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# ON m-JOINT DISTRIBUTION 

ANATOLIJ DVUREČENSKIJ

The $m$-joint distribution as a weak form of the compatibility of observables on a logic is studied and some results are proved. This notion allows to introduce multidimensional statistics of observables into the measurement theory of noncompatible observables.

Let $L$ be a $\sigma$-lattice with the first and the last elements 0 and 1 , respectively, with the orthocomplementation $\perp: a \mapsto a^{\perp}$ for which there hold (i) $\left(a^{\perp}\right)^{\perp}=a$ for all $a \in L^{\dot{\prime}}$; (ii) if $a<b$, then $b^{\perp}<a^{\perp}$; (iii) $a \vee a^{\perp}=1$ for all $a \in L$. Further we assume that if $a<b$, then $b=a \vee\left(b \wedge a^{\perp}\right)$. A $\sigma$-lattice $L$ satisfying the above axioms will be called a logic [5].
An element $0 \neq a \in L$ is an atom of $L$ if $b<a$ implies either $b=a$ or $b=0$. We say that the elements $a, b \in L$, are (i) orthogonal and we write $a \perp b$ if $a<b^{\perp}$; (ii) compatible if there are three mutually orthogonal elements $a_{1}, b_{1}, c \in L$ such that $a=a_{1} \vee c, b=b_{1} \vee c$ and we shall write $a \leftrightarrow b$.
A state is a map $m$ from $L$ into $\langle 0,1\rangle$ such that $m(1)=1$ and $m\left(\bigvee_{i} a_{i}\right)$
$=\sum_{i} m\left(a_{i}\right)$ if $a_{i} \perp a_{j}, i \neq j$. An element $a \in L$ is a carrier of a state $m$ if $m(b)=0$ iff $b \perp a$. If a carrier of $m$ exists, then it is unique. A logic is full if there is a system $M$ of states such that $a=b$ iff $m(a)=m(b)$ for all $m \in M$.

An observable is a map $x: B\left(R_{1}\right) \rightarrow L$ such that (i) $x\left(R_{1}\right)=1$; (ii) $x(E) \perp x(F)$ if $E \cap F=\emptyset$; (iii) $x\left(\bigcup_{i} E_{i}\right)=\bigvee_{i} x\left(E_{i}\right)$ if $E_{t} \cap E_{j}=\emptyset, i \neq j$. We denote by $\sigma(x)$ the
smallest closed set $E \subset R_{1}$ such that $x(E)=1$. An observable $x$ is purely atomic if (i) $\sigma(x)=\left\{\lambda_{1}, \lambda_{2}, \ldots\right\}$; (ii) $x(\{\lambda\})$ is an atom in $L$ for any $\lambda \in \sigma(x)$. The range of an observable $x$ is the set $R(x)=\left\{x(E): E \in B\left(R_{1}\right)\right\}$. The observables $x$ and $y$ are compatible and we write $x \leftrightarrow y$ if $x(E) \leftrightarrow y(F)$ for all $E, F \in B\left(R_{1}\right)$. If $m$ is a state and $x$ an observable, then $m_{x}: E \mapsto m(x(E)), E \in B\left(R_{1}\right)$ is a probabilty measure on $B\left(R_{1}\right)$.

Let $m$ be a state on a logic $L$. We shall say that the observables $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution if there exists a probability measure $\mu_{n}$ on $B\left(R_{n}\right)$ such that

$$
\begin{gather*}
\mu_{n}\left(E_{1} \times \ldots \times E_{n}\right)=m\left(\bigwedge_{j=1}^{n} x_{j}\left(E_{j}\right)\right),  \tag{1}\\
E_{j} \in B\left(R_{1}\right), \quad j=1, \ldots, n .
\end{gather*}
$$

In this case we may study the statistical properties of the observables $x_{1}, \ldots, x_{n}$.
It is evident that if $x_{1}, \ldots, x_{n}$ are mutually compatible and $m$ is an arbitrary state, then $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution [5, Theorem 6.17]; in the opposite case this does not hold in general (Theorem 7).

Theorem 1. Let $L$ be a full logic, then $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution for every $m \in M$ iff $x_{1}, \ldots, x_{n}$ are mutually compatible.

