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# THE DIAMETER OF PAIRED-DOMINATION VERTEX CRITICAL GRAPHS 

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#### Abstract

In this paper we continue the study of paired-domination in graphs introduced by Haynes and Slater (Networks 32 (1998), 199-206). A paired-dominating set of a graph $G$ with no isolated vertex is a dominating set of vertices whose induced subgraph has a perfect matching. The paired-domination number of $G$, denoted by $\gamma_{\mathrm{pr}}(G)$, is the minimum cardinality of a paired-dominating set of $G$. The graph $G$ is paired-domination vertex critical if for every vertex $v$ of $G$ that is not adjacent to a vertex of degree one, $\gamma_{\mathrm{pr}}(G-v)<\gamma_{\mathrm{pr}}(G)$. We characterize the connected graphs with minimum degree one that are paired-domination vertex critical and we obtain sharp bounds on their maximum diameter. We provide an example which shows that the maximum diameter of a paired-domination vertex critical graph is at least $\frac{3}{2}\left(\gamma_{\mathrm{pr}}(G)-2\right)$. For $\gamma_{\mathrm{pr}}(G) \leqslant 8$, we show that this lower bound is precisely the maximum diameter of a paired-domination vertex critical graph.


Keywords: paired-domination, vertex critical, bounds, diameter
MSC 2010: 05C69

## 1. Introduction

Domination and its variations in graphs are now well studied. The literature on this subject has been surveyed and detailed in the two books by Haynes, Hedetniemi, and Slater [11], [12]. Brigham, Chinn, and Dutton [1] began the study of vertex domination critical graphs where the domination number decreases by the removal of any vertex. Further properties of these graphs were explored in [7], [8], [21], [22], [23], [24], but they have not been characterized. In [10] the same concept was introduced for total domination. In this paper we investigate paired-domination vertex critical graphs first studied by Edwards [5].

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A matching $M$ in a graph $G$ is a set of independent edges in $G$. The number of edges in a maximum matching of $G$ is called the matching number of $G$ which we denote by $\alpha^{\prime}(G)$. A vertex of $G$ incident with an edge of the matching $M$ is said to be matched by $M$, or simply $M$-matched. The matching $M$ is called a perfect matching in $G$ if every vertex of $G$ is $M$-matched. A paired-dominating set, abbreviated PDS, of a graph $G$ is a set $S$ of vertices of $G$ such that every vertex is adjacent to some vertex in $S$ and the subgraph $G[S]$ induced by $S$ contains a perfect matching $M$ (not necessarily induced). Two vertices joined by an edge of $M$ are said to be paired and are also called partners in $S$. Every graph without isolated vertices has a PDS since the end-vertices of any maximal matching form such a set. The paired-domination number of $G$, denoted by $\gamma_{\mathrm{pr}}(G)$, is the minimum cardinality of a PDS. A PDS of cardinality $\gamma_{\mathrm{pr}}(G)$ we call a $\gamma_{\mathrm{pr}}(G)$-set. Paired-domination was introduced by Haynes and Slater $[14],[15]$ as a model for assigning backups to guards for security purposes, and is studied, for example, in [2], [3], [4], [6], [9], [13], [16], [17], [18], [19], [20] and elsewhere.

For notation and graph theory terminology we in general follow [11]. Specifically, let $G=(V, E)$ be a graph with vertex set $V$ of order $n$ and edge set $E$. The open neighborhood of $v \in V$ is $N(v)=\{u \in V: u v \in E\}$ and the closed neighborhood of $v$ is $N[v]=\{v\} \cup N(v)$. For a set $S \subseteq V, N(S)=\bigcup_{v \in S} N(v)$ and $N[S]=N(S) \cup S$. For sets $S, T \subseteq V$, we say that $S$ dominates $T$ if $T \subseteq N[S]$ and that $S$ paired-dominates $T$ if $S$ dominates $T$ in $G$ and $G[S]$ contains a perfect matching.

We denote the degree of a vertex $v$ in $G$ by $d_{G}(v)$, or simply by $d(v)$ if the graph $G$ is clear from context. The minimum and maximum degrees of the graph $G$ are denoted by $\delta(G)$ and $\Delta(G)$, respectively. An end-vertex is a vertex of degree one and a support vertex is one that is adjacent to an end-vertex. The set of support vertices in $G$ is denoted by $S(G)$, while the complement of $G$ is denoted by $\bar{G}$. Two vertices at maximum distance apart in $G$ are called diametrical vertices of $G$.

