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A search for Tribonacci-Wieferich primes

Jiří Klaška

Abstract. Such problems as the search for Wieferich primes or Wall-Sun-Sun primes are intensively studied and often discused at present. This paper is devoted to a similar problem related to the Tribonacci numbers.

1 Introduction

Let T_n denote the *n*-th Tribonacci number defined by $T_{n+3} = T_{n+2} + T_{n+1} + T_n$ with $T_0 = 0, T_1 = 0$, and $T_2 = 1$. This number has been examined by many authors. First by A. Agronomof [1] in 1914 and subsequently by many others. See, for example, [2], [5], [7], [8], [9], [10]. It is well known that $(T_n \mod m)_{n=0}^{\infty}$ is periodic for any modulus m > 1. The least positive integer h satisfying $[T_h, T_{h+1}, T_{h+2}] \equiv [T_0, T_1, T_2] \pmod{m}$ is called a period of $(T_n \mod m)_{n=0}^{\infty}$ and denoted by h(m).

Two problems remain open: 1. Is there a prime p satisfying $h(p) = h(p^2)$ (M. E. Waddill 1978, [10])? 2. Is there a prime p such that $h(p) \neq h(p^2)$ and $\operatorname{ord}_p(\alpha) = \operatorname{ord}_{p^2}(\alpha)$ where $\alpha \in \mathbb{Z}$ is a solution of $x^3 - x^2 - x - 1 \equiv 0 \pmod{p^2}$ (J. Klaška 2007, [5])? Here, $\operatorname{ord}_{p^t}(\alpha)$ denotes the order of α in the multiplicative group of the ring $\mathbb{Z}/p^t\mathbb{Z}$, $t \in \mathbb{N}$. See also [6, Problem 3.2]. In [6], the primes p satisfying $h(p) = h(p^2)$ are called Tribonacci-Wieferich primes and the primes for which $h(p^2) \neq h(p)$ and $\operatorname{ord}_p(\alpha) = \operatorname{ord}_{p^2}(\alpha)$ where $\alpha \in \mathbb{Z}$ is a solution of $x^3 - x^2 - x - 1 \equiv 0 \pmod{p^2}$ are called Tribonacci-Wieferich primes of the second kind. In [6] we proved that neither of this problems has a solution for $p < 10^9$. In the present paper we substantially extend these results focussing on the case of the Tribonacci characteristic polynomial $t(x) = x^3 - x^2 - x - 1$ being irreducible modulo p.

2 Tribonacci modulo p^2 – an irreducible case

Let $I = \{3, 5, 23, 31, ...\}$ be the set of all primes p for which t(x) is irreducible over $\mathbb{F}_p = \mathbb{Z}/p\mathbb{Z}$. Let K be the splitting field of t(x) over \mathbb{F}_p , $p \in I$ and α, β, γ the roots of t(x) in K. Clearly, $K = GF(p^3)$ and the multiplicative group of K has $p^3 - 1$

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elements. Using the Frobenius authomorphism, we can easily prove that $\beta = \alpha^p$ and $\gamma = \alpha^{p^2}$. This implies that α, β, γ have the same order in the multiplicative group of K. It is well known, see e.g. [5], [6], [8], that for any prime $p \neq 2, 11$:

$$h(p) = \operatorname{lcm}(\operatorname{ord}_L(\alpha), \operatorname{ord}_L(\beta), \operatorname{ord}_L(\gamma))$$
(2.1)

where L is the splitting field of t(x) over \mathbb{F}_p and $\operatorname{ord}_L(\alpha)$, $\operatorname{ord}_L(\beta)$, $\operatorname{ord}_L(\gamma)$ are the orders of α, β, γ in the multiplicative group of L. Consequently, for $p \in I$, we can state

Lemma 2.1. Let $p \in I$. Then $h(p) = \operatorname{ord}_K(\alpha)$ where α is any root of t(x) in a splitting field K of t(x) over \mathbb{F}_p .

Lemma 2.2. For any prime $p \in I$ we have $h(p)|p^2 + p + 1$.

Proof. The Viète equation $\alpha\beta\gamma = 1$ together with $\beta = \alpha^p$ and $\gamma = \alpha^{p^2}$ yields $\alpha^{p^2+p+1} = 1$. This implies $\operatorname{ord}_K(\alpha)|p^2 + p + 1$ and the relation $h(p)|p^2 + p + 1$ follows from Lemma 2.1.

Remark 2.3. In the relation $h(p)|p^2 + p + 1$ it is often, but not always, true that $h(p) = p^2 + p + 1$. For example, $h(3) = 3^2 + 3 + 1 = 13$ but $h(31) = (31^2 + 31 + 1)/3 = 331$.

In 1978, M. E. Waddill [10, Theorem 8] proved that for any prime p:

If
$$h(p) \neq h(p^2)$$
, then $h(p^t) = p^{t-1}h(p)$ for any $t \in \mathbb{N}$. (2.2)

Consequently, we have either $h(p^2) = p \cdot h(p)$ or $h(p^2) = h(p)$. If we combine Waddill's result (2.2) with Lemma 2.2, we obtain

Lemma 2.4. For any prime $p \in I$, $h(p) = h(p^2)$ if and only if $h(p^2)|p^2 + p + 1$.

