \end{array}$

 $k = 9, n = 189, r = 2, A = \{1, 5, 25, 39, 51, 52, 57, 68, 76, 86, 91, 93, 94, 119, 124, 125, 133, 138, 162, 163, 184\}.$

 $k = 10, n = 260, r = 3, A = \{1, 10, 18, 29, 32, 42, 52, 55, 62, 74, 98, 99, 100, 101, 106, 114, 176, 180, 197, 201, 219, 226, 231, 235, 237, 255\}.$

To check that they are S_n -sets is a matter of routine. The rest of the proof follows from Lemma 5.

Remark. It can be easily found that even g(1) = 1, g(2) = 2, g(3) = 6. By a systematic examination we can also establish that g(4) = 20, but, on the other hand, g(5) = 30. (The inequality $g(5) \leq 30$ follows from the fact that $A = \{1, 5, 6, 11, 14, 29\}$, $3 \circ A$, $3^2 \circ A$, $3^3 \circ A$ and $3^4 = A$ are disjoint S_{30} -sets.)

Theorem 2. We have: $f_2(2) \leq 5, f_3(2) \leq 13, f_4(2) \leq 41, f_5(2) \leq 61, f_6(2) \leq 157, f_7(2) \leq 193, f_8(2) \leq 193, f_9(2) \leq 379, f_{10}(2) \leq 521.$

Proof. For $k \neq 5$, $k \neq 7$ the upper estimation of $f_k(2)$ follows from Lemmas 3 and 6. For k = 5 it suffices to apply Lemma 3 and the preceding remark. For k = 7 we proceed thus: Evidently $f_7(2) \leq f_8(2)$, because from a decomposition of a complete graph into 8 factors with diameter two we obtain a decomposition into 7 factors with diameter two by unifying edges of any two of the 8 given factors leaving the other 6 factors without any change. Since $f_8(2) \leq 193$, we have $f_7(2) \leq 193$ as well.

Lemma 7. There exists a natural number N such that for all naturals n > Nwe have: The number A_n of all factors of $\langle n$ with t [] $3n^3 \log n$] edges and with a diameter greater than two is less than

$$\frac{1}{n} \begin{pmatrix} \binom{n}{2} \\ t \end{pmatrix}.$$

Proof uses methods similar to those used in [2].

³

(I) Pick a vertex x of $\langle n \rangle$. Let i be an integer for which

$$0 \leq i \leq t$$

holds. Denote by a_i the number of factors of $\langle n \rangle$ with t edges, in which the degree of x is *i*. Evidently, we have:

$$a_i = {n-1 \choose i} \left({n-1 \choose 2 \choose t-i}
ight),$$

(II) Put $l [] 3n \log n$]. Prove that there is a number N_1 such that for i = 0, 1, 2, ..., l and for every natural $n > N_1$ we have

$$rac{a_i}{a_{2l}} < rac{1}{n^3}$$

It is easy to see that for any natural n the inequalities

$$nl \leqslant t, \\ 2l \leqslant t$$

are valid. Now, we have:

$$\begin{aligned} \frac{a_i}{a_{2l}} &= \frac{\binom{n-1}{i}\binom{\binom{n-1}{2}}{t-i}}{\binom{n-1}{2}} = \\ &= \frac{\binom{n-1}{2l}\binom{\binom{n-1}{2}}{t-i}}{\binom{n-1}{2l}\binom{\binom{n-1}{2}}{t-2l}} = \\ &= \frac{(i+1)(i+2)\dots 2l}{(n-i-1)(n-i-2)\dots(n-2l)} \times \\ &\times \frac{\left(\binom{n-1}{2}-t+2l\right)\left(\binom{n-1}{2}-t+2l-1\right)\dots\left(\binom{n-1}{2}-t+i+1\right)}{(t-2l+1)(t-2l+2)\dots(t-i)} < \\ &\leq \frac{(i+1)(i+2)\dots 2l}{(n-i-1)(n-i-2)\dots(n-2l)} \cdot \frac{\binom{n^2}{2}^{2l-i}}{(t-2l+1)(t-2l+2)\dots(t-i)} = \\ &= \frac{(i+1)(i+2)\dots 2l}{2^{2l-i}} \cdot \binom{n}{t}^{2l-i} \cdot \frac{n^{2l-i}}{(n-i-1)(n-i-2)\dots(n-2l)} \times \end{aligned}$$