We call a vertex $v \in V$ paired-critical if $\gamma_{\mathrm{pr}}(G-v)<\gamma_{\mathrm{pr}}(G)$. Since paireddomination is undefined for a graph with isolated vertices, we say that a graph $G$ is paired-domination-vertex-critical, or $\gamma_{\mathrm{pr}}$-vertex-critical, if every vertex of $V \backslash S(G)$ is paired-critical. If $G$ is $\gamma_{\mathrm{pr}}$-vertex-critical and $\gamma_{\mathrm{pr}}(G)=k$, then we say that $G$ is $k-\gamma_{\mathrm{pr}}$-vertex-critical. For example, the 5 -cycle is $4-\gamma_{\mathrm{pr}}$-vertex-critical. A graph is $\gamma_{\mathrm{pr}}$-vertex-critical if and only if each of its components is $\gamma_{\mathrm{pr}}$-vertex-critical. Also, $K_{2}$ is trivially $2-\gamma_{\mathrm{pr}}$-vertex-critical. So henceforth we consider only connected graphs of order at least 3. The removal of a vertex can decrease the paired-domination number by at most two. Hence:

Observation 1. If $G$ is a $\gamma_{\mathrm{pr}}$-vertex-critical graph, then $\gamma_{\mathrm{pr}}(G-v)=\gamma_{\mathrm{pr}}(G)-2$ for every $v \in V(G) \backslash S(G)$. Furthermore, a $\gamma_{\mathrm{pr}}(G-v)$-set contains no neighbour of $v$.

In Section 2 we characterize the connected $\gamma_{\mathrm{pr}}$-vertex-critical graphs that have an end-vertex, and we obtain sharp bounds on their maximum diameter. In Section 3 we show that the maximum diameter of a $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph is at least $\frac{3}{2}(k-2)$. For $k \leqslant 8$ we show in Section 4 that this maximum diameter is achieved.

## 2. Graphs with end-vertices

We can readily characterize the $\gamma_{\text {pr-vertex-critical graphs with end-vertices. For }}$ this purpose, we recall that the corona $\operatorname{cor}(H)$ of a graph $H$ (also denoted $H \circ K_{1}$ in [11]) is the graph obtained from $H$ by adding a pendant edge to each vertex of $H$.

Theorem 2. Let $G$ be a connected graph of order at least 3 with at least one endvertex. Then $G$ is $\gamma_{\mathrm{pr}}$-vertex-critical if and only if $G=\operatorname{cor}(H)$ for some connected graph $H$ satisfying $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$.

Proof. First we consider sufficiency. Suppose $G=\operatorname{cor}(H)$ for some connected graph $H$ satisfying $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$. Since every minimal PDS contains every support vertex in the graph, and since $S(G)=V(H)$,

$$
\begin{equation*}
\gamma_{\mathrm{pr}}(G)=2 \alpha^{\prime}(H)+2\left(|V(H)|-2 \alpha^{\prime}(H)\right)=2\left(|V(H)|-\alpha^{\prime}(H)\right) \tag{1}
\end{equation*}
$$

To show that $G$ is $\gamma_{\mathrm{pr}}$-vertex-critical, let $u \in V(G)-S(G)$. Then $d_{G}(u)=1$ and $u$ is adjacent to a unique vertex $v$ of $H$. Let $M_{v}$ be a maximum matching in $H-v$. Then $\left|M_{v}\right|=\alpha^{\prime}(H-v)=\alpha^{\prime}(H)$. Let $V_{1}$ be the set of vertices in $H$ incident with an edge of $M_{v}$ and let $V_{2}=V(H) \backslash\left(V_{1} \cup\{v\}\right)$. Then $\left|V_{1}\right|=2 \alpha^{\prime}(H)$, $\left|V_{2}\right|=|V(H)|-2 \alpha^{\prime}(H)-1$ and $V_{2}$ is an independent set. Let $V_{2}^{\prime}$ be the set of endvertices of $G$ dominated by $V_{2}$; thus, $\left|V_{2}^{\prime}\right|=\left|V_{2}\right|$. Notice that since $H$ is a connected graph, $v$ is adjacent to at least one other vertex of $H$. Therefore, $(V(H) \backslash\{v\}) \cup V_{2}^{\prime}$ is a PDS of $G-u$, so that
(2) $\gamma_{\mathrm{pr}}(G-u) \leqslant|V(H)|-1+\left|V_{2}\right|=2\left(|V(H)|-\alpha^{\prime}(H)\right)-2=\gamma_{\mathrm{pr}}(G)-2 \leqslant \gamma_{\mathrm{pr}}(G-u)$.

Hence equality holds throughout the inequality chain (2) and by Observation $1, G$ is $\gamma_{\mathrm{pr}}$-vertex-critical. This establishes sufficiency.

Next we consider necessity. Suppose that $G$ is a $\gamma_{\mathrm{pr}}$-vertex-critical graph that contains an end-vertex. Let $v^{\prime}$ be an end-vertex and let $v$ be its neighbor. Suppose there exists $w \in N(v) \backslash\left\{v^{\prime}\right\}$ with $w \notin S(G)$. Then by Observation 1, there is a $\gamma_{\mathrm{pr}}(G-w)$-set not containing $v$, but since $v$ is a support vertex in $G-w$, the vertex $v$ belongs to every $\gamma_{\mathrm{pr}}(G-w)$-set, a contradiction. Thus each vertex in $N(v) \backslash\left\{v^{\prime}\right\}$ is
a support vertex. It follows that $G=\operatorname{cor}(H)$ for some connected graph $H$. Thus, as in $(1), \gamma_{\mathrm{pr}}(G)=2\left(|V(H)|-\alpha^{\prime}(H)\right)$.