Proof. Let $m \in M$. It is easy to see that for all $x_{i}, x_{i}$ there is the $m$-joint distribution $m_{x i x}$. Then for all $E, F \in B\left(R_{1}\right)$ we have

$$
\begin{gathered}
m\left(x_{i}(E)\right)=m\left(x_{i}(E) \wedge x_{i}\left(R_{1}\right)\right)=m_{x_{i x}(i)}\left(E \times R_{i}\right)= \\
=m_{x_{i, x}( }\left(E \times F^{c}\right)+m_{x_{i x}( }(E \times F)=m\left(x_{i}(E) \wedge x_{i}^{\perp}(F)\right)+m\left(x_{i}(E) \wedge x_{i}(F)\right)= \\
=m\left(\left[x_{i}(E) \wedge x_{i}^{\perp}(F)\right] \vee\left[x_{i}(E) \wedge x_{j}(F)\right] .\right.
\end{gathered}
$$

Similarly, $m\left(x_{j}(F)=m\left(\left[x_{j}(F) \wedge x_{i}^{\perp}(E)\right] \vee\left[x_{i}(E) \wedge x_{j}(F)\right]\right)\right.$. From the fullness of $L$ we have $x_{i}(E) \leftrightarrow x_{j}(F)$ for all $E, F \in B\left(R_{1}\right)$. Q.E.D.

We have seen that if $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution, then any pair $x_{i}, x_{j}$ has it. There arises a natural question regarding the converse implication. We shall show in the following that the answer in some cases is positive.

A valuation on a logic $L$ is a map $v: L \rightarrow R_{1}$ such that

$$
\begin{gather*}
v(a \vee b)+v(a \wedge b)=v(a)+v(b), a, b \in L ;  \tag{i}\\
v(a) \leqslant v(b) \text { if } a<b . \tag{ii}
\end{gather*}
$$

(Any state on $L$ has the property (ii).)
Then a functional $\varrho_{v}: \varrho_{v}(a, b)=v(a \vee b)-v(a \wedge b)$ is a pseudometric on $L$ [1, p. 230]. In this case we may define the quotient logic $\tilde{L}=L / \varrho_{v}$ identifying the elements $\tilde{a}=\tilde{b}$ iff $\varrho_{v}(a, b)=0, a, b \in L$. If $x$ is an observable on $L$, then $\tilde{x}(E) \equiv \tilde{x}(E), E \in B\left(R_{1}\right)$ is an observable on $L$; similarly, if $m$ is a state on $L$, then $\tilde{m}$ defined by the subscription $\tilde{m}(\tilde{a})=m(a)$ is a state on $\tilde{L}$.

Theorem 2. Let a state $m$ on $L$ be a valuation; then the observables $x$ and $y$ have an $m$-joint distribution iff $\tilde{x} \leftrightarrow \tilde{y}$.

Proof. If $x, y$ have an $m$-joint distribution, then for $E, F \in B\left(R_{1}\right)$ we have (i) $m(x(E))=m\left(\left[x(E) \wedge y^{\perp}(F)\right] \vee[x(E) \wedge y(F)]\right)$; (ii) $m(y(F))=$ $=m\left(\left[y(F) \wedge x^{\perp}(E)\right] \vee[x(E) \wedge y(F)]\right)$. If $a=x(E)$ and $b$ is the expression in parentheses on the right-hand side in (i), then $a>b$ and $\varrho_{m}(a, b)=0$. Therefore

$$
\begin{aligned}
& \tilde{x}(E)=\left(\tilde{x}(E) \wedge \tilde{y}^{1}(F)\right) \vee(\tilde{x}(E) \wedge \tilde{y}(F)), \\
& \tilde{y}(F)=\left(\tilde{y}(F) \wedge \tilde{x}^{1}(E)\right) \vee(\tilde{x}(E) \wedge \tilde{y}(F)),
\end{aligned}
$$

that is, $\tilde{x} \leftrightarrow \tilde{y}$.

