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# **ON SOME PROPERTIES OF TRANSFORMATIONS OF A LOGIC**

#### ANATOLIJ DVUREČENSKIJ

In the paper the notion of the ergodicity on a logic will be introduced and the different types of the transformations of a logic will be characterized and the recurrence theorems will be proved.

## 1. Ergodic properties of homomorphisms

Let L be a  $\sigma$ -lattice with the first and the last elements O and 1, respectively and an orthocomplementation  $\bot : a \mapsto a^{\perp}$ , which satisfies (i)  $(a^{\perp})^{\perp} = a$  for all  $a \in L$ ; (ii) if a < b, then  $b^{\perp} < a^{\perp}$  for all  $a, b \in L$ ; (iii)  $a \vee a^{\perp} = 1$  for all  $a \in L$ . We say that a, b are orthogonal and write  $a \perp b$  if  $a < b^{\perp}$ . We further assume that if  $a, b \in L$  and a < b, then there exists an element  $c \in L$  such that  $a \perp c$ and  $a \vee c = b$ . A  $\sigma$ -lattice satisfying the above axioms will be called a logic (see [1]).

A state is a map m from L into  $\langle 0, 1 \rangle$  such that m(1) = 1 and  $m(\bigvee_{i=1}^{\infty} a_i) = \sum_{i=1}^{\infty} m(a_i)$  if  $a_i \perp a_j$  for  $i \neq j$ . A logic is full in the case: (i) if  $a \neq b$ , there exists a state m such that  $m(a) \neq m(b)$ ; (ii) if  $a \neq 0$ , there exists a state m such that  $m(a) \neq m(b)$ ; (ii) if  $a \neq 0$ , there exists a state m such that m(a) = 1. An observable is a map x from the Borel sets  $B(R_1)$  of  $R_1$  into a logic L, which satisfies (i)  $x(R_1) = 1$ ; (ii)  $x(E) \perp x(F)$  if  $E \cap F = \emptyset$ ; (iii)  $x(\bigcup_{i=1}^{\infty} E_i) = \bigvee_{i=1}^{\infty} x(E_i)$  if  $E_i \cap E_j = \emptyset$ ,  $i \neq j$ ,  $E_i \in B(R_1)$ .

Let x be an observable and m be a state. Then we shall say that x is

- (i) a constant (a constant a.e. [m]) if there is a real number  $\lambda$  such that  $x(\{\lambda\}) = 1$   $(m(x(\{\lambda\}) = 1);$
- (ii) bounded (bounded a. e. [m]) if there is a compact set K with the property x(K) = 1 (m(x(K)) = 1).

We denote by  $\sigma(x)$  ( $\sigma_m(x)$ ) the smallest closed set E such that x(E) = 1 (m(x(E)) = 1).

A homomorphism of a logic L is a map T from L into L such that TO = O;

 $T(a^{\perp}) = (Ta)^{\perp}$  for all  $a \in L; T(\bigvee_{i=1}^{\infty} a_i) = \bigvee_{i=1}^{\infty} Ta_i$ . We shall say that a homo morphism T of a logic L is (i) invariant in a state m if m(Ta) = m(a) for all  $a \in L$ ; (ii) ergodic in a state m if the equality Ta = a implies  $m(a) \in \{0, 1\}$ .

Let T be a homomorphism of L and x be an observable. We shall say that x is T-invariant if T(x) = x, where  $(T(x))(E) = T(x(E)), E \in B(R_1)$ .

**Theorem 1.1.** A homomorphism T of a full logic L is ergodic in every state iff the constants are the only T-invariant observables.

Proof. For sufficiency, let the constants be the only *T*-invariant observables and let Ta = a. We define an observable  $q_a : q_a(\{0\}) = a^{\perp}, q_a(\{1\}) = a$ . It follows that  $q_a$  is *T*-invariant and hence  $q_a(\{1\}) = a$  is either 1 or *O*. Then  $m(a) \in \{0, 1\}$  for all *m*.