$$\times \frac{t^{2l-i}}{(t-2l+1)(t-2l+2)\dots(t-i)} \leqslant \frac{(i+1)(i+2)\dots 2l}{(2l)^{2l-i}} \times \\ \times \left(\frac{n}{n-2l}\right)^{2l-i} \cdot \left(\frac{t}{t-2l+1}\right)^{2l-i} \leqslant \frac{l+1}{2l} \cdot \frac{l+2}{2l} \dots \frac{2l}{2l} \cdot \left(\frac{n}{n-2l}\right)^{2l} \times \\ \times \left(\frac{t}{t-2l+1}\right)^{2l} \leqslant \left(\frac{3}{4}\right)^{l-1} \cdot \left(\frac{1}{2} / \frac{5}{4}\right)^{2l} \cdot \left(\frac{1}{2} / \frac{5}{4}\right)^{2l} = \frac{5}{4} \cdot \left(\frac{15}{16}\right)^{l-1} < \\ < \frac{5}{4} \left(\frac{15}{16}\right)^{\sqrt{n}} < \frac{1}{n^3}$$

for every natural $n > N_1$, if N_1 is a sufficiently large constant.

(III) Let us prove that the number $B_n(x)$ of the factors of $\langle n \rangle$ with t edges, in which the degree of x does not exceed l, is less than

$$\frac{1}{2} \frac{\binom{\binom{n}{2}}{t}}{n^2}$$

for every sufficiently large n.

Obviously, according to (II) for $n > N_1$ we have:

$$\begin{aligned} \frac{n^2 B_n(x)}{\binom{n}{2}} &= n^2 \frac{a_0 + a_1 + \dots + a_l}{\binom{n}{2}} \\ &\leq n^2 \frac{a_0 + a_1 + \dots + a_l}{a_{2l}} = n^2 \left(\frac{a_0}{a_{2l}} + \frac{a_1}{a_{2l}} + \dots + \frac{a_l}{a_{2l}}\right) < \\ &< n^2(l+1)\frac{1}{n^3} = \frac{\lfloor \sqrt{3n \log n} \rfloor + 1}{n}. \end{aligned}$$

Evidently, the last expression tends to zero for $n \rightarrow \infty$. Therefore

$$\frac{\left[\sqrt{3n\log n}\right]+1}{n} < \frac{1}{2}$$

 $\mathbf{22}$

for $n > N_2$, where N_2 is a sufficiently large constant so that

$$rac{n^2B_n(x)}{\left(inom{n}{2}
ight)\atopt}<rac{1}{2},$$

1. e.

$$B_n(x) < rac{1}{2} rac{igg(inom{n}{2}igg)}{n^2}$$

for $n > \max\{N_1, N_2\}$.

(IV) We prove now that the number B_n of the factors of $\langle n \rangle$ with t edges containing a vertex of degree $\leq l$, is less than

$$\frac{1}{2n} \binom{\binom{n}{2}}{t}$$

for $n > \max\{N_1, N_2\}$.

Evidently, we have

$$B_n \leqslant \sum_x B_n(x),$$

where x runs through the vertex set of $\langle n \rangle$. Therefore, using (III) we obtain

$$B_n \leq \sum_x B_n(x) < n - \frac{1}{2} \frac{1}{n^2} \left(\binom{n}{2} \\ t \right) = \frac{1}{2n} \left(\binom{n}{2} \\ t \right)$$

for $n > \max\{N_1, N_2\}$.

(V) Fix now two different vertices x and y of $\langle n \rangle$ and two integers i and j satisfying the relations l < i < n, l < j < n.

Denote by $D_n(x, y, i, j)$ the number of factors of $\langle n \rangle$ with t edges in which x has degree i, y has degree j, and x is not joined with y by an edge. We have:

$$D_n(x, y, i, j) = \binom{n-2}{i} \binom{n-2}{j} \binom{\binom{n-2}{2}}{t-i-j}.$$

Further, denote by $E_n(x, y, i, j)$ the number of factors of $\langle n \rangle$ with t edges in which x has degree i, y has degree j, and the distance of x and y is greater than two. Evidently,

$$E_n(x, y, i, j) = \binom{n-2}{i} \binom{n-2-i}{j} \binom{\binom{n-2}{2}}{t-i-j}.$$

