## PANG 14

Radek Kučera; Sitka Machalová; Pavel Ženčák
Interior point algorithms for 3D contact problems

In: Jan Chleboun and Petr Přikryl and Karel Segeth and Tomáš Vejchodský (eds.): Programs and Algorithms of Numerical Mathematics, Proceedings of Seminar. Dolní Maxov, June 1-6, 2008. Institute of Mathematics AS CR, Prague, 2008. pp. 111-117.

Persistent URL: http://dml.cz/dmlcz/702864

## Terms of use:

© Institute of Mathematics AS CR, 2008

Institute of Mathematics of the Czech Academy of Sciences provides access to digitized documents strictly for personal use. Each copy of any part of this document must contain these Terms of use.


This document has been digitized, optimized for electronic delivery and stamped with digital signature within the project DML-CZ: The Czech Digital Mathematics Library http://dml.cz

# INTERIOR POINT ALGORITHMS FOR 3D CONTACT PROBLEMS 

Radek Kučera, Jitka Machalová, Pavel Ženčák

## 1. Introduction

We shall be concerned with solving

$$
\begin{array}{ll}
\operatorname{minimize} & \frac{1}{2} x^{\top} A x-x^{\top} b, \\
\text { subject to } & x_{1, i} \geq l_{i}, x_{2, i}^{2}+x_{3, i}^{2} \leq g_{i}^{2}, i=1, \ldots, m  \tag{1}\\
& x=\left(x_{1}^{\top}, x_{2}^{\top}, x_{3}^{\top}\right)^{\top} \in \mathbb{R}^{n}
\end{array}
$$

where $\left|x_{1}\right|=\left|x_{2}\right|=\left|x_{3}\right|=m, n=3 m, A \in \mathbb{R}^{n \times n}$ is the symmetric, positive definite Hessian matrix, $b \in \mathbb{R}^{n}$, and $l, g \in \mathbb{R}^{m}$. This problem arises, e.g., in duality based methods for the solution of 3D contact problems of linear elasticity with Tresca friction. As a widely used approach of contact problems with (more realistic) Coulomb friction is based on a sequence of Tresca friction problems [2], an efficient solver for (1) is of crucial importance. In this contribution we shall test algorithms based on an "interior point" idea.

## 2. Description of algorithms

The solution to (1) exists and it is necessarily unique. We denote it by $x^{*}$. It is well-known [1] that $x^{*}$ is fully determined by the Karush-Kuhn-Tucker (KKT) conditions. The basic idea of interior point methods consists in applying Newton iterations to equalities in the KKT conditions while inequalities are satisfied strictly by damping Newton steps.

Let us introduce the Lagrangian $\mathcal{L}: \mathbb{R}^{n} \times \mathbb{R}^{2 m} \mapsto \mathbb{R}$ associated with (1) by

$$
\mathcal{L}(x, \lambda, \mu)=\frac{1}{2} x^{\top} A x-x^{\top} b+\lambda^{\top}\left(l-x_{1}\right)+\mu^{\top}\left(X_{2}^{2}+X_{3}^{2}-G^{2}\right) e,
$$

where $X_{2}, X_{3}, G \in \mathbb{R}^{m \times m}$ are defined by $X_{2}=\operatorname{diag}\left(x_{2}\right), X_{3}=\operatorname{diag}\left(x_{3}\right), G=\operatorname{diag}(g)$, and $e=(1, \ldots, 1)^{\top} \in \mathbb{R}^{m}$. There is $y^{*}:=\left(\lambda^{*}, s^{*}, \mu^{*}, d^{*}\right) \in \mathbb{R}^{4 m}$ so that the pair $\left(x^{*}, y^{*}\right)$ is the unique solution to the following system:

$$
\begin{align*}
& \frac{\partial \mathcal{L}}{\partial x}(x, \lambda, \mu)=0, \frac{\partial \mathcal{L}}{\partial \lambda}(x, \lambda, \mu)+s=0, \lambda^{\top} s=0, \frac{\partial \mathcal{L}}{\partial \mu}(x, \lambda, \mu)+d=0, \mu^{\top} d=0 \\
& \lambda \geq 0, s \geq 0, \mu \geq 0, d \geq 0 \tag{2}
\end{align*}
$$

[^0]Here, $\lambda$, and $\mu$ are the Lagrange multipliers while $s$, and $d$ are the slack variables. The classical KKT conditions can be derived from (2), (3) by eliminating the slack variables.