Conversely, let T be ergodic in every state and let T(x) = x, hence  $m(x(E)) \in \{0, 1\}$  for all m. If  $O \neq x(E) \neq 1$  for some  $E \in B(R_1)$ , then there exist two states  $m_1$ ,  $m_2$  such that  $m_1(x(E)) = 1$ ,  $m_2(x(E)^{\perp}) = 1$ . Thus if  $m = \frac{1}{2}(m_1 + \frac{1}{2})$ 

 $(+ m_2)$ , we have  $m(x(E)) = \frac{1}{2}$ . This is a contradiction and hence x(E) is either O or 1. Let us denote

$$\mathscr{C} = \{ E \in B(R_1) : E \supset \sigma(x) \text{ or } E \cap \sigma(x) = \emptyset \}.$$

If a < b, then either  $\langle a, b \rangle$  or (a, b) is in  $\mathscr{C}$ . But  $\mathscr{C}$  is a  $\sigma$ -algebra and hence it equals  $B(R_1)$ . Hence it follows that there is a  $\lambda \in R_1$  such that  $x(\{\lambda\}) = 1$ . q.e.d

**Theorem 1.2.** A homomorphism T of a logic L (L is arbitrary) is ergodic in a state m iff the constants a. e. [m] are the only T-invariant observables bounded a. e. [m].

Proof. Let T(x) = x, x be bounded a. e. [m] and let T be ergodic in m, then  $m(x(E)) \in \{0, 1\}$  for all  $E \in B(R_1)$ . If we denote  $a = \inf \sigma_m(x)$ ,  $b = \sup \sigma_m(x)$ , we shall have  $m(x(\langle a, b \rangle)) = 1$  and by application of the Weierstrass method of dividing repeatedly the bounded interval  $\langle a, b \rangle$  into halves we shall obtain a sequence  $\{\langle a_n, b_n \rangle\}$  of intervals such that  $\langle a, b \rangle \supset \langle a_1, b_1 \rangle \supset$  $\supset \langle a_2, b_2 \rangle \supset \ldots$  and  $m(x(\langle a_n, b_n \rangle)) = 1$  for  $n = 1, 2, \ldots$  Hence there is a  $\lambda \in R_1$  such that  $\{\lambda\} = \bigcap_{n=1}^{\infty} \langle a_n, b_n \rangle$  and consequently  $m(x(\{\lambda\})) = 1$ . The sufficient condition is trivial.

The sufficient condition is trivial.

**Corollary 1.2.1.** A homomorphism T of a logic L is ergodic in a state m iff the constants a. e. [m] are the only T-invariant observables (not necessarily bounded).

Proof. Only necessity. For  $\sigma_m(x)$  we have  $\sigma_m(x) = \bigcup_{n=-\infty}^{\infty} (\sigma_m(x) \cap \langle n, n + 1)) = 1$ . The set  $E = \sigma_m(x) \cap \langle n, n + 1 \rangle$  is bounded and as above there is a  $\lambda \in R_1$  such that  $m(x(\{\lambda\})) = 1$ .

Remark 1. Theorem 1.2. will be valid if the assumption of the boundedness a. e. [m] of x is omitted, provided that  $x \in O_p(m) = \{x: |\int \lambda^p m(x(d\lambda))| < \infty\}$ for  $1 \leq p < \infty$ . In fact, if Ta = a, then the observable  $q_a$  is in  $O_p(m)$  and  $\int \lambda^p m(q_a(d\lambda)) = m(a) \in \{0, 1\}$ . On the other hand, the necessity is easily seen from Corollary 1.2.1.

Remark 2. Let L be now a logic in the sense [5], that is, L is not a lattice in general. Then the Theorems 1.1., 1.2., the Corollary 1.2.1. and the Remark 1 will be valid, too.

