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## **RESULTS ON F-CONTINUOUS GRAPHS**

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Abstract. For any nontrivial connected graph F and any graph G, the F-degree of a vertex v in G is the number of copies of F in G containing v. G is called F-continuous if and only if the F-degrees of any two adjacent vertices in G differ by at most 1; G is F-regular if the F-degrees of all vertices in G are the same. This paper classifies all  $P_4$ -continuous graphs with girth greater than 3. We show that for any nontrivial connected graph F other than the star  $K_{1,k}$ ,  $k \ge 1$ , there exists a regular graph that is not F-continuous. If F is 2-connected, then there exists a regular F-continuous graph that is not F-regular.

*Keywords*: continuous, *F*-continuous, *F*-regular, regular graph *MSC 2010*: 05C12, 05C78

#### 1. INTRODUCTION

Chartrand et al. in [1] consider the general case of integer-valued functions f defined on a metric space of objects associated with a particular graph G. Such a function is *continuous* if and only if  $|f(x) - f(y)| \leq 1$  for every two adjacent elements x and y in the metric space. When the metric space is the vertex set of G, a continuous function defined on V(G) is, in fact, a labeling of the vertices of G with nonnegative integers such that the labels of any two vertices v and u connected with an edge differ by at most 1. Such a labeling is referred to as a *continuous labeling*. Degree-continuous graphs provide an example of graphs with a certain type of a continuous labeling. A graph G is called *degree-continuous* if  $|\deg(v) - \deg(v')| \leq 1$  for every pair  $\{v, v'\}$  of adjacent vertices of G. For more information on degree-continuous graphs see [5].

This paper is concerned with graphs G = (V, E) together with a different continuous labeling. Given any nontrivial connected graph F, and any vertex  $v \in V(G)$ , the *F*-degree of v in G, denoted F-deg<sub>G</sub>(v), is the number of copies (not necessarily induced) of F in G containing v. Thus, the degree of v, denoted deg<sub>G</sub>(v), and the  $P_2$ -degree of v are the same where  $P_n$  denotes the path on n vertices. When no confusion is possible, we write F-deg(v) instead of F-deg $_G(v)$ , and deg(v) instead of deg $_G(v)$ . A graph G is F-continuous (or F-degree continuous) if and only if the F-degrees of any two adjacent vertices in V(G) differ by at most 1. If, in addition, F-deg(v) = r for all  $v \in V(G)$ , then, G is F-regular of degree r.

Without loss of generality we can assume that G, as well as F, is nontrivial and connected; we do not allow loops or multiple edges. If no copy of F can be found in G, then F-deg(v) = 0 for all  $v \in V(G)$ , and trivially, G is F-continuous and even Fregular. The girth g(G) of a graph G is the minimum among all cycle lengths taken over all cycles in G; the circumference c(G) of G is the length of the largest cycle appearing in G. If G has no cycles, by default  $g(G) = \infty$ . The distance between any two vertices of G is the length of the shortest path between them; the diameter d(G) of G is the largest over all distances between pairs of vertices in G.

The concept of *F*-degree was introduced by Chartrand et al. [2] in 1987; results on *F*-continuous graphs can be found in [3]. In addition to determining all  $P_3$ continuous graphs, Chartrand, Jarrett et al. [3] show that if *G* is *F*-continuous for all nontrivial connected graphs *F*, then,  $G = P_n$  or *G* is regular. However, there are nontrivial connected graphs *F* such that there exists a regular graph *G* that is not *F*-continuous. Certainly, if  $F = K_{1,k}$ ,  $k \ge 2$ , and *G* is an *r*-regular graph, then  $K_{1,k}$ -deg $(v) = (k+1) {r \choose k}$  for every  $v \in V(G)$ . Thus, there is no regular graph which is not  $K_{1,k}$ -continuous. In the case when *F* is a 2-connected graph, however, Chartrand et al. construct a regular graph that is not *F*-continuous [3].

In Section 3, we extend the above result from 2-connected graphs F to all nontrivial connected graphs other than  $K_{1,k}$ ,  $k \ge 2$ , confirming a conjecture in [3]. Furthermore, we show that for every 2-connected graph F, there exists a regular F-continuous graph that is not F-regular. We begin, in Section 2, by classifying all  $P_4$ -continuous graphs that contain no triangles.