Let us divide $A, b$ into blocks $A_{i j}, b_{i}, i, j \in\{1,2,3\}$ consistently with the partition of $x$ onto $x_{1}, x_{2}, x_{3}$. We can equivalently rewrite (2), (3) as

$$
\begin{equation*}
F(x, y)=0, \quad y \geq 0 \tag{4}
\end{equation*}
$$

where $y:=\left(\lambda^{\top}, s^{\top}, \mu^{\top}, d^{\top}\right)^{\top}$, and $F: \mathbb{R}^{n+4 m} \mapsto \mathbb{R}^{n+4 m}$,

$$
F(x, y):=\left(\begin{array}{l}
A_{11} x_{1}+A_{12} x_{2}+A_{13} x_{3}-\lambda-b_{1} \\
A_{21} x_{1}+\left(A_{22}+2 M\right) x_{2}+A_{23} x_{3}-b_{2} \\
A_{31} x_{1}+A_{32} x_{2}+\left(A_{33}+2 M\right) x_{3}-b_{3} \\
\hline-x_{1}+s+l \\
\Lambda S e \\
\left(X_{2}^{2}+X_{3}^{2}-G^{2}\right) e+d \\
M D e
\end{array}\right)
$$

with $\Lambda, S, M, D \in \mathbb{R}^{m \times m}, \Lambda=\operatorname{diag}(\lambda), S=\operatorname{diag}(s), M=\operatorname{diag}(\mu), D=\operatorname{diag}(d)$. The Jacobi matrix to $F$ reads as follows:

$$
J(x, y)=\left(\begin{array}{rrr|rrrr}
A_{11} & A_{12} & A_{13} & -I & 0 & 0 & 0  \tag{5}\\
A_{21} & A_{22}+2 M & A_{23} & 0 & 0 & 2 X_{2} & 0 \\
A_{31} & A_{32} & A_{33}+2 M & 0 & 0 & 2 X_{3} & 0 \\
\hline-I & 0 & 0 & 0 & I & 0 & 0 \\
0 & 0 & 0 & S & \Lambda & 0 & 0 \\
0 & 2 X_{2} & 2 X_{3} & 0 & 0 & 0 & I \\
0 & 0 & 0 & 0 & 0 & D & M
\end{array}\right) .
$$

Let $\left(x^{(k)}, y^{(k)}\right), y^{(k)}>0$ be a known approximation of $\left(x^{*}, y^{*}\right)$. The Newton direction is the solution to the linear system

$$
\begin{equation*}
J\left(x^{(k)}, y^{(k)}\right)\binom{\Delta x^{(k+1)}}{\Delta y^{(k+1)}}=-F\left(x^{(k)}, y^{(k)}\right) \tag{6}
\end{equation*}
$$

and the new iterate is

$$
\begin{equation*}
\left(x^{(k+1)}, y^{(k+1)}\right)=\left(x^{(k)}, y^{(k)}\right)+\alpha^{(k)}\left(\Delta x^{(k+1)}, \Delta y^{(k+1)}\right), \tag{7}
\end{equation*}
$$

where $\alpha^{(k)}=\min _{\Delta y_{i}^{(k+1)}<0}\left\{1,-\delta y_{i}^{(k)} / \Delta y_{i}^{(k+1)}\right\}$ with $\delta \in(0,1]$ providing $y^{(k+1)}>0$ (typically $\delta=0.999$ ).

The computations based on (6), (7) can take short steps before violating $y^{(k+1)}>0$ so that the convergence rate can be slow. Therefore we use two modifications [7] called path following method, and (Mehrotra's) predictor-corrector method keeping iterates deeper in the feasible region so that it enables us to perform longer steps.