Let now $\tilde{x} \leftrightarrow \tilde{y}$; then for a state $\tilde{m}$ on $\tilde{L}=L / \varrho_{m}$ there is an $\tilde{m}$-joint distribution $\mu$ such that $\mu(E \times F)=\tilde{m}(\tilde{x}(E) \wedge \tilde{y}(F))$. But on the other hand $\tilde{m}(\tilde{x}(E) \wedge \tilde{y}(F))$ $=\tilde{m}\left([x(E) \wedge y(F)]^{\sim}\right)=m(x(E) \wedge y(F))$.
Q.E.D.

Theorem 3. If a state $m$ is a valuation on $L$, then the observables $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution iff any pair $x_{i}, x_{i}, i, j=1, \ldots, n$ has an $m$-joint distribution.

Proof. The necessity is evident. The sufficient condition follows from Theorem 2, because then we have $\tilde{x}_{i} \leftrightarrow \tilde{x}_{j}, i, j=1, \ldots, n$. and there exists a probability measure $\mu_{n}$ on $B\left(R_{n}\right)$ such that $\mu_{n}\left(E_{1} \times \ldots \times E_{n}\right)=\tilde{m}\left(\bigwedge_{i=1}^{n} \tilde{x}_{i}\left(E_{i}\right)\right)$ $=m\left(\bigcap_{i=1}^{n} x_{i}\left(E_{i}\right)\right)$.

Remark. If any pair $x_{i}, x_{j}$ from $x_{1}, \ldots, x_{n}$ has an $m$-joint distribution, then $m$ on $\bigcup_{i=1}^{n} R\left(x_{i}\right)$ has the property (2) of the valuation. There arises the question whether $m$ has the property (2) on the minimal sublogic $L_{0}$ generated by $R\left(x_{i}\right), i=1, \ldots, n$. In the case of a positive answer the proposition of Theorem 3 holds without the assumption on the valuation. The partial answer is given in the following (Corollary 6) and therefore the limitation on the valuation is not an extremely strong restriction in the study of an $m$-joint distribution.

Theorem 4. Let a be the carrier of a state $m$. Purely atomic observables $x$ and $y$ with the spectra $\sigma(x)=\left\{\lambda_{1}, \lambda_{2}, \ldots\right\}$ and $\sigma(y)=\left\{\mu_{1}, \mu_{2}, \ldots\right\}$ have an $m$-joint distribution iff
(i) there are the index sets $I$ and $J$ such that for any $i \in I$ there is $j \in J$ with

$$
\begin{equation*}
b_{i} \equiv y\left(\left\{\mu_{i}\right\}\right)=x\left(\left\{\lambda_{i}\right\}\right) \equiv a_{i} \tag{3}
\end{equation*}
$$

and for any $j \in J$ there is $i \in I$ such that (3) holds;
(ii)

$$
a<\bigvee_{i \in I} a_{i}=\bigvee_{j \in J} b_{j}
$$

Proof. Let $x$ and $y$ have an $m$-joint distribution. Then we have $1=\sum_{i, j} m\left(a_{i} \wedge b_{i}\right)$ $=m\left(\bigvee_{i, j} a_{i} \wedge b_{j}\right)$. Hence $a<\bigvee_{i, j} a_{i} \wedge b_{i}$. Since $a_{i}, b_{j}$ are atoms of $L$, either $a_{i} \wedge b_{j}=0$ or $a_{i}=b_{j}$. Let us denote by $I$ the set of such indexes $i$ for which there is $j$ such that (3) holds. Analogically we define J. Hence (i) is satisfied.

The property (i) implies $a<\bigvee_{i, j} a_{i} \wedge b_{j}=\bigvee_{i \in I} a_{i}=\bigvee_{i \in J} b_{j}$.
Conversely, let (i) and (ii) hold. Then $\sum_{i, j} m\left(a_{i} \wedge b_{j}\right)=m\left(\bigvee_{i, j} a_{i} \wedge b_{j}\right)=m\left(\bigvee_{i \in I} a_{i}\right)$ $\geqslant m(a)=1$, which is the necessary and sufficient condition for $x$ and $y$ to have an $m$-joint distribution.
Q.E.D.