#### 2. $P_4$ -continuous graphs

This section is entirely devoted to the case of  $F = P_4$ . All  $P_2$ -continuous graphs have been studied in [5], and all  $P_3$ -continuous graphs have been classified in [3]. We determine all  $P_4$ -continuous graphs with girth greater than 3.

Let H and K denote the graphs on five and four vertices, respectively, shown in Figure 1. Our main result is given below.

**Theorem 2.1.** Let G be a connected  $P_4$ -continuous graph with girth g(G) > 3and minimum degree  $\delta$ . Then, G is isomorphic to one of H,  $P_n$ ,  $K_{1,n}$ , for some integer  $n \ge 1$ , or G is  $\delta$ -regular.



Figure 1. Two  $P_4$ -continuous graphs: acyclic graph H on 5 vertices and graph K on 4 vertices

Before we prove Theorem 2.1, we consider some special cases.

**Lemma 2.2.** Let G be a connected  $P_4$ -continuous graph, and let  $C_3$  denote the cycle on 3 vertices.

- (i) If P<sub>4</sub>-deg(v) = 0 for some vertex v of G, then G ≃ C<sub>3</sub> or K<sub>1,n</sub> for some integer n ≥ 1.
- (ii) If  $P_4$ -deg(v) = 1 for some vertex v of G, then  $G \cong H$  or  $P_n$  for some integer  $n \ge 4$ .
- (iii) If deg(v) = 1 for some vertex v of G and G contains a copy of  $P_4$ , then  $G \cong H$ , K or  $P_n$  for some integer  $n \ge 4$ .

Proof. (i) The distance between any two vertices x and y of G is less than or equal to the length of a path from x to y passing through v. Since G is connected, such a path always exists; it must be that  $d(G) \leq 2$  and the result follows.

(ii) Since v is contained in a copy of  $P_4$ , there exists a vertex u adjacent to vwith deg(u) > 1; i.e.  $\{u, w\} \in E(G)$  for some vertex w other than v. For any vertex x adjacent to v other than u,  $\langle w, u, v, x \rangle$  is a copy of  $P_4$ . Therefore, deg(v) = 1or deg(v) = 2. If deg(v) = 2, no new edges or vertices can be added without contradicting  $P_4$ -deg(v) = 1. It must be that  $G \cong P_4$ . Suppose that deg(v) = 1 and let  $\langle v, u, w, y_1 \rangle$  be the copy of  $P_4$  containing v. Now, no new edges adjacent to w can be added; u can only be adjacent to a new vertex  $y_2$  in which case  $G \cong H$  and no additional edges are present. Otherwise, it must be that  $G \cong P_4$  or a new vertex  $y_2$ is adjacent to  $y_1$ . Again, either  $G \cong P_5$  or there is a new vertex  $y_3$  that can only be adjacent to  $y_2$ . Repeating the same procedure we see that  $G \cong P_n$  for some integer  $n \ge 4$ .

(iii) Let u be the only vertex adjacent to v. Denote by G' the graph with vertex set V(G) - v and edge set  $E(G) - \{v, u\}$ . A copy of  $P_4$  in G that contains v must necessarily contain u as well. Any copy of  $P_4$  that contains u but does not contain v must lie entirely in G'. Therefore,  $P_4 - \deg_{G'}(u) = 0$  or 1. If  $P_4 - \deg_{G'}(u) = 0$ , then, by (i), either  $G' \cong C_3$  and  $G \cong K$ , or  $G' \cong K_{1,n}$  and G is  $P_4$ -continuous and contains a copy of  $P_4$  only if  $G \cong P_4$  or  $G \cong H$ . If  $P_4 - \deg_{G'}(u) = 1$ , then, by (ii),  $G' \cong H$  or

 $G' \cong P_n, n \ge 4$ . It is easy to see that the only way for G to be  $P_4$ -continuous in this case is if  $G \cong H$  or  $G \cong P_n, n \ge 5$ .

Proof of Theorem 2.1. By Lemma 2.2, we may assume that  $\delta \ge 2$  and that G contains a copy of  $P_4$ . Let v be a vertex of G of degree  $\delta$ , and let  $u_1, \ldots, u_{\delta} \in V(G)$  denote the neighbours of v where deg $(u_i) := d_i + 1$ . For each i, let  $u_{i,1}, \ldots, u_{i,d_i} \in V(G)$  denote the  $d_i$  neighbours of  $u_i$  other than v with deg $(u_{i,j}) := d_{i,j} + 1$  for  $j = 1, \ldots, d_i$ . Certainly not all  $u_{i,j}$  have to be distinct. Define

$$c_i := \sum_{j=1}^{d_i} d_{i,j}$$

and without loss of generality assume that  $c_1 \ge c_i$  for  $i = 2, \ldots, \delta$ .