We shall find a natural number N_3 such that for every $n > N_3$ we have

$$\frac{E_n(x, y, i, j)}{D_n(x, y, i, j)} < \frac{1}{n^3}.$$

Obviously, we have:

$$\frac{E_n(x, y, i, j)}{D_n(x, y, i, j)} - \frac{n - i - 2}{n - 2} \cdot \frac{n - i - 3}{n - 3} \dots \frac{n - i - j - 1}{n - j - 1} < \left(\frac{n - i - 2}{n - 2}\right)^j \leqslant \left(\frac{n - 3 - l}{n - 2}\right)^{l-1}.$$

It is easy to see that there exists a natural number N_3 such that for all $n > N_3$ we have

$$\frac{n-2}{l+1} > 1.$$

Evidently, it suffices to prove that for every $n > N_3$ we have:

$$\left(\frac{n-2}{n-3-l}\right)^{l+1} > n^3.$$

But for $n > N_3$ we have:

$$\left(\frac{1+\frac{1}{n-2}}{l+1}-1\right)^{n-2} > e.$$

It follows that

$$\left(\frac{n-2}{n-3-l}\right)^{l+1} = \left(\left(1 + \frac{1}{\frac{n-2}{l+1} - 1}\right)^{n-2} \right)^{l+1} \frac{(l+1)^2}{n-2} > e^{\frac{(l+1)^2}{n-2}} > e^{\frac{(l+1)^2}{n-2}} > e^{\frac{(l+1)^2}{n-2}} = e^{\frac{(l+1)^2}{n-2}} - n^3.$$

24

(V1) Let C_n be the number of factors of $\langle n \rangle$ with t edges in which all the vertices have degrees greater than l and with diameters greater than two. From (V) it follows that for every $n > N_3$ we have:

$$C_{n} \leq \sum_{(x,y)} \sum_{(i,j)} E_{n}(x, y, i, j) \leq < < \sum_{(x,y)} \sum_{(i,j)} \frac{D_{n}(x, y, i, j)}{n^{3}} = \frac{1}{n^{3}} \sum_{(x,y)} \sum_{(i,j)} D_{n}(x, y, i, j) < < \sum_{(x,y)} \frac{\binom{n}{2}}{t} = \binom{n}{2} \cdot \binom{n}{2} \cdot \binom{n}{2} \cdot \binom{n}{2} = \binom{n}{2} \cdot \binom{$$

where (x, y) runs through the set of all unordered pairs of different vertices of n (i, j) runs through the set of all ordered pairs of integers such that $l < i < n, \ l < j < n.$

(VII) Put $N = \max \{N_1, N_2, N_3\}$. Then, according to (IV) and (VI) forevery natural number n > N we have:

$$A_n \leqslant B_n + C_n < egin{pmatrix} \binom{\binom{n}{2}}{t} & \binom{\binom{n}{2}}{t} & = rac{\binom{\binom{n}{2}}{t}}{t} \\ 2n & + rac{2n}{2n} & = rac{\binom{\binom{n}{2}}{t}}{n}. \end{cases}$$

The lemma follows.

Lemma 8. A natural number M exists such that for every integer n > M we have: n contains

$$\left[\left. \right| \right/ \frac{n-2}{12\log n} \right]$$

edge-disjoint factors with diameter two.

Proof. According to Lemma 7 there exists a positive integer N such that for every integer n > N we have:

$$A_n < \frac{1}{n} \begin{pmatrix} \binom{n}{2} \\ t \end{pmatrix}.$$

25

 \mathbf{Put}

$$u = \left[\frac{\binom{n}{2}}{t}\right].$$

Evidently there is a natural number N_4 such that for every $n > N_4$ we have u < n. Put $M = \max\{N, N_4\}$. Obviously for $n \ge 2$ we have:

$$u = \left[\frac{n(n-1)}{2\left[\sqrt[]{3n^3\log n}\right]}\right] \ge \left[\frac{n(n-1)}{2\sqrt[]{3n^3\log n}}\right] =$$
$$= \left[\sqrt[]{\frac{n^2 - 2n + 1}{12n\log n}}\right] \ge \left[\sqrt[]{\frac{n^2 - 2n}{12n\log n}}\right] = \left[\sqrt[]{\frac{n-2}{12\log n}}\right].$$

. Therefore it suffices to prove that for n > M the graph $\langle n \rangle$ contains u edgedisjoint factors with diameter two.