Theorem 5. Let $\left\{x_{t}, t \in T\right\}$ be purely atomic observables on a modular logic of the rank 3 (i.e., any set of nonzero mutually orthogonal elements of $L$ has at most three elements). If there is $a \in L$ such that for any $t \in T$ there is $\lambda_{t} \in \sigma\left(x_{t}\right)$ with

$$
x_{t}\left(\left\{\lambda_{t}\right\}\right)=a,
$$

then the minimal sublogic generated by $L_{0}=\bigcup_{t \in T} R\left(x_{t}\right)$ is equal to $L_{0}$.
Proof. Let $x$ and $y$ be two arbitrary observables of the given system of observables. If we show that for any $a, b \in R(x) \cup R(y) a \vee b^{\perp}$ is an element of $R(x) \cup R(y)$, then $R(x) \cup R(y)$ is a sublogic of $L$ and consequently $L_{0}$ is a sublogic of $L$, too.

There are two cases (i) $x \leftrightarrow y$, hence $R(x)=R(y)$ and the above proposition is true ; (ii) $x \nleftarrow y$. Let us put $a_{i}=x\left(\left\{\lambda_{i}\right\}\right), b_{j}=y\left(\left\{\mu_{j}\right\}\right), i, j=1,2,3$. Let $a_{1}=b_{1}$; then $a_{2} \wedge b_{2}=a_{2} \wedge b_{3}=a_{3} \wedge b_{2}=a_{3} \wedge b_{3}=0$ and the following table holds for the join $\vee$

| $\vee$ | $b_{1}$ | $b_{2}$ | $b_{3}$ | $b_{1} \vee b_{2}$ | $b_{1} \vee b_{3}$ | $b_{2} \vee b_{3}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $a_{1}$ | $a_{1}$ | $b_{1} \vee b_{2}$ | $b_{1} \vee b_{3}$ | $b_{1} \vee b_{2}$ | $b_{1} \vee b_{3}$ | 1 |
| $a_{2}$ | $a_{1} \vee a_{2}$ | $a_{2} \vee a_{3}$ | $a_{2} \vee a_{3}$ | 1 | 1 | $a_{2} \vee a_{3}$ |
| $a_{3}$ | $a_{1} \vee a_{3}$ | $a_{2} \vee a_{3}$ | $a_{2} \vee a_{3}$ | 1 | 1 | $a_{2} \vee a_{3}$ |
| $a_{1} \vee a_{2}$ | $a_{1} \vee a_{2}$ | 1 | 1 | 1 | 1 | 1 |
| $a_{1} \vee a_{3}$ | $a_{2} \vee a_{3}$ | 1 | 1 | 1 | 1 | 1 |
| $a_{2} \vee a_{3}$ | 1 | $a_{2} \vee a_{3}$ | $a_{2} \vee a_{3}$ | 1 | 1 | $a_{2} \vee a_{3}$ |

We show, for example, $a_{2} \vee b_{2}=a_{2} \vee a_{3}$. Let $v$ be a canonical dimension function [5, p. 26]. Since $a_{2} \perp a_{1}, b_{2} \perp a_{1}$, then $a_{2} \vee b_{2} \perp a_{1}$ and $v\left(a_{2} \vee b_{2} \vee a_{1}\right)=v\left(a_{2} \vee b_{2}\right)$ $+v\left(a_{1}\right)=v\left(a_{2}\right)+v\left(b_{2}\right)+v\left(a_{1}\right)=3$. Hence $a_{2} \vee b_{2} \vee a_{1}=1$ and $a_{2} \vee b_{2}=a_{1}^{\perp}=$ $a_{2} \vee a_{3}$. Similarly for other cases. Hence for any $a, b \in R(x) \cup R(y) a \vee b^{\perp}$ is in $R(x) \cup R(y)$.
Q.E.D.

Corollary 6. Let $\left\{x_{t}, t \in T\right\}$ be a system of purely atomic observables on a modular logic of the rank 3. Let $a \in L$ be such an element that for any $t \in T$ there is $\lambda_{t} \in \sigma\left(x_{t}\right)$ with $x_{t}\left(\left\{\lambda_{t}\right\}\right)=a$. If, moreover, a is the carrier of a state $m$, then
(i) for any $s, t \in T x_{s}$ and $x_{t}$ have an $m$-joint distribution;
(ii) $m$ on $L_{0}=L_{0}\left(\bigcup_{t \in T} R\left(x_{t}\right)\right)$ is a valuation;
(iii) every finite subsystem $x_{t}, \ldots, x_{t n}, t_{i} \in T$, has an $m$-joint distribution.