It remains for us to show that $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$. Let $v \in V(H)$ and let $u$ be the end-vertex adjacent to $v$. Let $M_{v}$ be a maximum matching in $H-v$. Then $\left|M_{v}\right|=\alpha^{\prime}(H-v)$. Let $V_{1}$ be the set of vertices in $H$ incident with an edge of $M_{v}$ and let $V_{2}=V(H) \backslash\left(V_{1} \cup\{v\}\right)$. Then $\left|V_{1}\right|=2 \alpha^{\prime}(H-v)$, $\left|V_{2}\right|=|V(H)|-2 \alpha^{\prime}(H-v)-1$ and $V_{2}$ is an independent set. Let $V_{2}^{\prime}$ be the set of end-vertices dominated by $V_{2}$; thus, $\left|V_{2}^{\prime}\right|=\left|V_{2}\right|$. Let $S=(V(H) \backslash\{v\}) \cup V_{2}^{\prime}$. Then $S$ is a minimum PDS of $G-u$. Hence, $\gamma_{\mathrm{pr}}(G-u)=|S|=|V(H)|-1+$ $\left|V_{2}\right|=2\left(|V(H)|-\alpha^{\prime}(H-v)\right)-2$. However, since $G$ is a $\gamma_{\mathrm{pr}}$-vertex-critical graph, $\gamma_{\mathrm{pr}}(G-u)=\gamma_{\mathrm{pr}}(G)-2=2\left(|V(H)|-\alpha^{\prime}(H)\right)-2$. Consequently, $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$, as desired.

We remark that there are infinite families of connected graphs $H$ satisfying $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$. For example, let $H$ be any hamiltonian graph of odd order. We observe further that the diameter of such graphs $H$ cannot be too large.

Proposition 3. If $H$ is a connected graph satisfying $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$, then every maximum matching in $H-v$ matches every neighbor of $v$. In particular, $H$ is a 2-edge-connected graph.

Proof. Suppose that $H-v$ contains a maximum matching $M$ that does not match a neighbor $u$ of $v$. Then $M \cup\{u v\}$ is a matching in $H$, and so $\alpha^{\prime}(H) \geqslant$ $|M|+1=\alpha^{\prime}(H-v)+1$, a contradiction. Hence every maximum matching in $H-v$ matches every neighbor of $v$.

Suppose that $H$ has a bridge $e=u v$. Let $H_{u}$ and $H_{v}$ be the two components of $H-e$, where $u \in V\left(H_{u}\right)$ and $v \in V\left(H_{v}\right)$. Then $\alpha^{\prime}(H) \geqslant \alpha^{\prime}\left(H_{u}\right)+\alpha^{\prime}\left(H_{v}\right)$. Since every maximum matching of $H-u$ matches every neighbor of $u$, the vertex $v$ is matched in every maximum matching of $H-u$. This implies that $\alpha^{\prime}\left(H_{v}-v\right)=\alpha^{\prime}\left(H_{v}\right)-1$. But then $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)=\alpha^{\prime}\left(H_{u}\right)+\alpha^{\prime}\left(H_{v}-v\right)=\alpha^{\prime}\left(H_{u}\right)+\alpha^{\prime}\left(H_{v}\right)-1$, producing a contradiction. Hence, $H$ is 2-edge-connected.

Proposition 4. If $H$ is a connected graph of order $n$ satisfying $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$, then $\operatorname{diam}(H) \leqslant \frac{1}{2}(n-1)$.

Proof. We proceed by induction on the number of blocks $b(H)$ in $H$. Suppose $b(H)=1$. Let $u$ and $v$ be two diametrical vertices in $H$, and so $\operatorname{diam}(H)=d(u, v)$. Since $H$ is 2-connected, every two vertices of $H$ lie on a common cycle of $H$. In particular, there is a cycle $C$ containing $u$ and $v$. Hence, $|V(C)| \geqslant 2 d(u, v)=$ $2 \operatorname{diam}(H)$. On the one hand, if $|V(C)| \geqslant 2 \operatorname{diam}(H)+1$, then $n \geqslant|V(C)| \geqslant$
$2 \operatorname{diam}(H)+1$. On the other hand, suppose $|V(C)|=2 \operatorname{diam}(H)$. Since $\alpha^{\prime}(H)=$ $\alpha^{\prime}(H-w)$ for every $w \in V(H)$, the graph $H$ is not a hamiltonian graph of even order. Thus $H$ contains at least one vertex not on $C$, implying that $n \geqslant|V(C)|+1=$ $2 \operatorname{diam}(H)+1$. In both cases, $n \geqslant 2 \operatorname{diam}(H)+1$, or, equivalently, $\operatorname{diam}(H) \leqslant$ $\frac{1}{2}(n-1)$. This establishes the base case.