Since G contains no triangles, the  $P_4$ -degree of a vertex in G depends only on the degrees of all vertices of distance two or less from the given vertex. If A denotes the number of copies of  $P_4$  in G that contain both v and  $u_1$ , then

$$P_{4}-\deg(v) = A + \sum_{i=2}^{\delta} c_{i} + (\delta - 2) \sum_{i=2}^{\delta} d_{i},$$
  
$$P_{4}-\deg(u_{1}) \ge A + c_{1}(d_{1} - 1) + c_{1}(\delta - 1) = A + c_{1}(d_{1} + \delta - 2)$$

since each neighbour of  $u_{1,j}$ ,  $j = 1, ..., d_1$ , must have degree at least  $\delta$ . For each  $i = 2, ..., \delta$ ,  $c_1 \ge c_i \ge d_i(\delta - 1)$ , leading to

$$c_1 \geqslant \sum_{i=2}^{\delta} d_i$$

It must be that

$$P_4$$
-deg $(v) \leq A + c_1(\delta - 1) + c_1(\delta - 2) = A + c_1(2\delta - 3)$ 

and since  $d_1 \ge \delta - 1$ ,

$$1 \ge P_4 - \deg(u_1) - P_4 - \deg(v) \ge c_1(d_1 - \delta + 1) \ge d_1(\delta - 1)(d_1 - \delta + 1).$$

The above inequality does not hold when  $d_1 > \delta - 1$  since  $\delta \ge 2$ ; we must have  $d_1 = \delta - 1$  and  $\deg(u_1) = \delta$ . Then,

(1) 
$$P_4 \operatorname{-deg}(v) \leqslant A + c_1(2\delta - 3) \leqslant P_4 \operatorname{-deg}(u_1).$$

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If equality holds in the first part of (1), then  $c_1 = c_i$  for all  $i = 2, ..., \delta$ , and by the same argument as above applied to  $c_i$ ,  $\deg(u_i) = \delta$ . All neighbours of the arbitrary vertex v of degree  $\delta$  must also have degree  $\delta$ , showing that G is  $\delta$ -regular.

Otherwise, equality must hold in the second part of (1). Assume that  $c_k < c_1$  for some  $k, 2 \leq k \leq j$ . Since G is  $P_4$ -continuous,  $P_4$ -deg $(v) = A + c_1(2\delta - 3) - 1$ , and then,

$$A + \sum_{i=2}^{\delta} c_i + (\delta - 2) \sum_{i=2}^{\delta} d_i = A + c_1(2\delta - 3) - 1,$$
  
$$c_k + c_1(\delta - 2) + c_1(\delta - 2) \ge c_1(2\delta - 3) - 1,$$
  
$$c_k \ge c_1 - 1.$$

Thus,  $c_k = c_1 - 1$  and  $c_i = c_1$  for all  $i \neq k$ . Our argument, then, applies to all  $c_i$ ,  $i \neq k$ , and shows that  $\deg(u_i) = \delta$  for all  $i \neq k$ . Then,  $d_k \ge d_i$  for all  $i = 1, \ldots, \delta$  since  $d_i = \delta - 1$  is the smallest possible when  $i \neq k$ . We get,

$$A + c_1(2\delta - 3) - 1 = A + c_1 - 1 + c_1(\delta - 2) + (\delta - 2) \sum_{i=2}^{\delta} d_i$$
$$c_1 = \sum_{i=2}^{\delta} d_i,$$
$$c_1 \le d_k(\delta - 1).$$

But then,  $c_k \ge d_k(\delta - 1) \ge c_1$  which contradicts the fact that  $c_k = c_1 - 1$ . Therefore, it must be that  $c_1 = c_i$  for all  $i = 2, ..., \delta$  and as before G is  $\delta$ -regular.

To complete the classification of  $P_4$ -continuous graphs of girth other than three, we conclude this section with a closer look at regular graphs.

**Lemma 2.3.** Let  $n \ge 4$  be a positive integer. Let G be an r-regular connected graph with  $g(G) \ge n-1$ . Then, for every  $v \in V(G)$ ,

$$P_n \operatorname{-deg}(v) = \frac{nr(r-1)^{n-2}}{2} - n C_{n-1} \operatorname{-deg}(v)$$

where  $C_{n-1}$  is the cycle on n-1 vertices.

