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## A note on finite sets of terms closed under subterms and unification

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*Abstract.* The paper contains two remarks on finite sets of groupoid terms closed under subterms and the application of unifying pairs.

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By a term we shall mean a groupoid term. Let us write  $u + v \sim t$  if  $t = f(u) = g(v)$  for a unifying pair  $f, g$  of the terms  $u$  and  $v$ , i.e., if  $t$  is a substitution instance of both  $u$  and  $v$  and any term that is a substitution instance of both  $u$  and  $v$ , is a substitution instance of  $t$ . (A survey of unification theory is contained, for example, in Dershowitz and Jouannaud [1].)

Let us call a set  $S$  of terms SU-closed if it is closed with respect to subterms and whenever  $u + v \sim t$  for two terms  $u, v \in S$ , then  $t \sim t' \in S$  for some  $t'$ .

**Theorem 1.** *There is no finite, SU-closed set of terms containing the following three terms:*

$$(xy \cdot z)x, \quad x(yz \cdot u), \quad x \cdot yx.$$

PROOF: Let us define an infinite sequence  $a_0, a_1, \dots$  of terms as follows:  $a_0$  is a variable;  $a_{i+1} = a_i x$  for a variable  $x$  not occurring in  $a_i$ . So,  $a_i = (((x_0 x_1) x_2) \dots) x_i$ , where  $x_0, \dots, x_i$  is a sequence of pairwise distinct variables. Also, put  $b_i = y a_i$ , where  $y$  is a variable not occurring in  $a_i$ . Hence  $b_2 \sim x(yz \cdot u)$ . It is easy to see that

$$(xy \cdot z)x + b_i \sim ((a_i x)y)a_i \supseteq a_{i+2}$$

for  $i \geq 2$  (where  $x, y$  are two distinct variables not occurring in  $a_i$ , and

$$x \cdot yx + a_{i+1} \sim a_i \cdot x a_i \supseteq b_i$$

for  $i \geq 3$ . □

The depth of a term is defined inductively as follows: the depth of a variable is 0; the depth of  $t_1 t_2$  is  $1 + \max(d_1, d_2)$ , where  $d_1$  is the depth of  $t_1$  and  $d_2$  is the depth of  $t_2$ . So,  $xy \cdot zu$  is of depth 2.

**Theorem 2.** *There exists a finite, SU-closed set of terms containing (up to similarity) all terms of depth at most 2.*

PROOF: The set consists of the terms of depth 2, plus the twelve terms

$$\begin{aligned}
 xx \cdot (xx \cdot xx) &\sim x \cdot xx + xx \cdot y \\
 xy \cdot (xy \cdot xy) &\sim x \cdot xx + xy \cdot z \\
 xx \cdot (xx \cdot y) &\sim x \cdot xy + xx \cdot y \\
 xy \cdot (xy \cdot z) &\sim x \cdot xy + xy \cdot z \\
 xx \cdot (xx \cdot x) &\sim x \cdot xy + xx \cdot yx \\
 xy \cdot (xy \cdot x) &\sim x \cdot xy + xy \cdot zx \\
 xy \cdot (xy \cdot y) &\sim x \cdot xy + xy \cdot zy \\
 xx \cdot (y \cdot xx) &\sim x \cdot yx + xx \cdot y \\
 xy \cdot (z \cdot xy) &\sim x \cdot yx + xy \cdot z \\
 xx \cdot (x \cdot xx) &\sim x \cdot yx + xx \cdot xy \\
 xy \cdot (x \cdot xy) &\sim x \cdot yx + xy \cdot xz \\
 xy \cdot (y \cdot xy) &\sim x \cdot yx + xy \cdot yz
 \end{aligned}$$

and their duals. □

#### REFERENCES

- [1] Dershowitz N., Jouannaud J.-P., *Rewrite systems*, Chapter 6, 243–320 in J. van Leeuwen, ed., *Handbook of Theoretical Computer Science, B: Formal Methods and Semantics*, North Holland, Amsterdam, 1990.

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