Assume that $b \geqslant 1$ and that if $H^{\prime}$ is a connected graph of order $n^{\prime}$ satisfying $b\left(H^{\prime}\right) \leqslant b$ and $\alpha^{\prime}\left(H^{\prime}\right)=\alpha^{\prime}\left(H^{\prime}-v\right)$ for every $v \in V\left(H^{\prime}\right)$, then $\operatorname{diam}\left(H^{\prime}\right) \leqslant \frac{1}{2}\left(n^{\prime}-1\right)$. Let $H$ be a connected graph of order $n$ satisfying $b(H)=b+1$ and $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$. Let $B$ be an end-block of $H$ and $v$ the unique cut-vertex of $H$ contained in $B$. Let $F=H-(V(B) \backslash\{v\})$. Then $F$ is a connected graph satisfying $b(F)=b$. We proceed further with three claims.

Claim 1. $\alpha^{\prime}(H)=\alpha^{\prime}(B)+\alpha^{\prime}(F)$.
Proof. We show first that $\alpha^{\prime}(B)=\alpha^{\prime}(B-v)$. Suppose $\alpha^{\prime}(B)>\alpha^{\prime}(B-v)$. Then $\alpha^{\prime}(B)=\alpha^{\prime}(B-v)+1$ and every maximum matching of $B$ matches the vertex $v$. Let $e=u v$ be an edge of such a maximum matching $M_{B}$ of $B$. Then $M_{B} \backslash\{e\}$ is a maximum matching of $B-v$ that does not match the vertex $u$. But every maximum matching of $B-v$ can be extended to a maximum matching of $H-v$ by adding to it the edges of a maximum matching of $F-v$. Hence we have shown that there is a maximum matching of $H-v$ that does not match the neighbor $u$ of $v$, contradicting Proposition 3. Hence, $\alpha^{\prime}(B)=\alpha^{\prime}(B-v)$. Similarly, $\alpha^{\prime}(F)=\alpha^{\prime}(F-v)$. Thus since the graph $H$ is $\gamma_{\mathrm{pr}}$-vertex-critical, $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)=\alpha^{\prime}(B-v)+\alpha^{\prime}(F-v)=$ $\alpha^{\prime}(B)+\alpha^{\prime}(F)$, as claimed.

Claim 2. $\operatorname{diam}(F) \leqslant \frac{1}{2}(|V(F)|-1)$.
Proof. Let $w \in V(F)$. Then, by Claim 1, $\alpha^{\prime}(B)+\alpha^{\prime}(F)=\alpha^{\prime}(H)=\alpha^{\prime}(H-w) \leqslant$ $\alpha^{\prime}(B)+\alpha^{\prime}(F-w)$, and so $\alpha^{\prime}(F) \leqslant \alpha^{\prime}(F-w)$. Consequently, $F$ is a connected graph with $b(F)=b$ such that $\alpha^{\prime}(F)=\alpha^{\prime}(F-w)$ for every vertex $w \in V(F)$. Applying the inductive hypothesis to $F$, we conclude that $\operatorname{diam}(F) \leqslant \frac{1}{2}(|V(F)|-1)$.

The proof of the following claim is similar to the proof of Claim 2 and is omitted.
Claim 3. $\operatorname{diam}(B) \leqslant \frac{1}{2}(|V(B)|-1)$.
The desired upper bound on the diameter of $H$ now follows readily from Claims 2 and 3 and the observations that $\operatorname{diam}(H) \leqslant \operatorname{diam}(B)+\operatorname{diam}(F)$ and $|V(B)|+$ $|V(F)|=n+1$. This completes the proof of Proposition 4.

As a consequence of Theorem 2 and Propositions 3 and 4, we have the following results.

Theorem 5. No tree is $\gamma_{\mathrm{pr}}$-vertex-critical.

Theorem 6. If $G$ is a connected $\gamma_{\mathrm{pr}}$-vertex-critical graph with at least one endvertex, then $\operatorname{diam}(G) \leqslant \frac{1}{2}\left(\gamma_{\mathrm{pr}}(G)+2\right)$, and this bound is sharp.

Proof. By Theorem 2, $G=\operatorname{cor}(H)$ for some connected graph $H$ satisfying $\alpha^{\prime}(H)=\alpha^{\prime}(H-v)$ for every $v \in V(H)$. Hence, $\operatorname{diam}(G)=2+\operatorname{diam}(H)$. Suppose $\gamma_{\mathrm{pr}}(G)=k$. Since $H$ does not have a perfect matching, $|V(H)| \leqslant k-1$. By Proposition 4, $\operatorname{diam}(H) \leqslant \frac{1}{2}(|V(H)|-1) \leqslant \frac{1}{2}(k-2)$. Hence, $\operatorname{diam}(G)=2+$ $\operatorname{diam}(H) \leqslant 2+\frac{1}{2}(k-2)=\frac{1}{2}(k+2)$. To see that this bound is sharp, take $H=C_{k-1}$.