Now we show that to calculate the powers of α in the multiplicative group of K we need to calculate with Tribonacci numbers.

Lemma 2.5. For any positive integer $n \ge 3$ we have the identity

$$x^{n} = T_{n}x^{2} + (T_{n-1} + T_{n-2})x + T_{n-1} + s_{n}(x)t(x) \text{ where } s_{n}(x) = \sum_{k=1}^{n} T_{k}x^{n-k}.$$
(2.3)

Proof. Using induction on n.

Reducing the identity (2.3) by the double modulus modd(m, t(x)) where m > 1 is an arbitrary positive integer, we obtain the congruence

$$x^{n} \equiv T_{n}x^{2} + (T_{n-1} + T_{n-2})x + T_{n-1} \pmod{m, t(x)}.$$
(2.4)

From (2.4) now it follows that

 $x^n \equiv 1 \pmod{m, t(x)}$ if and only if $[T_n, T_{n+1}, T_{n+2}] \equiv [0, 0, 1] \pmod{m}$. (2.5)

Particulary, if $m = p, p \in I$ and $x = \alpha$ where α is any root of t(x) in K, (2.5) implies Lemma 2.1.

Example 2.6. Let p = 3. Then $p^2 + p + 1 = 13$ and by (2.4) we have $x^{13} \equiv 504x^2 + 423x + 274 \equiv 4 \neq 1 \pmod{3^2, t(x)}$. From (2.5) now it follows that $h(3) \neq h(3^2)$ and thus p = 3 is not a Tribonacci-Wieferich prime. Moreover, from Lemma 2.2 and $h(3) \neq 1$, it follows that h(3) = 13 and by (2.2) we have $h(3^2) = 39$.

Let $q \in I$. By I_q denote the set of all primes $p \in I$ not exceeding q. Theoretically, we have two posibilites when searching for Tribonacci-Wieferich primes in I_q . First, we can calculate a finite sequence $(T_n)_{n=0}^{q^2+q+1}$ and, subsequently, for any particular primes $p \in I_q$, test whether $[T_{p^2+p+1}, T_{p^2+p+2}, T_{p^2+p+3}] \equiv [0, 0, 1] \pmod{p^2}$. Second, we compute the reduced sequences $(T_n \mod p^2)_{n=0}^{p^2+p+1}$ for any $p \in I_q$.

Let us now show that the first possibility is virtually excluded as it uses an enormous amount of computer memory. It can be easily proved that the Tribonacci polynomial t(x) has one real root

$$\tau = \frac{1}{3} \left(\sqrt[3]{19 + 3\sqrt{33}} + \sqrt[3]{19 - 3\sqrt{33}} + 1 \right) \approx 1.839\ 286\ 755\ 214\ 161\ 132\ \cdots \ (2.6)$$

and two complex roots $\sigma, \overline{\sigma}$ ($\overline{\sigma}$ is the complex conjugate of σ) where

$$\sigma = \frac{1}{6} \left(2 - \sqrt[3]{19 + 3\sqrt{33}} - \sqrt[3]{19 - 3\sqrt{33}} \right) + \frac{\sqrt{3}i}{6} \left(\sqrt[3]{19 + 3\sqrt{33}} - \sqrt[3]{19 - 3\sqrt{33}} \right).$$
(2.7)

Put $\varepsilon = \tau^2/|\tau - \sigma|^2 \approx 0.618$ 419 922 319 392 550 · · · . In [7], W. R. Spickerman proved that for T_n we have

$$T_n = [\varepsilon \cdot \tau^n + 0.5]. \tag{2.8}$$

Here [x] denotes the greatest integer not exceeding x. Clearly, if x is positive, then [x] is simply the integer part of x. Note that, in [7], σ is incorrect. See [7, p. 119]. From (2.8) it follows that, for log T_n , we have

$$\log T_n \approx n \cdot \log \tau$$
 where $\log \tau = 0.264\ 649\ 443\ 484\ 250\ 871\cdots$. (2.9)

Evidently, T_n has exactly k digits for n > 1 if and only if $k - 1 \le \log T_n < k$. This, together with (2.9) yields an estimate for the number of digits of T_n . The following example may provide a more precise idea of the greatness of Tribonacci numbers T_n .

Example 2.7. The Tribonacci number T_{100} has 26 digits, T_{1000} has 264 digits, and T_{10000} has 2646 digits. Consider now the greatest prime p from the interval $[2, 10^9]$ for which t(x) is irreducible modulo p. This p is equal to 999999929. To test whether $h(p) = h(p^2)$ we need to find $[T_q, T_{q+1}, T_{q+2}]$ where $q = p^2 + p + 1 =$ 999999859000004971. Since, by (2.9), T_q has more than $5 \cdot 10^{15}$ digits, we need about 10^6 GB of memory for T_q , assuming that one byte is needed for one digit.