Proof. The validity of proposition (i)-(iii) is obtained by using Theorems 4, 5, 3. Q.E.D.

One of the most important examples of logics is a logic $L(H)$ of the separable Hilbert space $H$ (real or complex), where the elements of $L(H)$ are all closed subspaces of $H$. The Gleason theorem [5] asserts that every state $m$ on $L(H)$, $\operatorname{dim} H \geqslant 3$, is of the form $m(M)=\sum_{i} \lambda_{i}\left(M \Phi_{i}, \Phi_{i}\right), M \in L(H)$, where $\lambda_{i}>0, \sum_{i} \lambda_{i}=$ 1, or equivalently, $m(M)=m_{T}(M)=\operatorname{tr}(T M)$, where $T=\sum_{i} \lambda_{i}\left(\cdot, \Phi_{i}\right) \Phi_{i}$ is a Hermitian operator of the trace class ( $\left\{\Phi_{i}\right\}_{i}$ is an orthonormal system).

Gudder [3] proved the following theorem:
Theorem 7. Let $x, y$ be observables in $L(H)$. Then $x, y$ have an $m$-joint distribution in a state $m=m_{T}$ iff

$$
\begin{equation*}
x(E) y(F) \Phi_{i}=y(F) x(E) \Phi_{i} \tag{4}
\end{equation*}
$$

for all $E, F \in B\left(R_{1}\right), i=1,2, \ldots$.
We shall say that a subspace $H_{0}$ reduces an observable $x$ on $L(H)$ if $x(E) H_{0}$ $=H_{0} x(E)$ for all $E \in B\left(R_{1}\right)$. Then the map $x_{0}: x_{0}(E) \equiv x(E) H_{0}=x(E) \wedge H_{0}$ is an observable on the logic $L\left(H_{0}\right)$.

Theorem 8. If $x, y$ have an $m$-joint distribution in a state $m=m_{T}$ and if a subspace $H_{0}$ generated by $\left\{\Phi_{i}\right\}_{i}$ reduces the observables $x$ and $y$, then the observables $x_{0}$ and $y_{0}$ are compatible on $L\left(H_{0}\right)$.

Proof. If $M, N \in L(H)$, then $M \leftrightarrow N$ iff $M N=N M$ [5] and therefore it suffices to verify whether $x_{0}(E) y_{0}(F)=y_{0}(F) x_{0}(E), E, F \in B\left(R_{1}\right)$. Due to Theorem 7 we have $x_{0}(E) y_{0}(F)=x(E) H_{0} y(F) H_{0}=x(E) y(F) H_{0}=y(F) x(E) H_{0}$ $=y(F) H_{0} x(E) H_{0}=y_{0}(F) x_{0}(E) . \quad$ Q.E.D.

Theorem 9. Let a subspace $H_{0}$ generated by $\left\{\Phi_{i}\right\}_{i}$ reduce the observables $x_{1}, \ldots, x_{n}$. Then $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution in a state $m=m_{T}$ iff any pair $x_{k}, x_{i}, k, j=1, \ldots, n$ has an $m$-joint distribution.

Proof. Only the sufficient condition. According to Theorem 8, the observables $x_{j 0}(E)=x_{j}(E) H_{0}, j=1, \ldots, n$ are compatible on the logic $L\left(H_{0}\right)$. The state $m=m_{T}$ may be considered on $L\left(H_{0}\right)$, too. Therefore there is a probability measure $\mu_{n}$ on $B\left(R_{n}\right)$ for which we have $\mu_{n}\left(E_{1} \times \ldots \times E_{n}\right)=m_{T}\left(x_{10}\left(E_{1}\right) \wedge \ldots \wedge x_{n 0}\left(E_{n}\right)\right)$. By using the methods of the Hilbert space [4] the right-hand side is equal to