## 3. $\gamma_{\mathrm{pr}}$-VERTEX-CRITICAL GRAPHS WITH LARGE DIAMETER

In this section we provide a construction of $\gamma_{\mathrm{pr}}$-vertex-critical graphs with large diameter. First we give a way of constructing a $\gamma_{\mathrm{pr}}$-vertex-critical graph from two smaller $\gamma_{\text {pr-vertex-critical graphs. }}$

Lemma 7. Let $F$ and $H$ be a $j-\gamma_{\mathrm{pr}}$-vertex-critical and a $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph, respectively, with minimum degrees at least two, and let $G$ be a graph formed by identifying a vertex of $F$ with a vertex of $H$. If $\gamma_{\mathrm{pr}}(G)=j+k-2$, then $G$ is $\gamma_{\mathrm{pr}}$-vertex-critical.

Proof. Note that since $\delta(F) \geqslant 2$ and $\delta(H) \geqslant 2, S(G)=\emptyset$. Label the identified vertex $v$. Let $u \in V(G)$. Without loss of generality, $u \in V(F)$. Since $F$ is $j$ - $\gamma_{\mathrm{pr}}$-vertex-critical, $\gamma_{\mathrm{pr}}(F-u)=j-2$. If $u \neq v$, then every $\gamma_{\mathrm{pr}}(F-u)$-set dominates $v$ and can be extended to a PDS of $G-u$ by adding to it $\gamma_{\mathrm{pr}}(H-v)=k-2$ vertices from $H-v$. Hence, $\gamma_{\mathrm{pr}}(G-u) \leqslant j-2+k-2=\gamma_{\mathrm{pr}}(G)-2$. If $u=v$, then $\gamma_{\mathrm{pr}}(G-v)=\gamma_{\mathrm{pr}}(F-v)+\gamma_{\mathrm{pr}}(H-v)=j-2+k-2=\gamma_{\mathrm{pr}}(G)-2$. Thus, $\gamma_{\mathrm{pr}}(G-u)<\gamma_{\mathrm{pr}}(G)$ and $G$ is $\gamma_{\mathrm{pr}}$-vertex-critical.

Next we establish a lower bound on the maximum diameter of a $k-\gamma_{\mathrm{pr}}$-vertexcritical graph. For this purpose, following the notation of Goddard et al. [10] we define a graph as pointed if there are two designated diametrical vertices called Left and right. Then, for two pointed graphs $G$ and $H$, we define $G \circ H$ as the pointed graph obtained by identifying and undesignating the RIGHT-vertex from $G$ and the LEFT-vertex from $H$. Note that the operator $\circ$ is associative.

For a graph $G=(V, E)$ with $\operatorname{diam}(G)=d$ we also define the following subsets of $V$, and use this notation throughout the rest of the paper. Fix a diametrical
vertex $v$ of $G$. For $i=0,1, \ldots, d$, define

$$
\begin{equation*}
V_{i}=\{u \in V: d(u, v)=i\}, \quad V_{\leqslant i}=\bigcup_{j=0}^{i} V_{j} \quad \text { and } \quad V_{\geqslant i}=\bigcup_{j=i}^{d} V_{j} . \tag{3}
\end{equation*}
$$

Note that $V_{0}=\{v\}$ and $V_{1}=N(v)$.

Theorem 8. For every even integer $k \geqslant 4$ there exists a connected $k$ - $\gamma_{\mathrm{pr}}$-vertexcritical graph of diameter $\frac{3}{2}(k-2)$.

Proof. We begin by constructing a $4-\gamma_{\mathrm{pr}}$-vertex-critical graph with diameter 3 . Let $H_{1}$ be a copy of $P_{4}$ and let $H_{2}$ be a copy of $\overline{H_{1}}$. Let $F$ be the pointed graph obtained from $H_{1} \cup H_{2}$ by adding all edges between $H_{1}$ and $H_{2}$ except for a perfect matching between the corresponding vertices of $H_{1}$ and $H_{2}$, and then adding two new vertices, LEFT and RIGHT, such that LEFT is joined to every vertex in $H_{1}$ and RIGHT is joined to every vertex in $H_{2}$. The graph $F$ is shown in Fig. 1 where for clarity we omit the edges between $H_{1}$ and $H_{2}$. Then $F$ is $4-\gamma_{\mathrm{pr}}$-vertex-critical with diameter 3.