Let us replace (4) by

$$
\begin{equation*}
F(x, y)=\left(0^{\top}, 0^{\top}, 0^{\top}, 0^{\top}, \tau e^{\top}, 0^{\top}, \tau e^{\top}\right)^{\top}, \quad y>0 \tag{8}
\end{equation*}
$$

where $\tau>0$. Solutions ( $x^{\tau}, y^{\tau}$ ) to (8) define in $\mathbb{R}^{n} \times \mathbb{R}_{+}^{4 m}$ a curve $\mathcal{C}(\tau)$ called central path that leads to $\left(x^{*}, y^{*}\right)$ as $\tau$ tends to zero. The next algorithm combines Newton iterations for the equation in (8) with changes of $\tau$ so that the (modified) Newton steps follow $\mathcal{C}(\tau)$.
Algorithm PF: Given $x^{(0)} \in \mathbb{R}^{n}, y^{(0)} \in \mathbb{R}_{+}^{4 m}, \sigma_{l}, \sigma_{q}, \delta \in[0,1]$, and $\epsilon \geq 0$. Set $k:=0$.
(1) Compute $\beta_{l}^{(k)}=\lambda^{(k)^{\top}} s^{(k)} / m, \beta_{q}^{(k)}=\mu^{(k)^{\top}} d^{(k)} / m$, and $\tau_{l}^{(k)}=\sigma_{l} \beta_{l}^{(k)}, \tau_{q}^{(k)}=\sigma_{q} \beta_{q}^{(k)}$. Solve

$$
\begin{equation*}
J\left(x^{(k)}, y^{(k)}\right)\binom{\Delta x^{(k+1)}}{\Delta y^{(k+1)}}=-F\left(x^{(k)}, y^{(k)}\right)+\left(0^{\top}, 0^{\top}, 0^{\top}, 0^{\top}, \tau_{l}^{(k)} e^{\top}, 0^{\top}, \tau_{q}^{(k)} e^{\top}\right)^{\top} \tag{9}
\end{equation*}
$$

and generate $\left(x^{(k+1)}, y^{(k+1)}\right)$ by (7).
(2) If $\left\|\left(\Delta x^{(k+1)}, \Delta y^{(k+1)}\right)\right\|_{\mathbb{R}^{n+4 m}} \leq \epsilon$, return $(\bar{x}, \bar{y})=\left(x^{(k+1)}, y^{(k+1)}\right)$, else set $k:=$ $k+1, \quad$ and go to step (1).

The parameters $\beta_{l}^{(k)}, \beta_{q}^{(k)}$, and $\sigma_{l}, \sigma_{q}$ are called duality measures, and centering parameters, respectively. Let us note that $\sigma_{l}=\sigma_{q}=0$ reduces Algorithm PF to the standard (damped) Newton method. Our choices of $\sigma_{l}$, and $\sigma_{q}$ are based on the rule proposed in [4]:

$$
\sigma_{l}=\left(\min \left\{2 \cdot 10^{-3}, 5 \cdot 10^{-5}\left(1-\xi_{l}\right) / \xi_{l}\right\}\right)^{3},
$$

where $\xi_{l}=\min _{i=1, \ldots, m}\left\{\lambda_{i}^{(k)} s_{i}^{(k)}\right\} / \beta_{l}^{(k)}$, and analogously for $\sigma_{q}$.
The second algorithm calculates centering parameters adaptively using second order information (curvature) of the central path $\mathcal{C}(\tau)$. First, in the predictor stage, we compute duality measures $\beta_{l}^{P}, \beta_{q}^{P}$ for the longest step of the (standard) Newton direction. Then, in the corrector stage, we set $\sigma_{l}, \sigma_{q}$ near 0 , when the good progress along the predicted direction is made or near 1 conversely.
Algorithm PC: Given $x^{(0)} \in \mathbb{R}^{n}, y^{(0)} \in \mathbb{R}_{+}^{4 m}, \delta \in(0,1)$, and $\epsilon \geq 0$. Set $k:=0$.
(1) Solve

$$
J\left(x^{(k)}, y^{(k)}\right)\binom{\Delta x^{P}}{\Delta y^{P}}=-F\left(x^{(k)}, y^{(k)}\right)
$$

compute $\alpha^{P}=\min _{\Delta y_{i}^{P}<0}\left\{1,-y_{i}^{(k)} / \Delta y_{i}^{P}\right\}$, and

$$
\begin{aligned}
\beta_{l}^{P} & =\left(\lambda^{(k)}+\alpha^{P} \Delta \lambda^{P}\right)^{\top}\left(s^{(k)}+\alpha^{P} \Delta s^{P}\right) / m, \\
\beta_{q}^{P} & =\left(\mu^{(k)}+\alpha^{P} \Delta \mu^{P}\right)^{\top}\left(d^{(k)}+\alpha^{P} \Delta d^{P}\right) / m .
\end{aligned}
$$