Proof. Fix  $v \in V(G)$ . If  $g(G) \ge n$ , then  $C_{n-1}$ -deg(v) = 0, and for v to be at position i of the path,  $1 \le i \le n$ , we have r choices for the first edge incident to v and r-1 choices for each additional edge of  $P_n$ . Finally, since there are n possible positions for v and since  $P_n$  is symmetric, the result follows. If g(G) = n - 1, we are counting illegitimate copies of  $P_n$  whenever v lies on a copy of  $C_{n-1}$ . Moreover, every such false copy of  $P_n$  is counted exactly n times.

**Corollary 2.4.** Let  $n \ge 4$  be a positive integer. A regular connected graph G with  $g(G) \ge n-1$  is  $P_n$ -continuous if and only if it is  $C_{n-1}$ -regular.

**Corollary 2.5.** An r-regular graph G with girth g(G) > 3 is always  $P_4$ -continuous and, in fact,  $P_4$ -regular of degree  $2r(r-1)^2$ . An r-regular graph with girth equal to 3 is  $P_4$ -continuous if and only if it is  $C_3$ -regular. There does not exist a regular  $P_4$ -continuous graph that is not  $P_4$ -regular.

**Open Problem 2.6.** Determine all  $P_4$ -continuous graphs with girth 3 and minimum degree at least 2.

## 3. F-CONTINUOUS GRAPHS AND REGULAR GRAPHS

In this section we examine F-continuous and F-regular graphs for a general graph F. Using a counting argument similar to the one used in the proof of Lemma 2.3 we can consider the case when F is any tree.

**Lemma 3.1.** Let T be a tree with diameter  $d(T) = d \ge 3$  and let G be an r-regular connected graph with  $g(G) \ge d + 1$ . Then, G is T-regular.

Proof. Fix  $v_0 \in V(G)$ . When  $v_0$  is contained in a copy T' of T in G,  $v_0$  is identified with a vertex t' of T'. Think of T' as a rooted tree with root t' and say that  $v_0$  lies in a copy of T in position T'. There exists a set of rooted trees  $T_1, T_2, \ldots, T_a$  that satisfy

- 1.  $T_i$  is isomorphic to T as undirected graphs for i = 1, 2, ..., a, and
- 2. For any graph H, and any vertex  $v \in V(H)$ ,

$$T$$
-deg<sub>H</sub> $(v) = \sum_{i=1}^{a} n_i(H, v)$ 

where  $n_i(H, v)$  denotes the number of times v lies in a copy of T in H in position  $T_i$ .

The integer a depends only on the structure of T. In the case of the tree  $P_4$ , for example, a = 3 and Figure 2 shows the set of three rooted trees.



Figure 2. Set of three rooted trees for  $P_4$ 

We want to show that T-deg<sub>G</sub> $(v_0)$  is constant. For each i = 1, 2, ..., a,

$$n_i(G, v_0) = \binom{r}{\deg_{T_i}(t_i)} \prod_{u \in V(T_i), \ u \neq t_i} \binom{r-1}{\deg_{T_i}(u) - 1}$$

where  $t_i \in V(T_i)$  denotes the root of  $T_i$ . The correctness of the counting argument is guaranteed by  $g(G) \ge d+1$  which is large enough to never mistake a cyclic graph in G for a copy of T. Therefore,  $n_i(G, v_0)$ , i = 1, 2, ..., a, is a function of r and the structure of T showing that T-deg<sub>G</sub> $(v_0)$  will remain the same irrespective of the choice of vertex  $v_0$  of G.

We make use of the following result of Erdös and Sachs, the proof of which can be found in [4].

**Lemma 3.2** [4]. For every two integers  $r \ge 2$  and  $g \ge 3$ , there exists an *r*-regular graph *G* with g(G) = g.

The next theorem solves an open problem posed in [3].

**Theorem 3.3.** Given any nontrivial connected graph F other than the star  $K_{1,k}$ ,  $k \ge 1$ , there exists a regular graph that is not F-continuous.

Proof. Chartrand et al. in [3] have resolved the case of 2-connected graphs F. It will suffice, then, to construct a regular graph with the desired property for any other possible F, falling into two categories.

Case 1: F is a tree.