$$
\begin{aligned}
& \sum_{i} \lambda_{i}\left(x_{10}\left(E_{1}\right) \wedge \ldots \wedge x_{n 0}\left(E_{n}\right) \Phi_{i}, \Phi_{i}\right)= \\
& \sum_{i} \lambda_{i}\left(x_{1}\left(E_{1}\right) H_{0} \wedge \ldots \wedge x_{n}\left(E_{n}\right) H_{0} \Phi_{i}, \Phi_{i}\right)=
\end{aligned}
$$

$$
\begin{align*}
& \sum_{i} \lambda_{i}\left(x_{1}\left(E_{1}\right) \wedge H_{0} \wedge \ldots \wedge x_{n}\left(E_{n}\right) \wedge H_{0} \Phi_{i}, \Phi_{i}\right)= \\
& \sum_{i} \lambda_{i}\left(x_{1}\left(E_{1}\right) \wedge \ldots \wedge x_{n}\left(E_{n}\right) \wedge H_{0} \Phi_{i}, \Phi_{i}\right)= \\
& \sum_{i} \lambda_{i}\left(x_{1}\left(E_{1}\right) \wedge \ldots \wedge x_{n}\left(E_{n}\right) \Phi_{i}, \Phi_{i}\right)=m_{T}\left(\bigcap_{j=1}^{n} x_{j}\left(E_{j}\right)\right) .
\end{align*}
$$

Finally, an independence of observables in a state $m$ will be investigated in this contribution. The observables $x_{1}, \ldots, x_{n}$ are independent in a state $m$ if

$$
\begin{equation*}
m\left(\bigcap_{i=1}^{n} x_{i}\left(E_{j}\right)\right)=\prod_{i=1}^{n} m\left(x_{j}\left(E_{i}\right)\right), \tag{5}
\end{equation*}
$$

for all $E_{j} \in B\left(R_{1}\right), j=1, \ldots, n$.
If $L$ is a sum logic $[2,3]$ we shall say that the summable observables $x_{1}, \ldots, x_{n}$ are strongly independent in a state $m$ if for any finite system of bounded observables $f_{1}, \ldots, f_{n}$ there holds

$$
\begin{equation*}
m_{f_{10} \times x_{1}+\ldots+f_{n} \times x_{n}}=m_{f_{1} \times x_{1}} * \ldots * m_{f_{n} \times x_{n}}, \tag{6}
\end{equation*}
$$

where the $*$ is the convolution. (If $f$ is a Borel function and $x$ is an observable, then a map $f \circ x: E \mapsto x\left(f^{-1}(E)\right), E \in B\left(R_{1}\right)$, is an observable.)
These two notions of independence coincide in the case of compatible observables. Gudder [2] showed that the strong independence in $m$ implies the independence in $m$. The equivalency was proved by him only for question observables in a pure state on a logic $L(H)$.

Theorem 10. Let the subspace $H_{0}$ generated by $\left\{\Phi_{i}\right\}_{i}$ reduce the observables $x_{1}$, $\ldots, x_{n}$. Then the independence of the observables in a state $m=m_{T}$ implies the strong independence.

Proof. We may easily see that $x_{1}, \ldots, x_{n}$ have an $m$-joint distribution in a state $m=m_{r}$. Therefore $x_{10}, \ldots, x_{1 n}$ are compatible on $L\left(H_{0}\right)$ and independent in $m=m_{\mathrm{T}}$. Hence they are strongly independent, too, and there holds

$$
\begin{aligned}
& m_{T, f_{1} \times x_{1}+\ldots+f_{n} \times x_{n}}=m_{T,\left(f_{1} \times x_{1}+\ldots+f_{n} \times x_{n) 0}\right.}= \\
& m_{T, f_{1} \times x_{10}+\ldots+f_{n} \times x_{0}}=m_{T, f_{1} \times x_{10} * \ldots * m_{T, f_{n} \times x_{n} 0}}= \\
& m_{T, f 10 x_{1} * \ldots * * m_{r, f_{n} \times x_{n}} .} \text { Q.E.D. }
\end{aligned}
$$

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О $m$-СОВМЕСТНОМ РАСПРЕДЕЛЕНИИ
Анатолий Двуреченский
Резюме
Исследуется $m$-совместное распределение как слабая форма совместности наблюдаемых на логике и доказываются некоторые результаты. Это понятие допускает введение многомерной статистики наблюдаемых в теории несовместных наблюдаемых.