(2) Set $\sigma_{l}=\left(\beta_{l}^{P} / \beta_{l}^{(k)}\right)^{3}, \sigma_{q}=\left(\beta_{q}^{P} / \beta_{q}^{(k)}\right)^{3}$, compute ( $\left.\Delta x^{(k+1)}, \Delta y^{(k+1)}\right)$ solving (9) with the right-hand-side replaced by
$-F\left(x^{(k)}, y^{(k)}\right)+\left(0^{\top}, 0^{\top}, 0^{\top}, 0^{\top},-e^{\top} \Delta \Lambda^{P} \Delta S^{P}+\tau_{l}^{(k)} e^{\top}, 0^{\top},-e^{\top} \Delta M^{P} \Delta D^{P} e+\tau_{q}^{(k)} e^{\top}\right)^{\top}$, and generate $\left(x^{(k+1)}, y^{(k+1)}\right)$ by (7).
(3) If $\left\|\left(\Delta x^{(k+1)}, \Delta y^{(k+1)}\right)\right\|_{\mathbb{R}^{n+4 m}} \leq \epsilon$, return $(\bar{x}, \bar{y})=\left(x^{(k+1)}, y^{(k+1)}\right)$, else set $k:=k+1$, and go to step (1).

## 3. Numerical experiments

### 3.1. Model problem

Let us consider a steel brick in $\mathbb{R}^{3}$ lying on a rigid foundation. The brick occupies the domain $\Omega=(0,3) \times(0,1) \times(0,1)$, whose boundary $\partial \Omega$ split into three nonempty disjoint parts $\Gamma_{u}=\{0\} \times(0,1) \times(0,1), \Gamma_{c}=(0,3) \times(0,1) \times\{0\}$, and $\Gamma_{p}=\partial \Omega \backslash\left(\bar{\Gamma}_{u} \cup \bar{\Gamma}_{c}\right)$ with different boundary conditions. The zero displacements are prescribed on $\Gamma_{u}$, whereas the surface tractions act on $\Gamma_{p}$. On $\Gamma_{c}$ we consider the contact conditions, i.e., the non-penetration, and the effect of friction. The elastic behavior of the brick is described by Lamé equations that, after finite element discretization, lead to a symmetric positive definite stiffness matrix $K \in \mathbb{R}^{3 n_{c} \times 3 n_{c}}$ and to a load vector $f \in \mathbb{R}^{3 n_{c}}$. Moreover, we introduce full rank matrices $N, T_{1}, T_{2} \in \mathbb{R}^{m \times 3 n_{c}}$ projecting displacements at contact nodes to normal and tangential directions, respectively, and we denote $B=\left(N^{\top}, T_{1}^{\top}, T_{2}^{\top}\right)^{\top} \in \mathbb{R}^{3 m \times 3 n_{c}}$. For more details about this model problem we refer to [2].

Here, we shall use the dual formulation in terms of contact stresses. Considering only Tresca friction, our model problem reduces directly to (1), where $A=B K^{-1} B^{\top}$, $b=B K^{-1} f, l=0$, and $g_{i} \geq 0$ are given slip bound values at contact nodes. Let us note that unknowns $x_{1}$, and $x_{2}, x_{3}$ represent the normal, and tangential contact stresses, respectively.

### 3.2. Inner solver

Our algorithms require repeatedly to solve the linear systems

$$
\begin{equation*}
J(x, y)\binom{\Delta x}{\Delta y}=\binom{r_{x}}{r_{y}} \tag{10}
\end{equation*}
$$

with the Jacobi matrix given by (5), and $r_{x} \in \mathbb{R}^{n}, r_{y}=\left(r_{\lambda}^{\top}, r_{s}^{\top}, r_{\mu}^{\top}, r_{d}^{\top}\right)^{\top} \in \mathbb{R}^{4 m}$, $\Delta y=\left(\Delta \lambda^{\top}, \Delta s^{\top}, \Delta \mu^{\top}, \Delta d^{\top}\right)^{\top} \in \mathbb{R}^{4 m}$. First we compute the solution to the reduced system arising from (10) by eliminating increments with respect to the slack variables:

$$
\begin{equation*}
J_{R}(x, y)\binom{\Delta x}{\Delta z}=\binom{r_{x}}{r_{z}} \tag{11}
\end{equation*}
$$

where

$$
J_{R}(x, y)=\left(\begin{array}{rrr|rr}
A_{11} & A_{12} & A_{13} & -I & 0 \\
A_{21} & A_{22}+2 M & A_{23} & 0 & 2 X_{2} \\
A_{31} & A_{32} & A_{33}+2 M & 0 & 2 X_{3} \\
\hline-I & 0 & 0 & -\Lambda^{-1} S & 0 \\
0 & 2 X_{2} & 2 X_{3} & 0 & -M^{-1} D
\end{array}\right),
$$

and $r_{z}=\left(r_{\lambda}^{\top}-r_{s}^{\top} \Lambda^{-1}, r_{\mu}^{\top}-r_{d}^{\top} M^{-1}\right)^{\top} \in \mathbb{R}^{2 m}, \Delta z=\left(\Delta \lambda^{\top}, \Delta \mu^{\top}\right)^{\top} \in \mathbb{R}^{2 m}$. Then we obtain the eliminated components by $\Delta s=\Lambda^{-1}\left(r_{s}-S \Delta \lambda\right)$, and $\Delta d=M^{-1}\left(r_{d}-\right.$ $D \Delta \mu)$. It is easy to prove that the Schur complement with respect to the second diagonal block in $J_{R}(x, y)$ is positive definite provided $y>0$. Therefore $J_{R}(x, y)$ is non-singular but indefinite. In order to solve (11), one can use direct [6] or iterative methods. In this paper we apply the conjugate gradient method with an appropriate indefinite preconditioning [5].

### 3.3. Tests

We compare the algorithms PF, and PC with the one presented in [3], here denoted by QPC. For various numbers of the primal, and dual degrees of freedoms $\left(3 n_{c} / 3 m\right)$, we report the computational time (time), the total number of the matrixvector multiplications $\left(n_{A}\right)$, and, in case of the interior point algorithms, the number of the outer iterations (out), and the number of the full steps (full), i.e. the steps with $\alpha^{(k)}=1$. All computations are performed by Matlab 7 on Pentium $4,2.8 \mathrm{GHz}$ with 1GB RAM.

The first experiments in Table 1 demonstrate the computational strategy in which the Hessian matrix $A$ is assembled (only in PF, and PC). As the time consumed by assembling $A$ predominates for larger $3 n_{c} / 3 m$, this strategy seems to be nonacceptable form more realistic contact problems.

In Tables 2, and 3 we present the computational efficiency of PF, and PC with non-assembled $A$. We test two preconditioners of indefinite type for (10). The computation of Hessian matrix is based on an approximation of a Schur complement

|  | QPC |  | PF |  | PC |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $3 n_{c} / 3 m$ | time | $n_{A}$ | time | time $_{A}$ | time | time $_{A}$ |
| $162 / 54$ | 0.29 | 203 | 0.07 | $0.03(45 \%)$ | 0.12 | $0.03(26 \%)$ |
| $900 / 180$ | 2.08 | 311 | 0.68 | $0.34(50 \%)$ | 1.07 | $0.34(31 \%)$ |
| $2646 / 378$ | 12.91 | 347 | 5.85 | $3.46(59 \%)$ | 7.00 | $3.26(47 \%)$ |
| $5832 / 648$ | 53.4 | 384 | 27.1 | $18.1(67 \%)$ | 27.0 | $15.8(59 \%)$ |
| $10890 / 990$ | 126.2 | 408 | 79.7 | $58.5(73 \%)$ | 90.0 | $60.5(67 \%)$ |
| $18252 / 1404$ | 361.9 | 493 | 246.2 | $192.5(78 \%)$ | 274.0 | $184(67 \%)$ |
| $28350 / 1890$ | 809.4 | 478 | 620.5 | $493.0(79 \%)$ | 677.6 | $493.5(73 \%)$ |

Tab. 1: $A$ is assembled in PF, and PC (time $A_{A}$ is consumed by assembling $A$ ).