Let  $d(F) = d, d \ge 3$ , and |V(F)| = n. Note that d < 3 implies that F is a star graph. As hinted by Lemma 3.1, the idea is to construct a regular graph of girth dwhich contains exactly one copy of  $C_d$ . We start with a copy of F to avoid designing a regular graph that is trivially F-continuous because all of its F-degrees are zero. Pick two vertices x and y of F distance d apart and let the path P, passing through vertices  $x, v_1, v_2, \ldots, v_{d-1}, y$  in that order, be a path of length d. Denote by  $\Delta$  the highest degree of a vertex in F and set  $r = 4\Delta$ . We will construct an r-regular graph that is not F-continuous.

Attach a single cycle  $C_d$  to the vertex x of a copy of F by identifying x with a vertex on the cycle. Each vertex of this new graph, that we will call H, has a degree less than or equal to  $\Delta + 2 < r$ . Collectively, for the vertices in the copy of F we need additional

$$nr - 2 - \sum_{v \in V(F)} \deg_F(v)$$

edges, in order to make all of them have degree r in the new graph we are creating. For the vertices in the cycle  $C_d$ , excluding x, we need (d-1)(r-2) more edges. Note that

$$(d-1)(r-2) + nr - 2 - \sum_{v \in V(F)} \deg_F(v) = r(n+d-1) - 2d - \sum_{v \in V(F)} \deg_F(v) = 2q$$

is even since r is even. By Lemma 3.2, there exists an r-regular graph J with g(J) = d + 1. Take q distinct copies of J and remove the same edge  $\{s,t\}$  in each copy. Then, glue each of those graphs to H by adding the edges  $\{s,u\}$  and  $\{t,w\}$ , where u and w are vertices of H in such a way that will complete the degree of each vertex to r. Denote the new r-regular graph by G. Certainly G contains no cycles of length less than d + 1 except the single cycle  $C_d$  we started with. If G' is any r-regular graph with  $g(G') \ge d + 1$ , Lemma 3.1 will imply that F-deg<sub>G'</sub>(v) = A for all  $v \in V(G')$  and some positive constant A. Since we have the cycle  $C_d$  in G, however, the F-degree of some vertices of G will be less than A since A would count some cyclic graphs as copies of F.

Consider the adjacent vertices x and  $v_1$  of G. Despite the edges we added to H,  $v_1$  does not lie directly on the cycle  $C_d$ , and therefore, no double counting will occur and F-deg<sub>G</sub> $(v_1) = A$ . However, the same counting procedure applied to F-deg<sub>G</sub>(x) will consider the cycle  $C_d$  as an acyclic path of length d at least twice, once in either direction. Then, F-deg<sub>G</sub> $(x) \leq A - 2$ , making the F-degrees of x and  $v_1$  differ by more than 1; we have shown that G is not F-continuous.

Case 2: F is not a tree.

Let c(F) = c and say that F has m cycles  $C_c$ . For each  $v \in V(F)$ , define the proximity of v in F, denoted  $\operatorname{prox}_F(v)$ , to be the length of a shortest path from v to a vertex on any of the m cycles  $C_c$  in F. If v lies on one of the m cycles, then  $\operatorname{prox}_F(v) = 0$ . Also, let

$$p = \max\{\operatorname{prox}_F(v) \colon v \in V(F)\}.$$

Identify two copies  $F_1$  and  $F_2$  of F at the same vertex x, where  $\operatorname{prox}_F(x) = p$ . Add an additional vertex y and the edge  $\{x, y\}$ , and denote the resulting graph by H. Let r be the largest degree of a vertex in H. Using H, we will construct an r-regular graph G that is not F-continuous. In particular, our goal is to make F-deg<sub>G</sub>(y) = 0while F-deg<sub>G</sub> $(x) \ge 2$ .

Using Lemma 3.2, there exists an r-regular graph J of girth  $g = \max \{c+1, p\}$ . Note that

$$\sum_{e \in V(H), u \neq x, y} \left( r - \deg_H(u) \right) = 2 \sum_{u \in V(F), u \neq x} \left( r - \deg_F(u) \right) = 2q$$

u

is even. Let  $J_1, J_2, \ldots, J_q$  be q disjoint copies of J. Remove the same edge, say  $\{s, t\}$  from each copy. Then, glue each copy  $J_i - \{s, t\}$  to H by adding the edges  $\{s, v_1\}$  and  $\{t, v_2\}$  where  $v_1$  is a vertex of  $F_1$ ,  $v_2$  is the corresponding vertex of  $F_2$ , and  $v_1, v_2 \neq x$ . Continue to glue the copies of J until all vertices of H, except possibly x and y, have degree r.