Figure 1. The $4-\gamma_{\mathrm{pr}}$-vertex-critical graph $F$ of diameter 3
For $q \geqslant 1$ define the pointed graph $G_{q}=F \circ F \circ \ldots \circ F$ for $q$ copies of $F$. Then $\operatorname{diam}\left(G_{q}\right)=3 q$. We show that $G_{q}$ is a $2(q+1)$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph. We proceed by induction on $q$. When $q=1$, then $G_{q}=F$ which is a $4-\gamma_{\mathrm{pr}}$-vertex-critical graph. This establishes the base case. Assume then that $q \geqslant 2$ and that $G_{q^{\prime}}$ is a $2\left(q^{\prime}+1\right)-\gamma_{\mathrm{pr}}$-vertex-critical graph for $1 \leqslant q^{\prime}<q$. We now consider the graph $G_{q}$.

The graph $G_{q}$ is the pointed graph obtained from the pointed graphs $F$ and $G_{q-1}$; that is, $G_{q}=F \circ G_{q-1}$, where $F$ is a $4-\gamma_{\mathrm{pr}}$-vertex-critical graph and, by induction, $G_{q-1}$ is a $2 q-\gamma_{\mathrm{pr}}$-vertex-critical graph. Let $F_{1}$ denote the first copy of $F$ in $G_{q}$, and let $v$ and $w$ denote the LEFT-vertex and RIGHT-vertex from $F_{1}$.

The vertex $v$ is a diametrical vertex of $G_{q}$. Let $d=\operatorname{diam}\left(G_{q}\right)=3 q$. As in (3), $V_{0}=\{v\}$ and $V_{1}=N(v)$. Further, $V_{3}=\{w\}$, while $V_{2}$ is the neighborhood of $w$ in $F_{1}$ and $V_{4}$ is the neighborhood of $w$ in $G_{q-1}$.

Among all $\gamma_{\mathrm{pr}}\left(G_{q}\right)$-sets, let $S$ be one which contains as few vertices of $V_{\leqslant 2}$ as possible. To dominate $V_{0}$, we have that $\left|S \cap V_{\leqslant 2}\right| \geqslant 2$. Suppose that $\left|S \cap V_{\leqslant 2}\right| \geqslant 3$.

Then $\left|S \cap V_{\leqslant 3}\right| \geqslant 4$. Note that $V_{4} \nsubseteq S$, otherwise, if $x, x^{\prime} \in V_{4}$ are partners in $S$, then $S \backslash\left\{x, x^{\prime}\right\}$ is a PDS of $G_{q}$ of smaller cardinality than $S$, which is impossible. Replacing the vertices in $S \cap V_{\leqslant 3}$ by the two central vertices of the $P_{4}$ in $G_{q}\left[V_{1}\right]$ and the vertex $w$, and then adding to the resulting set a neighbor of $w$ from $V_{4}$ (to serve as a partner of $w$ ) produces a new $\gamma_{\mathrm{pr}}\left(G_{q}\right)$-set that contains fewer vertices from $V_{\leqslant 2}$ than does $S$, contradicting our choice of $S$. Hence, $\left|S \cap V_{\leqslant 2}\right|=2$. If follows that $S \cap V_{\geqslant 3}$ is a PDS of $G_{q-1}$ and that $\left|S \cap V_{\geqslant 3}\right|=\gamma_{\mathrm{pr}}\left(G_{q}\right)-2$. Hence, $\gamma_{\mathrm{pr}}\left(G_{q-1}\right) \leqslant \gamma_{\mathrm{pr}}\left(G_{q}\right)-2$. Every $\gamma_{\mathrm{pr}}\left(G_{q-1}\right)$-set can easily be extended to a PDS of $G_{q}$ by adding to it two vertices (namely, the two central vertices of the $P_{4}$ in $G_{q}\left[V_{1}\right]$ ), and so $\gamma_{\mathrm{pr}}\left(G_{q}\right) \leqslant \gamma_{\mathrm{pr}}\left(G_{q-1}\right)+2$. Consequently, $\gamma_{\mathrm{pr}}\left(G_{q}\right)=\gamma_{\mathrm{pr}}\left(G_{q-1}\right)+2=\gamma_{\mathrm{pr}}(F)+\gamma_{\mathrm{pr}}\left(G_{q-1}\right)-2$. Hence, by Lemma 7 , $G_{q}$ is $\gamma_{\mathrm{pr}}$-vertex-critical. By induction, $\gamma_{\mathrm{pr}}\left(G_{q-1}\right)=2 q$, and so $G_{q}$ is a $k$ - $\gamma_{\mathrm{pr}}$-vertexcritical graph where $k=2(q+1)$ with $\operatorname{diam}\left(G_{q}\right)=3 q=\frac{3}{2}(k-2)$.

## 4. Bounds on the Diameter

In this section we establish bounds on the diameter of a connected $k$ - $\gamma_{\mathrm{pr}}$-vertexcritical graph. First we mention a sufficient condition for a graph not to be $\gamma_{\mathrm{pr}}{ }^{-}$ vertex-critical. (We assume in what follows that $G$ has no end-vertex, for otherwise we have the upper bound given in Theorem 6.)