Lemma 1.3. An automorphism T of a logic L is ergodic in a state m iff  $m(\bigvee_{j=\infty}^{\infty} T^j a) = 1$  holds for each  $a \in L$ , m(a) > 0. Proof. The sufficiency is trivial. On the other hand let m(a) > 0, then for  $b = \bigvee_{j=-\infty}^{\infty} T^j a$  we have m(b) > 0. But Tb = b and hence m(b) = 1.

q.e.d.

q.e.d.

If we use Wigner's and Gleason's theorems (see [1]) about the representation of automorphisms and the states, respectively, in the case of a logic of all closed subspaces of a Hilbert space H we shall give an interesting example which is a generalization of a known proposition in the ergodic theory (see [2] p. 34).

Let L = L(H) be the logic of all closed subspaces of H and (.,.) be the inner product on H. Since there is a one-to-one correspondence between the closed subspaces M of H and their projectors  $P^M$ , we shall write M for an element as well as for its projector. Let U be a unitary operator on H and  $\varphi$  be a unit vector in H. Then  $T_U: M \mapsto UMU^{-1}, M \in L(H)$ , is an automorphism of L(H)and  $m_{\varphi}: M \mapsto (M\varphi, \varphi), M \in L(H)$  is a state of L(H).

Example. Let U be a unitary operator on a Hilbert space H and  $P = \{\xi \in H : U\xi = \xi\} \neq 0$ . Then an automorphism  $T_U(.) = U(.)U^{-1}$ , is invariant in a state  $m_{\varphi}, \varphi \in P$ ,  $||\varphi|| = 1$ , where  $m_{\varphi}(M) = (M\varphi, \varphi), M \in L(H)$ .

If dim P = 1 then, moreover,  $T_U$  is ergodic in a state  $m_{\varphi}$ , Conversely, if for each  $\varphi \in P$ ,  $||\varphi|| = 1$ ,  $T_U$  is ergodic in a state  $m_{\varphi}$ , then dim P = 1.

Proof. For invariancy:  $m_{\varphi}(T_U M) = (UMU^{-1}\varphi, \varphi) = (MU^{-1}\varphi, U^{-1}\varphi) =$ =  $(M\varphi, \varphi) = m_{\varphi}(M)$ . Now let dim P = 1 and  $T_U M = M$ , that is  $UMU^{-1} =$ = M, UM = MU. If  $\varphi$  is a unit vector in P, then  $UM\varphi = MU\varphi = M\varphi$ , i. e.  $M\varphi \in P$  and  $M\varphi = \alpha\varphi$ . But  $\alpha^2\varphi = M^2\varphi = M\varphi = \alpha\varphi$ , hence  $\alpha \in \{0, 1\}$  and consequently  $m_{\varphi}(M) = \alpha \in \{0, 1\}$ .

Conversely, let  $T_U$  be ergodic in all  $m_{\varphi}$ ,  $\varphi \in P$ ,  $||\varphi|| = 1$  and let dim P > 1then there exist two orthonormal vectors  $\varphi_1$ ,  $\varphi_2$  in P. Hence if  $\varphi = \frac{\sqrt{2}}{2}(\varphi_1 + \varphi_2)$ 

 $+ \varphi_2$ ) and M is a subspace generated by  $\varphi_1$ , then  $\varphi \in P$ ,  $||\varphi|| = 1$  and UM = MU because if  $\xi = \alpha \varphi_1 + y$ ,  $y \perp \varphi_1$ , then  $UM\xi = \alpha \varphi_1$ ,  $MU\xi = \alpha \varphi_1 + MUy$ . But  $(\varphi_1, Uy) = (U^*\varphi_1, y) = 0$  and hence  $MU\xi = \alpha \varphi_1$ . Finally  $m_{\varphi}(M) = \frac{1}{2}$  and it is a contradiction with our assumption and hence  $\dim P = 1$ .

q.e.d

### 2. Characterizing some types of transformations of a logic

For every two elements  $a, b \in L$  we shall write  $a - b = a \wedge b^{\perp}$ .