Next, we deal with the vertices x and y. Let  $b = \deg_F(x)$ . Note that  $\deg_H(x) = 2b + 1$  and  $\deg_H(y) = 1$ . Take r - (2b + 1) more copies of J, remove the same edge  $\{s, t\}$ , and attach each copy to H by adding the edges  $\{x, s\}$  and  $\{y, t\}$ . In the graph we have constructed so far, all vertices will have degree r, except possibly the vertex y that will have degree r - 2b. So, finally, take b copies of J, remove the same edge  $\{s, t\}$ , and glue each copy to our graph by the edges  $\{y, s\}$  and  $\{y, t\}$ . Denote the final graph by G. Certainly G is r-regular and the only cycles  $C_c$  in G are the 2m such cycles in  $F_1$  and  $F_2$ . Furthermore, since  $F_1$  and  $F_2$  contain the vertex x, it is clear that F-deg<sub>G</sub> $(x) \ge 2$ . We are left to show that F-deg<sub>G</sub>(y) = 0.

Assume on the contrary that y is contained in a copy F' of F, where F' is a subgraph of G. Then, F' must contain m of the 2m cycles  $C_c$  in G. By our definition of p,  $\operatorname{prox}_{F'}(y) \leq p$ .

However, if we remove the vertex x from G, G is no longer connected, and all of the cycles of type  $C_c$  will lie in a different component than the vertex y. Also, since  $g(J) \ge p$ , the shortest distance from x to a cycle  $C_c$  in G remains p. That is, any shortest path from y to a cycle  $C_c$  must start with the edge  $\{y, x\}$  and continue with a path from x to a cycle  $C_c$ . Thus,

$$\operatorname{prox}_G(y) \ge 1 + p$$

which is impossible because  $\operatorname{prox}_{F'}(y) \ge \operatorname{prox}_G(y)$ . Therefore, x and y are adjacent vertices of G whose F-degrees differ by more than 1; G is not F-continuous.

Chartrand et al. in [3] pose yet another open problem concerning regular graphs. They question whether for every nontrivial connected graph  $F, F \neq K_{1,k}$  for  $k \ge 1$ , there exists a regular F-continuous graph which is not F-regular. In [3] they answer this question in the affirmative if F is any nontrivial complete graph  $K_n$ . Here, we show that the answer is still affirmative if F is any 2-connected graph - a graph which remains connected after removing any two of its vertices and their adjacent edges.

**Theorem 3.4.** For every nontrivial 2-connected graph F, there exists a regular F-continuous graph that is not F-regular.

Proof. Let c(F) = c and take two disjoint copies  $F_1$  and  $F_2$  of F. Add a new vertex y and two new edges  $\{y, x_1\}$  and  $\{y, x_2\}$ , where  $x_1$  is a vertex in  $F_1$  and  $x_2$  is

the corresponding vertex in  $F_2$ . Denote the graph constructed so far by H. If  $\Delta(H)$  is the largest degree of a vertex in H, let  $r = 4\Delta(H)$ . We will add edges and vertices to H to convert it to an r-regular graph. Observe that

$$r-2+2\bigg(\sum_{v\in V(F)}r-\deg_F(v)\bigg)-2=2q$$

is even since r is even. Using q disjoint copies of an r-regular graph J with g(J) = c+1we can transform H into an r-regular graph G with girth c using the same approach as in the proof

of Theorem 3.3. The only cycles of length c in G would be the ones in  $F_1$  and  $F_2$ . This and the fact that F is 2-connected guarantees that  $F_1$  and  $F_2$  are the only copies of F in G. Then, F-deg<sub>G</sub>(y) = 0 while F-deg<sub>G</sub> $(x_1) = F$ -deg<sub>G</sub> $(x_2) = 1$  and there is no vertex in G that is contained in both  $F_1$  and  $F_2$ . Therefore, G is not F-regular but it is F-continuous.

When F is not 2-connected, however, the same result does not necessarily hold. In particular, when  $F = P_4$  there does not exist a regular  $P_4$ -continuous graph that is not  $P_4$ -regular as seen in Corollary 2.5.

**Open Problem 3.5.** For every integer  $n \ge 5$ , does there exist a regular  $P_n$ continuous graph that is not  $P_n$ -regular?

**Open Problem 3.6.** Given any nontrivial connected graph F that is not 2-connected, does there exist a regular F-continuous graph that is not F-regular?

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