Proposition 9 ([5, Proposition 5.4]). If a graph $G$ has nonadjacent vertices $u$ and $v$ with $N(u) \subseteq N(v)$, then $G$ is not a $\gamma_{\mathrm{pr}}$-vertex-critical graph.

We provide next a trivial upper bound on the diameter of a $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph. Throughout this section, for a graph $G=(V, E)$ and a vertex $x \in V$, we let $S_{x}$ denote a $\gamma_{\mathrm{pr}}(G-x)$-set.

Proposition 10. The diameter of a connected $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical $G$ graph with $\operatorname{diam}(G)=d$ is at most $2 k-8+(d \bmod 4)$.

Proof. Let $v$ be a diametrical vertex of $G$ and let $d=\operatorname{diam}(G)$. As in (3), $V_{0}=\{v\}$ and $V_{1}=N(v)$. By Observation $1,\left|S_{v}\right|=k-2$ and $S_{v} \cap V_{1}=\emptyset$. Hence to dominate $V_{1},\left|S_{v} \cap V_{2}\right| \geqslant 1$. In fact, by Proposition $9,\left|S_{v} \cap V_{2}\right| \geqslant 2$. Thus, $S=S_{v} \cup\left\{v, v_{1}\right\}$ is a $\gamma_{\mathrm{pr}}(G)$-set for any $v_{1} \in V_{1}$ and $\left|S \cap\left(V_{0} \cup V_{1} \cup V_{2}\right)\right| \geqslant 4$. For any $i \geqslant 3,\left|S \cap\left(V_{i} \cup \ldots \cup V_{i+3}\right)\right| \geqslant 2$. It follows that if $d=2+4 j+r$ where $0 \leqslant r \leqslant 3$, then $k=|S| \geqslant 4+2 j$ if $r \in\{0,1\}$ while $k \geqslant 4+2 j+2$ if $r \in\{2,3\}$. The desired result now follows from simple algebra.

Since $d \bmod 4 \in\{0,1,2,3\}$, as an immediate consequence of Proposition 10 we have the following result.

Corollary 11. The diameter of a connected $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph $G$ is at most $2 k-5$ with inequality if $\operatorname{diam}(G) \not \equiv 3(\bmod 4)$.

As an immediate consequence of Theorem 8, we have the following result.
Corollary 12. The maximum diameter of a connected $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph is at least $\frac{3}{2}(k-2)$.

Next we establish a sharp upper bound on the diameter of a connected $k-\gamma_{\mathrm{pr}}-$ vertex-critical graph for small $k$. Recall that for a graph $G=(V, E)$ and sets $S, T \subseteq V$, we say that $S$ paired-dominates $T$ if $S$ dominates $T$ in $G$ and $G[S]$ contains a perfect matching.

Theorem 13. For $k \leqslant 8$, the diameter of a connected $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph is at most $\frac{3}{2}(k-2)$.

Proof. Let $G=(V, E)$ be a connected $k$ - $\gamma_{\mathrm{pr}}$-vertex-critical graph. If $\delta(G)=1$, then the upper bounds follow from Theorem 6. Hence we may assume in what follows that $\delta(G) \geqslant 2$. We will show that the diameter of $G$ is at most the value given in Tab. 1.

| $k$ | 4 | 6 | 8 |
| :---: | :---: | :---: | :---: |
| $\operatorname{diam}(G)$ | 3 | 6 | 9 |

Table 1. The maximum value $\operatorname{od} \operatorname{diam}(G)$ for $k \leqslant 8$.
If $k=4$, then the upper bound follows from Corollary 11. Hence we may assume $\delta(G) \geqslant 2$ and $k \geqslant 6$. Let $v$ be a diametrical vertex of $G$ and let $d=\operatorname{diam}(G)$. For $S, T \subseteq V$ we write $S \succ_{\text {pr }} T$ if $S$ paired-dominates $T$ in $G$. Furthermore, we write $S \mapsto_{\mathrm{pr}} T$ if $S \cap T \succ_{\mathrm{pr}} T$. As before, for $x \in V$, let $S_{x}$ be a $\gamma_{\mathrm{pr}}(G-x)$-set.

Suppose that $k=6$ and assume that $d \geqslant 7$. Again using the notation defined in (3), let $u \in V_{1}$; then $\left|S_{u}\right|=4$. To paired-dominate $V_{0} \cup V_{3} \cup V_{4} \cup V_{7}$, it follows that $d=7$ and $\left|S_{u} \cap V_{j}\right|=1$ for $j \in\{1,2,5,6\}$. Thus, $\left|S_{u} \cap V_{\geqslant 4}\right|=2$ and $S_{u} \mapsto_{\mathrm{pr}} V_{\geqslant 4}$. By symmetry, for $w \in V_{6}$ it follows that $\left|S_{w} \cap V_{\leqslant 3}\right|=2$ and $S_{w} \mapsto_{\mathrm{pr}} V_{\leqslant 3}$. Therefore, $\left(S_{u} \cap V_{\geqslant 4}\right) \cup\left(S_{w} \cap V_{\leqslant 3}\right)$ is a PDS of $G$ of cardinality 4, which contradicts $\gamma_{\mathrm{pr}}(G)=6$. Hence, if $k=6$, then $d \leqslant 6$, as desired.