|  | QPC |  | PF, precond. 1 |  |  | PC, precond. 1 |  |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $3 n_{c} / 3 m$ | time | $n_{A}$ | time | $n_{A}$ | out | full | time | $n_{A}$ | out | full |
| $162 / 54$ | 0.29 | 203 | 0.25 | 97 | 19 | 11 | 0.30 | 195 | 27 | 16 |
| $900 / 180$ | 2.08 | 311 | 0.94 | 112 | 20 | 12 | 1.20 | 154 | 17 | 11 |
| $2646 / 378$ | 12.91 | 347 | 6.69 | 139 | 22 | 13 | 6.33 | 147 | 13 | 8 |
| $5832 / 648$ | 53.4 | 384 | 29.33 | 173 | 25 | 13 | 23.8 | 157 | 13 | 7 |
| $10890 / 990$ | 126.2 | 408 | 109.5 | 233 | 30 | 13 | 68.3 | 159 | 13 | 6 |
| $18252 / 1404$ | 361.9 | 493 | 265.3 | 244 | 31 | 12 | 177.7 | 183 | 14 | 7 |
| $28350 / 1890$ | 809.4 | 478 | 644.2 | 282 | 36 | 13 | 420.9 | 209 | 16 | 7 |

Tab. 2: $A$ is not assembled, the preconditioner 1 is assembled.

|  | QPC |  | PF, precond. 2 |  |  |  | PC, precond. 2 |  |  |  |
| ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: | ---: |
| $3 n_{c} / 3 m$ | time | $n_{A}$ | time | $n_{A}$ | out | full | time | $n_{A}$ | out | full |
| $162 / 54$ | 0.29 | 203 | 0.28 | 154 | 27 | 12 | 0.23 | 148 | 16 | 12 |
| $900 / 180$ | 2.08 | 311 | 0.98 | 112 | 20 | 14 | 1.05 | 138 | 13 | 8 |
| $2646 / 378$ | 12.91 | 347 | 6.61 | 133 | 22 | 13 | 5.5 | 122 | 12 | 8 |
| $5832 / 648$ | 53.4 | 384 | 26.6 | 150 | 24 | 13 | 22.6 | 145 | 13 | 8 |
| $10890 / 990$ | 126.2 | 408 | 105.4 | 218 | 30 | 13 | 66.4 | 151 | 13 | 7 |
| $18252 / 1404$ | 361.9 | 493 | 253.2 | 224 | 31 | 13 | 169.8 | 169 | 14 | 7 |
| $28350 / 1890$ | 809.4 | 478 | 584.7 | 251 | 33 | 12 | 397.1 | 190 | 15 | 7 |

Tab. 3: $A$ is not assembled, the preconditioner 2 is assembled.
to the stiffness matrix $K$. While the first preconditioner uses diagonal scaling, the second one requires more expensive computations. Let us note that both preconditioners are assembled.

## 4. Conclusions

In the contribution we present our first experience with solving contact problems by the interior point algorithms. As the results seem promising many questions are still open, namely the convergence proof of both algorithms. The future work consists also of applying non-assembled preconditioners, and of implementing theoretically supported inner solvers.

## References

[1] S. Boyd, L. Vandenberghe: Convex optimization. Cambridge University Press, Cambridge 2004.
[2] J. Haslinger, R. Kučera, Z. Dostál: An algorithm for the numerical realization of 3D contact problems with Coulomb friction. J. Comput. Appl. Math., 164-165 (2004), 387-408.
[3] R. Kučera: Convergence rate of an optimization algorithm for minimizing quadratic functions with separable convex constraints. SIAM J. Optim., 19 (2008), 846-862.
[4] J. Nocedal, A. Wächter, R.A. Waltz: Adaptive barrier strategies for nonlinear interior methods. TR RC 23563, IBM T.J. Watson Research Center, 2005.
[5] L. Lukšan, C. Matonoha, J. Vlček: Interior point methods for large-scale nonlinear programming. TR No. 917, Academy of Sciences of the Czech Republic 2004.
[6] J. Machalová, P. Ženčák, R. Kučera: Metody vnitřních bodů pro řešení úlohy nelineárnîho programování. ODAM 2007, pp. 4-17. Can be found online at: http://mant.upol.cz/soubory/odam/odam07sb.pdf
[7] J. Nocedal, S. J. Wright: Numerical optimization. Springer, New York 1999.


[^0]:    *Supported by grants GAČR 101/08/0574, and MSM6198910027.