**Theorem 2.1.** (Recurrence theorem) Let T be a homomorphism of L and let T be invariant in a state m. Then for all  $a \in L$  we have

(1) 
$$m(a - \bigvee_{j=1}^{\infty} T^{j}a) = 0.$$

Proof. Let  $b = a - \bigvee_{j=1}^{\infty} T^{j}a$ , then  $\{T^{j}b\}_{j=0}^{\infty}$  are orthogonal elements of Land therefore  $m(\bigvee_{j=0}^{\infty} T^{j}b) = \sum_{j=0}^{\infty} m(T^{j}b) = \sum_{j=0}^{\infty} m(b) < 1$ . Hence m(b) = 0. q.e.d.

A logic *L* is said to satisfy the finite chain condition (f.c.c.) if  $\{a_m\} \subset L$ with  $a_1 > a_2 > \ldots$  implies that there exists *N* such that  $a_n = a_N$  for n > N(see [3]). A logic *L* is  $\sigma$ -continuous if for  $\{a_n\} \subset L$  with  $a_1 < a_2 < \ldots$  we have  $a \land (\bigvee_{n=1}^{\infty} a_n) = \bigvee_{n=1}^{\infty} (a \land a_n)$  for all  $a \in L$ . It is easy to see that if *L* satisfies f.c.c. then it is  $\sigma$ -continuous. For  $\{a_n\} \subset L$  let  $\limsup_{n \to \infty} a_n - \bigwedge_{n=1}^{\infty} \bigvee_{n=1}^{\infty} a_j$ .

**Theorem 2.2.** (Strong recurrence theorem) Let L be  $\sigma$ -continuous and T be a homomorphism invariant in a state m. Then for all  $a \in L$  we have

(2) 
$$m(a - \limsup T^{j}a) = 0.$$

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Proof. Let us put  $b = a - \limsup T^j a$ , then  $b = a \wedge \bigvee_{n=1} (\bigvee_{j=n} T^j a)^\perp = -\bigvee_{n=1}^{\infty} (a \wedge (\bigvee_{j=n}^{\infty} T^j a)^\perp) = \bigvee_{n=1}^{\infty} (a - \bigvee_{j=n}^{\infty} T^j a) = \bigvee_{n=1}^{\infty} b_n$  where  $b_n = a - \bigvee_{j=n}^{\infty} T^j a$ ,  $n = 1, 2, \ldots$  Applying Theorem 2.1. to a map  $\Pi = T^n$  we get for  $b_n^* = a - -\bigvee_{j=1}^{\infty} \Pi^j a$ ,  $m(b_n^*) = 0$ . But  $b_n < b_n^*$ , therefore  $m(b_n) = 0$ ,  $n = 1, 2, \ldots$  and  $m(b) = \lim_n m(b_n) = 0$ .

q.e.d.

In the rest of this paper according to [4] some types of transformations will be introduced and relations among them will be shown.

Let T be a transformation  $L \to L$  and m be a state. Then we shall say that T is

- (i) incompressible in a state m: if  $a \in L$ , a < Ta implies m(Ta a) = 0;
- (ii) conservative in a state m: if  $a \in L$ ,  $a \perp T^n a$ , n = 1, 2, ... implies m(a) = 0;
- (iii) weakly conservative in a state m: if  $a \in L$ ,  $\{T^n a\}_{n=0}^{\infty}$  is a sequence of mutually orthogonal elements of L, then m(a) = 0;

(iv) recurrent in a state m: if  $a \in L$ , then  $m(a - \bigvee_{n=1}^{\infty} T^n a) = 0$ ;

(v) strongly recurrent in a state m: if  $a \in L$ , then  $m(a - \limsup T^n a) = 0$ .

Remark 3. If T is a homomorphism of L invariant in a state m, then T is conservative in m.

Theorem 2.3. Let L be 
$$\sigma$$
-continuous, then (v) implies (iv).  
Proof. Let  $a \in L$ , then  $a - \limsup T^n a = a - \bigwedge_{n=1}^{\infty} \bigvee_{k=n}^{\infty} T^k a = \bigvee_{n=1}^{\infty} (a - \sum_{k=n}^{\infty} T^k a) > a - (\bigvee_{k=1}^{\infty} T^k a)$  and hence  $m(a - \bigvee_{n=1}^{\infty} T^n a) = 0$ .  
q.e.d.