Suppose that $k=8$ and assume that $d \geqslant 10$. Let $u \in V_{1}$; then $\left|S_{u}\right|=6$. To paired-dominate $V_{\leqslant 2}$, we must have $\left|S_{u} \cap V_{\leqslant 2}\right| \geqslant 2$, while to paired-dominate $V_{\geqslant 8}$, we must have $\left|S_{u} \cap V_{\geqslant 8}\right| \geqslant 2$. Hence, to paired-dominate $V_{4} \cup V_{5} \cup V_{6}$, we must have $\left|S_{u} \cap V_{5}\right| \geqslant 1$ and $\left|S_{u} \cap\left(V_{4} \cup V_{5} \cup V_{6}\right)\right| \geqslant 2$. Hence, $\left|S_{u} \cap V_{\leqslant 2}\right|=2,\left|S_{u} \cap V_{\geqslant 8}\right|=2$, $\left|S_{u} \cap V_{5}\right| \geqslant 1$ and $\left|S_{u} \cap\left(V_{4} \cup V_{5} \cup V_{6}\right)\right|=2$. In particular, $\left|S_{u} \cap V_{\geqslant 4}\right|=4$ and $S_{u} \mapsto_{\mathrm{pr}} V_{\geqslant 4}$.

Let $w \in V_{9}$. By symmetry, $\left|S_{w} \cap V_{\leqslant 2}\right|=2,\left|S_{w} \cap V_{\geqslant 8}\right|=2,\left|S_{w} \cap V_{5}\right| \geqslant 1$ and $\left|S_{w} \cap\left(V_{4} \cup V_{5} \cup V_{6}\right)\right|=2$. In particular, $\left|S_{w} \cap V_{\leqslant 6}\right|=4$ and $S_{w} \mapsto_{\text {pr }} V_{\leqslant 6}$. If $S_{u} \cap V_{6}=\emptyset$,
then $S_{u} \mapsto_{\mathrm{pr}} V_{\geqslant 7}$, and $\left(S_{w} \cap V_{\leqslant 6}\right) \cup\left(S_{u} \cap V_{\geqslant 8}\right)$ is a PDS of $G$ of cardinality 6, which contradicts $\gamma_{\mathrm{pr}}(G)=8$. Thus we may assume that $\left|S_{u} \cap V_{6}\right|=1$, and so $\left|S_{u} \cap V_{5}\right|=1$; similarly, $\left|S_{w} \cap V_{4}\right|=\left|S_{w} \cap V_{5}\right|=1$.

Let $x \in V_{5}$. Then, as before, $\left|S_{x} \cap V_{\leqslant 2}\right|=2$ and $\left|S_{x} \cap V_{\geqslant 8}\right|=2$. Suppose there is another vertex in $V_{5}$. Then $\left|S_{x} \cap V_{5}\right| \geqslant 1$ and $\left|S_{x} \cap\left(V_{4} \cup V_{5} \cup V_{6}\right)\right|=2$. Without loss of generality, $S_{x} \cap V_{4}=\emptyset$, and so $S_{x} \mapsto_{\mathrm{pr}} V_{\leqslant 3}$. Therefore $\left(S_{x} \cap V_{\leqslant 2}\right) \cup\left(S_{u} \cap V_{\geqslant 4}\right) \succ_{\text {pr }} V$, which contradicts $\gamma_{\mathrm{pr}}(G)=8$. Hence there is no other vertex in $V_{5}$. But then $S_{x}$ contains at least one vertex in each of the sets $V_{0} \cup V_{1}$ (to dominate $V_{0}$ ), $V_{3} \cup V_{4}$ (to dominate $V_{4}$ ), $V_{6} \cup V_{7}$ (to dominate $V_{6}$ ), and $V_{9} \cup V_{10}$ (to dominate $V_{10}$ ). Thus, $S_{x}$ contains four vertices that are pairwise nonadjacent, implying that $\left|S_{x}\right| \geqslant 8$, a contradiction. Hence, if $k=8$, then $d \leqslant 9$, as desired.

We close with the following question about the maximum diameter of a connected $\gamma_{\mathrm{pr}}$-vertex-critical graph.

Question 1. If $G$ is a connected $\gamma_{\mathrm{pr}}$-vertex-critical graph, then is it true that

$$
\operatorname{diam}(G) \leqslant \frac{3}{2}\left(\gamma_{\mathrm{pr}}(G)-2\right) ?
$$

Note that by Theorem 13, Question 1 is true for $\gamma_{\mathrm{pr}}(G) \leqslant 8$. By Corollary 12, if Question 1 is true, then this bound is sharp.

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