If we assume the contrary, then each of the

$$p = rac{\displaystyle \prod_{i=0}^{u-1} \left(inom{n}{2} - it \ t
ight)}{u!}$$

systems S consisting of u edge-disjoint factors of $\langle n \rangle$, each with t edges, contains at least one factor with diameter greater than two. Any such factor with t edges and with diameter greater than two occurs just in

$$q = \frac{\prod_{i=1}^{u-1} \left(\binom{n}{2} - it \right)}{(u-1)!}$$

systems S. Therefore the number of factors of $\langle n \rangle$ with t edges and with a diameter greater than two is at least

$$rac{p}{q} = rac{1}{u} egin{pmatrix} n \ 2 \ t \end{pmatrix} > rac{1}{n} egin{pmatrix} n \ 2 \ t \end{pmatrix},$$

which contradicts Lemma 7. Thus Lemma 8 follows.

Theorem 3. There exists a positive integer K such that for any integer k > K we have:

26

$$f_k(2) \leqslant \left(\frac{49}{10}\right)^2 k^2 \log k$$

Proof. Pick a natural number K_1 such that for every $k > K_1$ we have

$$\left[\left(rac{49}{10}
ight)^2 \ k^2 \log k
ight] > M$$
 ,

where M is the constant from Lemma 8.

Pick a natural number K_2 in such a way that for any $k > K_2$

 $k^2 \log k \ge 750,$

and, consequently,

$$-3 \ge -\frac{1}{250} k^2 \log k.$$

Further, pick a natural number K_3 such that for every integer $k > K_3$ we have:

$$\left(\frac{49}{10}\right)^2 \log k \leqslant k^{\frac{1}{2000}}.$$

Put $K = \max \{K_1, K_2, K_3\}$. Pick an integer k > K. Put

$$n = \left[\left(\frac{49}{10} \right)^2 k^2 \log k \right].$$

Then we have:

$$\frac{n-2}{\log n} \ge \frac{\left(\left(\frac{49}{10}\right)^2 k^2 \log k - 1\right) - 2}{\log\left(\left(\frac{49}{10}\right)^2 k^2 \log k\right)} = \frac{\left(\frac{49}{10}\right)^2 k^2 \log k - 3}{2 \log k + \log\left(\left(\frac{49}{10}\right)^2 \log k\right)} \ge \frac{\left(\frac{49}{10}\right)^2 k^2 \log k - \frac{1}{250} k^2 \log k}{2 \log k + \log\left(k^{\frac{1}{2000}}\right)} = 12k^2.$$

It follows that

$$k \leqslant \sqrt{\frac{n-2}{12\log n}}$$
,

where n > M. From Lemma 8 it follows that $\langle n \rangle$ can be decomposed into k edge-disjoint factors with diameter two (the remaining edges may be added to any factor). Consequently,

$$f_k(2) \leqslant n \leqslant \left(\frac{49}{10}\right)^2 k^2 \log k.$$

The theorem follows.

Remark. It can be proved that there exist positive constants C_1 and C_2 such that

$$C_1 k^2 < g(k) < C_2 k^2 \log k$$

for every sufficiently large k; the left inequality is obvious; the right one can be obtained using similar methods as in our Theorem 3 and in [3]; this remains true even if we do not allow representations of the form 2n + 1 - (a + b). Now, using Lemma 3 we can again obtain that $f_2(k) < Ck^2 \log k$ for certain constant C and all sufficiently large k.

Problem 1. Is $g(k)/k^2$ bounded?

Problem 2. Determine $\lim_{k\to\infty} \frac{f_k(2)}{k}$.

REFERENCES

- Bosák J., Rosa A., Znám Š., On decompositions of complete graphs into factors with given diameters, Theory of Graphs, Proc. Colloq. Tihany 1966, Akadémiai Kiadó, Budapest 1968, 37-56.
- [2] Erdös P., Rényi A., On the evolution of random graphs, Magyar tud. akad. Mat. kutató int. közl. 5 (1960), 17 61.
- [3] Erdös P., Rényi A., Additive properties of random sequences of positive integers, Acta Arithm. 6 (1960), 83-110.
- [4] Erdös P., Rényi A., Egy gráfelméleti problémáról, Magyar tud. akad. Mat. kutató int. közl. 7 (1962), 623-641.
- [5] Виноградов И. М., Основы теории чисел, ГИТТЛ, Москва-Ленинград 1952.

A magyar tudományos akadémia Matematikai kutató intézet, Budapest

Matematický ústav Slovenskej akadémie vied, Bratislava

McMaster University Department of Mathematics, Hamilton