In this paper, we use a method based on matrix algebra to search for Tribonacci-Wieferich primes on a given set I_q using a computer. It is well known (see e.g. [5], [9]) that Tribonacci numbers can be computed by powers of the Tribonacci matrix ${\cal T}$ where

$$T = \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \\ 1 & 1 & 1 \end{bmatrix} \text{ and } T^{n+1} = \begin{bmatrix} T_n & T_{n-1} + T_n & T_{n+1} \\ T_{n+1} & T_n + T_{n+1} & T_{n+2} \\ T_{n+2} & T_{n+1} + T_{n+2} & T_{n+3} \end{bmatrix} \text{ for } n \in \mathbb{N}.$$

$$(2.10)$$

Clearly, h(m) is the period of $(T_n \mod m)_{n=0}^{\infty}$ if and only if h(m) is the smallest positive integer h for which $T^h \equiv E \pmod{m}$ where E is the 3×3 identity matrix. This, together with Lemma 2.4, yields

Lemma 2.8. For any $p \in I$ we have $h(p) = h(p^2)$ if and only if $T^{p^2+p+1} \equiv E \pmod{p^2}$ where E is the 3×3 identity matrix.

Now we briefly describe the algorithm used to prove the main theorem of this section.

Algorithm for testing $h(p) = h(p^2)$ for $p \in I$

First, we find a 2-adic expansion of $p^2 + p + 1 = c_0 + 2c_1 + 2^2c_2 + \dots + 2^kc_k$.

Second, we define the matrix $T \mod p^2$ and, subsequently, we compute k matrices $T^{2^i} \mod p^2$ for $i = 1, \ldots, k$.

Third, we compute the matrix

$$T^{p^2+p+1} \mod p^2 = \prod_{i=0}^k (T^{2^i} \mod p^2)^{c_i}.$$
 (2.11)

Finally, we test whether $T^{p^2+p+1} \mod p^2$ is equal to the identity matrix E. This process is repeated for every prime $p \in I$.

Implementing this algorithm in Pari GP, we have obtained the following result:

Theorem 2.9. For any prime $p \in I$, $p < 10^{11}$ we have $h(p) \neq h(p^2)$.

Let us remark that, achieving this result takes about 1500 hours of CPU time on a 1.6 GHz processor computer.

3 Searching for Tribonacci-Wieferich primes $p \not\in I$

In the case of $p \notin I$ we can use the criteria derived in [6] to search for Tribonacci-Wieferich primes. Moreover, when dealing with this case, Tribonacci-Wieferich primes of the second kind may also be found easily. Indeed, by [5], from $h(p) = h(p^2)$, we have $\operatorname{ord}_p(\xi) = \operatorname{ord}_{p^2}(\xi)$ for any solution $\xi \in \mathbb{Z}$ of $t(x) \equiv 0 \pmod{p^2}$. Next, according to [6], if $\alpha \in \mathbb{Z}$ is the unique root of t(x) modulo p with the property

$$3\alpha^{p+2} - 2\alpha^{p+1} - \alpha^p - 2\alpha^3 + \alpha^2 - 1 \equiv 0 \pmod{p^2}$$
(3.1)

or, equivalently, with the property

$$\alpha^{3p} - \alpha^{2p} - \alpha^p - 1 \equiv 0 \pmod{p^2} \tag{3.2}$$

then p is the Tribonacci-Wieferich prime of the second kind. It should be stressed that the criteria (3.1) and (3.2) make it possible to find Tribonacci-Wieferich primes of the second kind and thus also Tribonacci-Wieferich primes p with $p \notin I$ without having to calculate with Tribonacci numbers. The following result has been obtained using (3.1) in Pari GP.

Theorem 3.1. There is no prime $p \notin I$, $p < 10^{11}$ satisfying $\operatorname{ord}_p(\xi) = \operatorname{ord}_{p^2}(\xi)$ where $\xi \in \mathbb{Z}$ is a solution of $t(x) \equiv 0 \pmod{p^2}$. Consequently, there is no Tribonacci-Wieferich prime of the second kind less than 10^{11} .

Note that, as compared with Theorem 2.9, only about 700 hours of CPU time are needed to obtain Theorem 3.1 on the same computer.

Corollary 3.2. For any prime $p \notin I$, $p < 10^{11}$, we have $h(p) \neq h(p^2)$.

If we combine Corollary 3.2 with Theorem 2.9, we obtain the main theorem of this paper:

Theorem 3.3. There is no Tribonacci-Wieferich prime $p < 10^{11}$.

Moreover, based on (2.2), we can now state

Corollary 3.4. For any prime $p < 10^{11}$ and for any $t \in \mathbb{N}$, we have $h(p^t) = p^{t-1}h(p)$.

Remark 3.5. Like in the problem of finding Fibonacci-Wieferich primes (see [3], [4]) also in the Tribonacci case a question may be raised whether the probability of some primes being Tribonacci-Wieferich is greater than that of others. Using a reasoning similar to that used in [4], we can conclude that further search of the set I for $p > 10^{11}$ will virtually not increase the probability of finding a Tribonacci-Wieferich prime. Consequently, the chances of finding Tribonacci-Wieferich primes on a computer seem to be greater for primes not in I, particulary, for those for which t(x) can be factorized into linear terms over \mathbb{F}_p .

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