**Theorem 2.4.** Let T be a monotonic transformation, that is Ta < Tb if a < b, m be a state, then (v) implies (iv), and (ii) and (iv) are equivalent.

Proof. Let 
$$a \in L$$
 and let (v) hold, then  $a - \bigvee_{n=1}^{\infty} T^n a = (a - \bigvee_{n=1}^{\infty} T^n a) - \lim_{n \to \infty} \sup T^k (a - \bigvee_{n=1}^{\infty} T^n a)$ . Indeed, if  $b$  is the element on the right-hand side, then  $b < a$   $\bigvee_{n=1}^{\infty} T^n a$ . Since  $\limsup T^k (a - \bigvee_{n=1}^{\infty} T^n a) < \limsup T^k a < \bigvee_{k=1}^{\infty} T^k a_k$ .

we have  $b > (a - \bigvee_{n=1}^{\infty} T^n a) - \bigvee_{n=1}^{\infty} T^n a = a - \bigvee_{n=1}^{\infty} T^n a$  and therefore  $m(a - - \bigvee_{n=1}^{\infty} T^n a) = m(b) = 0$ . Let now (ii) hold. Then if  $a \in L$ , let  $b = a - \bigvee_{n=1}^{\infty} T^n a$ . For each m = 1, 2, ...we get  $T^m(a - \bigvee_{n=1}^{\infty} T^n a) < T^m a < a^{\perp} \lor \bigvee_{n=1}^{\infty} T^n a = (a - \bigvee_{n=1}^{\infty} T^n a)^{\perp}$ , therefore  $T^m b \perp b$  and hence m(b) = 0. On the other hand let (iv) be valid, then if  $a \perp T^n a, n = 1, 2, ...$ , we have  $a \perp \bigvee_{n=1}^{\infty} T^n a$  and  $a - \bigvee_{n=1}^{\infty} T^n a = a$ . Therefore

q.e.d.

We shall be able to say something more if we assume the following properties of T

(3) 
$$T(\bigvee_{n=1}^{\infty} a_n) = \bigvee_{n=1}^{\infty} Ta_n \text{ for } \{a_n\} \subset L$$

(4) 
$$T(a^{\perp}) > (Ta)^{\perp}$$
 for all  $a \in L$ .

If  $\{a_{\lambda}\}_{\lambda \in A}$  is a system of orthogonal elements from L, there is a Boolean  $\sigma$ -algebra  $A \subset L$  which contains the given system (see [1]). Therefore the distributive law holds for the orthogonal elements of L.

**Theorem 2.5.** Let T be a transformation  $L \rightarrow L$  with the properties (3), (4) and m be a state, then (i) implies (iii).

Proof. Let 
$$a \in L$$
 and (i) hold. If  $\{T^n a\}_{n=0}^{\infty}$  are orthogonal elements of  $L$ ,  
then for  $b = (\bigvee_{n=0}^{\infty} T^n a)^{\perp}$  we have  $T((\bigvee_{n=0}^{\infty} T^n a)^{\perp}) > (T(\bigvee_{n=0}^{\infty} T^n a))^{\perp} = \bigvee_{n=0}^{\infty} \{T^n a (\bigvee_{m=1}^{\infty} T^m a)^{\perp}\} = a \land (\bigvee_{m=1}^{\infty} T^m a)^{\perp} = a$ . We conclude finally that  $m(a) \leq m(Tb - b) = 0$ .

**Lemma 2.6.** Let T be a homomorphism of L and m be a state, then (ii) (iv) are equivalent.

Proof. The equivalency of (ii) and (iv) has been proved, Theorem 2.4., and (ii) and (iii) are equivalent as can easily be seen.

q.e.d.

q.e.d

(ii) holds, too.

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